

DATA HIDING IN GRAYSCALE IMAGES BY DYNAMIC PROGRAMMING BASED ON A HUMAN VISUAL MODEL

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ABSTRACT

A new method for data hiding in grayscale images based on a human vision model with distortion-minimizing capabilities is proposed for covert communication. Each of the eight bit planes of an input grayscale image is viewed as a binary image, into which message data are embedded horizontally. To minimize the image distortion, two optimization techniques, namely, block pattern coding and dynamic programming, are proposed. Furthermore, the proposed method can predict the PSNR value of the resulting image according to the size of the data to be embedded before the embedding process starts. Experimental results show good performances of the proposed method.

1. INTRODUCTION

Data hiding in images is a useful technique for covert communication. Many data hiding techniques have been proposed recently [1, 2, 3]. The most common approach is least-significant-bit (LSB) replacement, which embeds message data in a subset of the LSB planes of an image. The image into which a message is hidden is called a *cover image*, and the result a *stego-image*. Wang et al. [4] embedded a binary image in the fifth LSB bit plane of a cover image, and employed an optimal substitution process based on a genetic algorithm and a local pixel adjustment method to lower the distortion in the stego-image. Chang et al. [5] used dynamic programming to obtain an optimal solution for the LSB substitution method. Chan and Cheng [6, 7] presented an optimal pixel adjustment process to improve the image quality of the stego-image acquired by Wang's schemes. Thien and Lin [8] proposed a method for hiding data in images digit by digit using a modulus function. The method is better than simple LSB substitution not only in eliminating false contours but also in reducing image distortion. Lee and Chen [9] applied variable-sized LSB insertion to estimate the maximum embedding capacity by a human visual system (HVS) property, and to maintain image fidelity by removing false contours in smooth image regions. Liu et al. [10] presented a novel bit plane-wise data hiding scheme using variable-depth LSB substitution and employed post-processing to eliminate the resulting noticeable artifacts.

Most of the above methods lack consideration of using precise human visual models in improving the data hiding effect. Instead, Wu and Tsai [11] presented a method based on the HVS by modifying quantization scales according to variation insensitivity from smooth to contrastive to improve the stego-image quality. And Lie and Chang [12] presented an adjusted LSB technique with the number of LSBs adapting to the pixels of different grayscales.

In this study, we propose a method to embed data into a grayscale image for covert communication, based on the use of a new HVS model to estimate the number of usable bits of each pixel in the cover image. Furthermore, a block pattern encoding method is proposed to embed up to three data bits in a 2×2 block of the bit planes without yielding visible degrading of the stego-image quality. This is achieved by using two optimization techniques. The first technique utilizes multiple block pattern encoding tables, from which an optimal one is chosen for each input image, and the second technique uses dynamic programming to divide the message data stream into appropriate bit segments for optimal data bit embedding in the image blocks to minimize a cost function. Especially, the proposed method can predict the PSNR value of the stego-image according to the embedded data size before the embedding process is started. Moreover, the proposed method can extract embedded data without referencing the original image, and does not require post-processing to refine the stego-image quality.

In the remainder of this paper, we introduce the basic idea behind the proposed method in Section 2. In Section 3, we describe the adopted HVS model and the corresponding cost function. In Section 4, the proposed data hiding method is described. The corresponding data recovery process is proposed in Section 5. Some experimental results are given in Section 6, followed by conclusions in Section 7.

2. EMBEDDING DATA IN BIT PLANES OF GRAYSCALE IMAGES

Eight bits represent a pixel's intensity in a grayscale image. The *bit plane* formed by the same bit of each pixel in the grayscale image is a binary image. Figure 1 shows the eight bit planes of each of three given 128×128 grayscale

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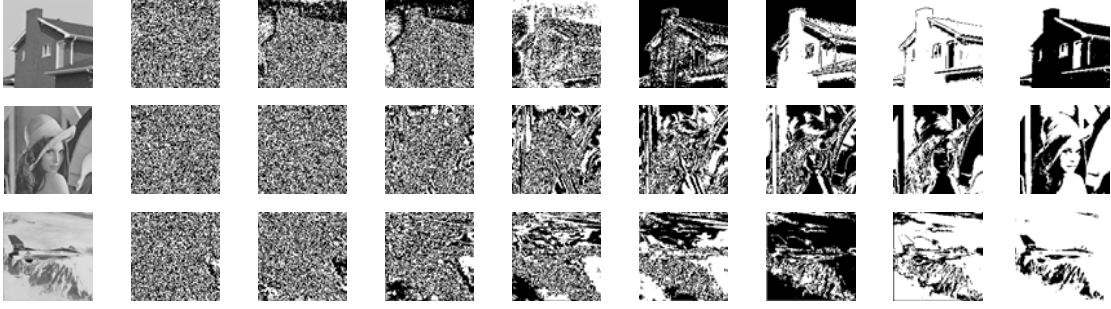


Figure 1. Three grayscale images and their 8 corresponding bit planes (from left to right, original images, bp_0 , bp_1 , bp_2 , ..., and bp_7 , respectively).

images. The image of each bit plane is zoomed out for comparison. It is observed that the content of the LSB plane bp_0 is almost fully randomized. If the message is embedded in bp_0 , the resulting image will appear to be almost unaltered to human eyes. On the contrary, there exist less random noise areas in the remaining more significant bit planes.

The most-significant-bit plane bp_7 contains almost no noise, and data cannot be embedded easily into it without causing significant visual changes. We may embed message data into bit planes in the order of bp_0 , bp_1 , ..., bp_7 . This scheme is termed *horizontal data hiding*, to be contrastive with traditional *vertical data hiding* methods which embed data into the bits b_7 , b_6 , ..., b_0 of each pixel in the order of b_0 through b_7 , where b_0 is the LSB of the pixel. Compared with the vertical data hiding method, horizontal data hiding can reduce more distortion in the stego-image, as revealed in the results of this study.

On the other hand, embedding data directly in bit planes will cause visible damages to the edges in the bit planes. To overcome this difficulty, in this study we design a new cost function which considers certain perception characteristics of the HVS, and adopt a method proposed in Lee and Tsai [13] for data embedding. Each bit plane is regarded to have a different weight in its capability for data hiding, and the new cost function is designed accordingly to measure the degree of distortion resulting from pixel value changes. The details are discussed in the following.

3. COST FUNCTION FOR DISTORTION MEASUREMENT

Since stego-images are viewed by human vision, the characteristics of the HVS must be exploited in designing a data embedding process. Two of such characteristics are useful here. First, human perception is more sensitive to grayscale changes in smooth areas than in texture areas in a grayscale image. Second, human perception is sensitive to relative luminance rather than absolute one. Designing the cost function for distortion measurement for data embedding must take these two characteristics into consideration, as elaborated in the following.

3.1 Computing Number of Data-Embeddable Bits with Consideration of Neighborhood Grayscale Value Change

For the first consideration, assume that a pixel P with grayscale value g is to be used to embed message data. Let MAX denote the maximum grayscale value, and MIN the minimum, in the 3×3 block with P as the center, which we call the *neighborhood* of P . Then, the maximum *between-pixel grayscale range* in this block is $\Delta = MAX - MIN$. According to the previous discussions, to avoid a significant change of the

smoothness degree with respect to the neighborhood of P , the new grayscale value g' resulting from the data embedding is restricted in this study to remain in the range of $g \pm \Delta/2$. Then, we define a *maximum number D of data-embeddable bits* at P as

$$D = \lfloor \log_2(\Delta/2) \rfloor = \lfloor (\log_2 \Delta) - 1 \rfloor = \lfloor \log_2(MAX - MIN) - 1 \rfloor. \quad (1)$$

3.2 Computing Number of Data-Embeddable Bits with Consideration of Pixel's Luminance Change

For the second consideration mentioned above, let f denote the luminance of a pixel P with grayscale value g where $1 \leq f \leq 100$. According to the Fechner law [14], the relative luminance property perceived by the HVS may be expressed as a contrast value c computed by

$$c = 50 \times \log_{10} f$$

where $0 \leq c \leq 100$. Moreover, according to the Weber law [14], the maximum allowable change Δc of the contrast value c according to the principle of "just noticeable difference (JND)" about the pixel's luminance change is about 2. That is, if the luminance of a pixel is changed too much so that Δc is larger than 2, the change will be noticeable to the HVS. Accordingly, we can compute in another way a maximum number of data-embeddable bits in the 8 bits of a pixel's grayscale value, as described next.

First, we want to compute the maximum luminance change $(\Delta f)_{max}$ in accordance with the maximum allowable contrast change $(\Delta c)_{max} = 2$. With c being the contrast of pixel P , let c_{max} denote the maximum possible contrast value. Then, we have

$$\begin{aligned} 2 &= (\Delta c)_{max} = c_{max} - c = 50 \times \log_{10} f_{max} - 50 \times \log_{10} f \\ &= 50 \times \log_{10} f_{max} / f, \end{aligned}$$

which can be reduced to be $f_{max}/f = 10^{(2/50)}$. So, the maximum allowable luminance change can be expressed as

$$(\Delta f)_{max} = f_{max} - f = (f_{max}/f - 1) \times f = (10^{(2/50)} - 1) \times f \approx 0.0965 \times f.$$

And so we may impose the following constraint to the value of f :

$$(\Delta f)_{max}/f \leq 0.0965. \quad (2)$$

On the other hand, in a monochrome image the luminance f in the range of [1 100] is represented by the grayscale value g in the range [0 255], such that g may be computed by the mapping $g = (f - 1) \times (255/99) \approx 2.576(f - 1)$, or equivalently, the mapping $f \approx 0.3882g + 1$, which specifies a linear relation between f and g . Hence, from Constraint (2), we can, after some derivations, get the following new constraint for grayscale changes according to the principle of JND:

$$0.0965 \geq (\Delta f)_{\max}/f = (\Delta g)_{\max}/(g + 2.576) \quad (3)$$

where $(\Delta g)_{\max}$, corresponding to $(\Delta f)_{\max}$, denotes the maximum grayscale change in the pixel's neighborhood. That is, if the above constraint (3) is set for data embedding, the changes of grayscales in the stego-image will not be detectable by human eyes according to the JND principle.

Now, we discuss how many bits can be utilized for data embedding for each possible grayscale value g . If 5 bits of the pixel's grayscale are used for embedding message data, the maximum grayscale change at the pixel will be $(\Delta g)_{\max} = 2^5 - 1 = 31$. And according to Constraint (3), g must be larger than 319, which, however, is out of the grayscale range [0, 255]. This means that embedding 5 or more bits of message data into a pixel is impractical according to the JND principle. As a result, bp_4 , bp_5 , bp_6 , and bp_7 are not used for data embedding in this study. If 4 LSBs of g are changed, then $(\Delta g)_{\max} = 2^4 - 1 = 15$, and by Constraint (3) we get $g > 153$. That is, when the following constraint $g > 153$ is satisfied, we can embed data into the 4 LSBs of g without causing a noticeable luminance change according to the JND principle.

However, the binary value of 153 is 10011001_2 . After the 4 LSBs of g are changed, the new value of g might become a value in the range of 10010000_2 through 10011000_2 , which is smaller than 153, causing a violation of Constraint (3). Therefore, we must change the above constraint $g > 153$ to be $g > 160$ where $160 = 10100000_2$ such that after any 4-bit data are embedded into the 4 LSBs of g , the resulting new value g' of g will always be larger than 160, thus satisfying Constraint (3). In other words, to meet Constraint (3), only when a given pixel's grayscale g satisfies $g \geq 160$ can the 4 LSBs of g be replaced by 4-bit message data. And in short, 4 bits are the upper limit to be embedded in a pixel's grayscale according to the JND principle.

Similarly, if 3 bits are changed, then $(\Delta g)_{\max} = 2^3 - 1 = 7$, and by Constraint (3) as well as a similar reasoning process, the constraint $g \geq 72$ should be satisfied, where $72 = 01001000_2$. If 2 bits are changed, the constraint $g \geq 32$ is required, where $32 = 00100000_2$. Finally, if 1 bit is changed, $g \geq 10$ is necessary, where $10 = 00001010_2$.

In summary, we embed an appropriate number B of message bits in a pixel's grayscale g according to the following rule to satisfy the JND principle:

$$\begin{aligned} &\text{if } g \geq 160, \text{ then } B = 4; \\ &\text{if } g \geq 72, \text{ then } B = 3; \\ &\text{if } g \geq 32, \text{ then } B = 2; \\ &\text{if } g \geq 10, \text{ then } B = 1; \\ &\text{otherwise, } B = 0. \end{aligned} \quad (4)$$

3.3 Combining Results of Two Considerations

To combine the results of the above two considerations, it is not difficult to figure out that the maximum number of data-embeddable bits at a pixel should be taken to be $E = \min(D, B)$ where D and B are as specified in (1) and (4), respectively.

Let the grayscale value g of a pixel P in binary form be denoted as $g = (g_7 g_6 g_5 g_4 g_3 g_2 g_1 g_0)_2$, and the replacement cost of g_i in the i -th bit plane be denoted as C_i where $0 \leq i \leq 3$. According to the previous discussions, C_i is defined in this study as:

$$\text{if } i \leq (E - 1), \text{ then } C_i = 8/2^{(E-1)-i}; \text{ otherwise, } C_i = \infty.$$

The above definition of the cost function gives more penalties to the replacement of more significant bits. In more details, we have the following results:

$$\begin{aligned} &\text{if } E = 4, \text{ then } C_0 = 1, C_1 = 2, C_2 = 4, C_3 = 8, C_4 \sim C_7 = \infty; \\ &\text{if } E = 3, \text{ then } C_0 = 2, C_1 = 4, C_2 = 8, C_3 \sim C_7 = \infty; \\ &\text{if } E = 2, \text{ then } C_0 = 4, C_1 = 8, C_2 \sim C_7 = \infty; \\ &\text{if } E = 1, \text{ then } C_0 = 8, C_1 \sim C_7 = \infty; \\ &\text{if } E = 0, \text{ then } C_0 \sim C_7 = \infty. \end{aligned}$$

4. PROPOSED HORIZONTAL DATA HIDING METHOD

The proposed method is implemented as an algorithm which includes two stages: (1) embedding of some *control data*, followed by (2) embedding of message data. The control data include the necessary information for use in the data recovery process. All data are embedded in the bit planes by a block pattern encoding method. As mentioned previously, each of the bit planes bp_0 through bp_3 can be viewed as a binary image and they together can be regarded as being concatenated into a sequence for data embedding. In this section, the idea to deal with the binary image is presented first, followed by the proposed process.

4.1 Block Pattern Encoding for Data Embedding

In order to embed a message into a binary image, every 2×2 image block is regarded as a *pattern* with a 4-bit *binary value* in which each bit of 0 corresponds to a black pixel and each 1 a white one. The proposed data embedding process is based on the use of a *block pattern encoding table* which maps each block pattern into a certain code with one, two, or three bits of the message data to be hidden. And data embedding is accomplished by changing the block bit values so that the corresponding code of the resulting block pattern become just some bits of the input message data to be embedded. A possible block pattern encoding table designed for use in this study is shown in Table 1. It is emphasized, by the way, that such a table is just one of the many possible tables which may be used for data hiding, and the proposed data embedding process will choose from them an *optimal one* for each specific input binary image, as described later.

Suppose that we want to embed one bit in a 2×2 block. The number of possible patterns in a 2×2 block are 16. This number is much larger than the required number of 2 to represent the two different message bits '0' and '1' in a block, so we may use *more than one* block pattern to represent a single message bit (0 or 1), allowing the possibility of choosing among the block patterns an *optimal one* to replace the original block in the data embedding process and thus reducing more distortion in the resulting block. On the other hand, we wish to embed more data in a block, not just a bit as just mentioned; and for this we can use a block pattern to represent more than one bit, as is done in this study. In short, we want to achieve both minimum-cost bit replacement and maximum-volume data embedding.

As an illustration, we may use either the block pattern $t_1 = 1011_2$ or the pattern $t_2 = 0111_2$ to represent the two-bit message value $s = 01_2$. In this way, when we want to embed, for example, the message value $s = 01_2$ into a block B with value $v = 1010_2$, we have the *two alternative* block patterns $t_1 = 1011_2$ and $t_2 = 0111_2$ to choose to replace $v = 1010_2$, instead of the conventional case of *just one*. And if we choose $t_1 = 1011_2$ to replace $v = 1010_2$, then less distortion of just a 1-bit error (occurring at the LSB position) will result. Contrastively, if only one block pattern, say, $t_2 = 0111_2$ is available, then an error of 3 bits will result, causing more distortion in the resulting block. It is such an allowance of multiple choices for

block pattern replacement that achieves more distortion reduction in the proposed method. By the way, the previously-mentioned *bit errors* are used just for convenience of illustrating the advantage of multiple choices of replacing blocks; they in fact should be the *replacement costs* defined previously.

Table 1. A block pattern encoding table proposed in this study.

Type	Block pattern	Corresponding binary value	Encoded message data
0		1111	1
2		1110	00
4		1101	00
6		1011	01
8		0111	01
10		0011	011
12		0101	011
14		1010	010
1		0000	0
3		0001	11
5		0010	11
7		0100	10
9		1000	10
11		0110	100
13		1001	101
15		1100	010

4.2 Data Embedding in Binary Images

The proposed data embedding process in binary bit-plane images consists of four major steps and includes two folds of distortion minimization, as described in the following.

- (1) *Computing bit costs for data embedding:* We calculate the replacement cost value for each bit in the image according to the cost function defined in Section 3.
- (2) *Dividing the input image into blocks:* We first divide each of bit planes bp_0 through bp_3 into non-overlap 2×2 blocks with every two neighboring blocks separated by a 1-pixel-wide line of pixels in between, as shown in Figure 2. And next, we select the first n “embeddable” blocks and concatenate them sequentially, where n is the length of the message data string to be embedded. A block is said to be *embeddable* in this study if the replacement cost value of any one bit of the block is not infinite.
- (3) *Using multiple block pattern encoding tables for first-fold distortion reduction:* We generate all possible block pattern encoding tables and select an *optimal* one for use in the data embedding process, in the sense of introducing the least distortion. The reason is that a single block pattern encoding table will not be suitable for every input binary image; if an image is destroyed seriously after data embedding using a specific table like Table 1, it will be appropriate to use another table with other combinations of block patterns to encode the message data. Specifically, we exchange the encoded message data of certain types in Table 1 with those of the other types in the following way:

- exchange the message data “0” with “1”;
- exchange the message data “00” with “01”;
- exchange the message data “10” with “11”;
- exchange the message data “010” with “011”;
- exchange the message data “100” with “101”;
- exchange the message data “00” and “01” with “10” and “11,” respectively;
- exchange the message data “010” and “011” with “100” and “101,” respectively.

By enumerating all possible cases in the above way, we can get the 128 distinct tables (numbered from 0 to 127) for selection to minimize the distortion.

- (4) *Applying search techniques for second-fold distortion reduction:* Finally, we apply the dynamic programming technique to segment the input message data stream *optimally* into a series of codes and embed them in the input image, according to the cost function proposed previously. This reduces the resulting distortion further in a global sense.

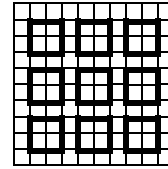


Figure 2. Division of an input image into 2×2 blocks with separating lines (grids with bold boundaries are 2×2 blocks for data embedding)

4.3 Search for Optimal Solutions

The *search cost* proposed in this study for use in the adopted search technique is the *total replacement cost* in the resulting stego-image, computed from the summation of the replacement costs of all the bit changes in the replaced blocks. In Table 1, block patterns can be used to encode one, two, or three message bits. Accordingly, when we embed a binary message value v , we have the three choices of embedding the first one, two, or three initial bits of v into a block. To determine how many bits should be embedded in a selected block, we may calculate first the cost for each of the three cases, and replace the selected block with the block pattern corresponding to the minimum cost. This method provides a quick way for data embedding; however, it is just a *greedy search* algorithm and in general *does not* yield an optimal solution.

To see this, for example, suppose that the message value v of 011_2 is to be embedded in three selected blocks with patterns $B_1 = 0100$, $B_2 = 0100$, and $B_3 = 1100$ according to Table 1. And as illustrated in Figure 3, suppose also that the costs of replacing the four bits are computed to be 2, 1, 1, and 2 for B_1 ; to be 1, 4, 4, and 1 for B_2 ; and to be 4, 4, 1, and 1 for B_3 . By the above-mentioned greedy search algorithm, we replace $B_1 = 0100$ with the block pattern 0000 of type 1 to embed the initial bit 0 of v . The replacement cost for this block is $2 \times 0 + 1 \times 1 + 1 \times 0 + 2 \times 0 = 1$ because a bit (the second bit) is flipped here with its corresponding cost being 1 and the other bits in the original block are not changed. This cost is a local minimum. Next, we replace $B_2 = 0100$ with the block pattern 0001 of type 3 to embed the last two bits 11_2 of v , and the

replacement cost is $1 \times 0 + 4 \times 1 + 4 \times 0 + 1 \times 1 = 5$. Therefore the total replacement cost for embedding v is $1 + 5 = 6$.

Now, if we do not use the greedy search algorithm at the beginning, and replace instead $B_1 = 0100$ by the block pattern 0101 of type 12 in Table 1 to embed the three bits 011₂ of v directly, then the total replacement cost value will be reduced to be $2 \times 0 + 1 \times 0 + 1 \times 1 + 2 \times 1 = 3$ which is *smaller* than the previously-computed total replacement cost of 6. This shows that there indeed exists at least one solution better than that found by the greedy search algorithm. Figure 3 illustrates the data embedding process for this example. This is also true for many other examples, as found by this study. And so the search of an optimal solution is meaningful, for which the proposed method is *dynamic programming*.

Message value $v = "011"$																
		1	2	3												
Selected block list																
Cost function of corresponding bit		<table border="1" style="font-size: small;"><tr><td>2</td><td>1</td></tr><tr><td>1</td><td>2</td></tr></table>	2	1	1	2	<table border="1" style="font-size: small;"><tr><td>1</td><td>4</td></tr><tr><td>4</td><td>1</td></tr></table>	1	4	4	1	<table border="1" style="font-size: small;"><tr><td>4</td><td>4</td></tr><tr><td>1</td><td>1</td></tr></table>	4	4	1	1
2	1															
1	2															
1	4															
4	1															
4	4															
1	1															
Replacing blocks (by greedy search)																
Hidden bit(s)		0	11													
Cost		1	5	(total replaced cost=6)												
Replacing block (by optimum search)																
Hidden bits		011		(total replaced cost=3)												
Cost		3														

Figure 3. An example of proposed data embedding process.

4.4 Dynamic Programming for Data Embedding

In the proposed dynamic programming algorithm (abbreviated as DPA hereafter), *edit distances* are defined for cost minimization in the search. Assume that the input message data to be embedded are in the form of an n -bit string S_1 with $S_1[i]$ denoting its i th bit. Also, let n 2×2 embeddable blocks be selected as a list in advance for data embedding and expressed as another string S_2 with $S_2[i]$ denoting its i th block. For convenience, let $S_k[i \sim j]$ denote a substring of S_k with bits or blocks $S_k[i]$ through $S_k[j]$, where $k=1, 2$ and $i, j=1, 2, \dots, n$.

Only one type of edit operation, namely, *replacement*, is used in the proposed DPA to specify the image block replacement operations involving S_1 and S_2 in the proposed data embedding process. The edit distance between S_1 and S_2 is defined, according to the previous discussions, as the minimum total replacement cost to transform S_2 into S_1 by editing operations according to a certain block pattern encoding table. Let C be an $n \times n$ cost matrix with its element $C[j, i]$ denoting the minimum total replacement cost to transform a substring $S_2[j \sim n]$ of S_2 into a substring $S_1[i \sim n]$ of S_1 . Then $C[1, 1]$ is the minimum total replacement cost to transform $S_2[1 \sim m]$ into $S_1[1 \sim n]$ (i. e., to transform the substring of S_2 into the entire string of S_1), where $1 \leq m \leq n$. Also, let RC be a cost function with each of its element $RC(j, i,$

$\ell)$ denoting the total replacement cost for replacing the j th block $S_2[j]$ of S_2 with the block pattern which encodes the initial ℓ bits of the substring $S_1[i \sim n]$ of S_1 with $\ell = 1, 2, \text{ or } 3$. By the above definitions, the value $C[j, i]$ is recursively just the minimum of all the possible values of $RC(j, i, \ell) + C[j+1, i+\ell]$, where $\ell = 1, 2$ or 3 . Then, according to dynamic programming, the *minimum search cost* and its corresponding solution may be computed by the following algorithm.

Algorithm 1. Computing minimum search cost for minimizing distortion by DPA.

Input: (1) an n -bit message data string S_1 ; (2) a string S_2 of n selected blocks; (3) a block pattern encoding table T ; (4) an $n \times n$ cost matrix $C[j, i]$, for $i, j = 1, 2, \dots, n$; (5) an $n \times n$ type matrix I with its element $I[j, i]$ used for recording the block pattern in T used for replacing $S_2[j]$ in calculating $C[j, i]$; and (6) an $n \times n$ segmentation matrix N with its element $N[j, i]$ used for recording the number of initial bits of $S_1[i \sim n]$ used in calculating $C[j, i]$.

Output: $C[j, i]$, $I[j, i]$, and $N[j, i]$ for all $i, j = 1, 2, \dots, n$.

Steps:

1. Set all $C[j, i]$ initially to be ∞ for all $i, j = 1, 2, \dots, n$.
2. Starting from $i = n$ and $j = n$, for each pair of (j, i) with $i, j = 1, 2, \dots, n$, perform the following steps.
 - 2.1 If $C[j, i]$ is equal to ∞ , continue the next step (Step 2.2); else increment i and j to calculate the next $C[j, i]$.
 - 2.2 Take $C[j, i]$ to be the minimum of the three replacement costs, $RC(j, i, 1) + C[j+1, i+1]$, $RC(j, i, 2) + C[j+1, i+2]$, and $RC(j, i, 3) + C[j+1, i+3]$; and record the corresponding number of the processed initial bits (1, 2, or 3) of $S_1[i \sim n]$ in $N[j, i]$, and the corresponding type of the used block pattern of T in $I[j, i]$.

In the above algorithm, the number of initial bits of $S_1[i \sim n]$ and the used block pattern type in each recursive step are recorded in matrices N and I , respectively, which are used in the data embedding process, as described next.

Algorithm 2. Data embedding using block pattern encoding tables and DPA.

Input: (1) a grayscale image G ; (2) a secret message data string S_1 with n bits; (3) a control message data string S_c with m bits, including a *table number* T_{opt} (specifying the block pattern encoding table used) with seven bits, followed by a value L_{opt} (specifying the number of selected blocks used) with $m-7$ bits; and (4) 128 block pattern encoding tables.

Output: a stego-image S .

Steps:

1. Compute the cost of each bit of G as mentioned previously.
2. Get a list B_m of m 2×2 embeddable blocks sequentially from the bit planes bp_0 through bp_3 of G in order for embedding the m bits of S_c . Also, get a list B_n of n 2×2 embeddable blocks sequentially for the n bits of S_1 . Let B_m and B_n also include the position information of each selected block.
3. For each block pattern encoding table T among the input 128 ones, with S_1, B_n , and T as input, apply Algorithm 1 to calculate the cost matrix $C[j, i]$, the type matrix $I[j, i]$, and the segmentation matrix $N[j, i]$ for all $i, j = 1, 2, \dots, n$.
4. Find the minimum C_{min} of the 128 values of $C[1, 1]$

- computed in the last step, and set T_{opt} to be the table number of the corresponding block pattern encoding table used in computing C_{min} .
5. Use the block pattern encoding table T_{opt} , the type matrix I_{min} and the segmentation matrix N_{min} corresponding to C_{min} , and the position information of each block in B_n , to embed the string S_1 into bp_0 through bp_3 of G to get the stego-image S .
 6. Set the value L_{opt} to be the number of the blocks used for embedding S_1 in the last step.
 7. Using S_c (including T_{opt} and L_{opt}), B_m , and $T=1$ as input, apply Algorithm 1 to calculate the cost matrix $C[j, i]$, the type matrix $I[j, i]$, and the segmentation matrix $N[j, i]$ for all $i, j = 1, 2, \dots, m$.
 8. Use the block pattern encoding table Table 1, the type matrix I and the segmentation matrix N in the last step, and the position information of each block in B_m , to embed the substring S_c into bp_0 through bp_3 of G to get the final stego-image S .

5. PROPOSED DATA RECOVERY PROCESS

The goal of data recovery is to extract the embedded message data from a stego-image, as described in the following.

Algorithm 3. Message data recovery

Input: a stego-image I' including a message bit stream S .

Output: the message bit stream S .

Steps:

1. Calculate the cost of every bit of G as mentioned previously.
2. Get m 2×2 embeddable blocks sequentially from bp_0 through bp_3 of I' as a list L_m .
3. For each 2×2 block P of L_m , compute the binary value v corresponding to the block pattern, and decode v by looking v up in the block pattern encoding table 1 to get the corresponding encoded message data bits as the data recovery result of P .
4. Concatenate the preface m data bits extracted in the last step into a sequence as a desired control message data S_c .
5. Get the preface 7 data bits of S_c as T_{opt} , and the remaining $m-7$ data bits of S_c as L_{opt} , which specify respectively (1) the optimal block pattern encoding table T_{opt} used in data embedding; and (2) the number of 2×2 blocks of I' used in embedding S in the bp_0 through bp_3 of I' .
6. Also, get L_{opt} 2×2 selected blocks sequentially from bp_0 through bp_3 of I' as a list L .
7. For each 2×2 block P of L , compute the binary value v corresponding to the block pattern, and decode v by looking v up in the block pattern encoding table T_{opt} to get the corresponding encoded message data bits as the data recovery result of P .
8. Concatenate all the data bits extracted in the last step into a sequence as the desired message S and exit.

For security consideration, we encrypt further the control message by a secret key before the data embedding process, and embed the result into bp_0 through bp_3 at bit positions randomly generated with a distinct secret key and a random number generator. The reverse process can be easily performed to get the original control message. The same method is also applied to the message data to get a higher degree of data protection.

6. EXPERIMENTAL RESULTS

Figures 4 and 5 illustrate some experimental results of applying the proposed method. The bit streams of message data in Figures 4 and 5 were generated randomly. The stego-images “House” and “Lena” both of size 256×256 with high PSNR values of 56.88 dB and 68.44dB, respectively, obtained by embedding 16440 bits (about 2KB) message data using the DPA and the optimal block pattern encoding table among the 128 ones are shown Figures 4(b) and 4(d), respectively. The cover images are depicted in Figures 4(a) and 4(c) respectively for comparison. The results show that the proposed method can be applied to embed message data in grayscale images and obtain good-quality stego-images without noticeable artifacts in smooth image regions.

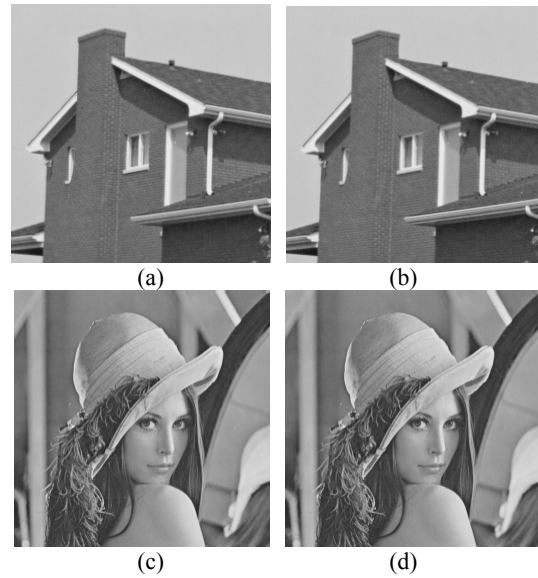


Figure 4. Some experimental results. (a) Cover image “House.” (b) Stego-image of (a). (c) Cover image “Lena.” (d) Stego-image of (c).

Figure 5 illustrates three grayscale stego-images “House”, “Lena” and “Jet.” and results of data embedding. The three stego-images of size 128×128 were obtained by embedding 1000 bytes of message data using the proposed DPA and the optimal encoding table. The PSNR values are 46.90 dB, 49.33dB and 48.80 dB, respectively. In each row of the figure, from left to right are an original image, bp_0 , bp_1 , bp_2 , and bp_3 of the original bit planes (copied from Figure 1), followed by the resulting stego-bit planes of bp_0 , bp_1 , bp_2 , and bp_3 , respectively. Compared with the cover images in Figure 1 and their four corresponding bit planes, it can be seen that the stego-images retain most significant textures.

Table 2 summarizes the statistical data of the stego-image “Lena” using the DPA and the optimal encoding table, including the message data length, the PSNR value, the selected block pattern encoding table, the numbers of used blocks, the minimum replacement cost values, and the average of the numbers of embedded bits per block. The message data bit stream in Table 2 was generated randomly.

In more details, the result of Table 2 is transformed into Fig. 6. When the amount of the embedded data is smaller than 1000B, the PSNR values in Table 2 are all larger than 49dB. And the differences in the stego-images are not noticeable.

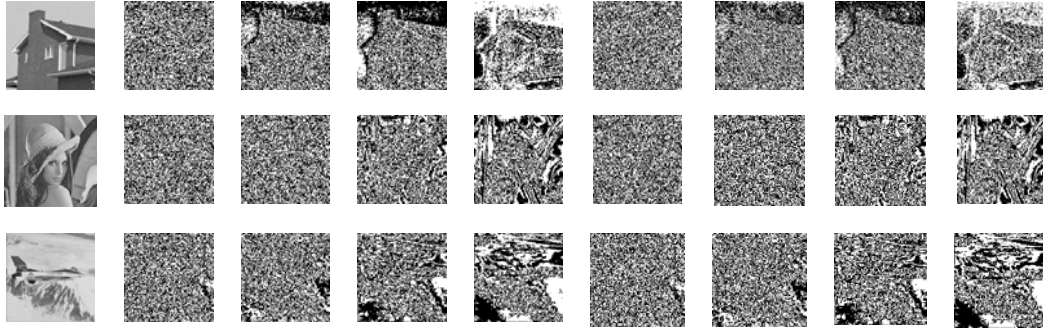


Figure 5. Experimental results of three images. In each row, from left to right are the original image, bp_0 , bp_1 , bp_2 , and bp_3 of original bit planes (copied from Fig. 1), followed by resulting stego-bit planes of bp_0 , bp_1 , bp_2 , and bp_3 .

Figure 6 also reveals that the relation between the PSNR value yielded by the proposed method and the embedded data amount is approximately *linear*. The PSNR value of the DPA decreases about 3.3225 dB when the embedded data size increases 200 bytes. Thus, the proposed DPA method can predict the PSNR value before the data embedding process starts according to the message data size. From Table 2, the PSNR value of the DPA can be estimated by a simple line fitting method to be $PSNR = 62.5870 - (m - 1) \times 3.3225$ (dB), $1 \leq m \leq 5$, where m denotes the size of message data in the unit of 200 bytes for 128×128 grayscale images. Similar results can be observed for the other images. Moreover, the equation of the PSNR value can be extended and used for grayscale images of any sizes by applying the proposed DPA. For a grayscale image of size $H \times W$, if the above value of m denotes the size of the message data in the unit of $(200 \times H \times W) / (128 \times 128)$ bytes, then the resulting PSNR value still can be estimated using the above equation. Note that this merit of predictable PSNR values enables a user of the proposed method to determine in advance how large a cover image should be selected for a certain given amount of message data.

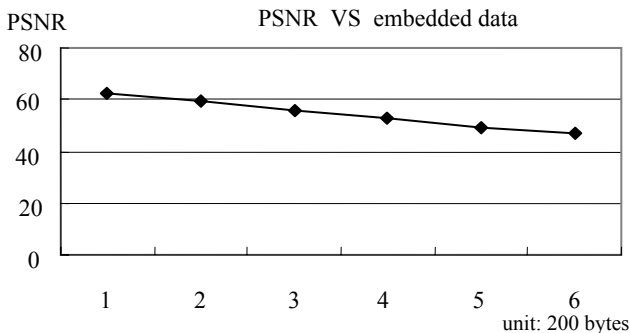


Figure 6. PSNR values of stego-image “Lena” using DPA.

Furthermore, we may compute a *distortion rate* for each stego-image. This rate is computed in this study as the ratio of the number of bit flippings (changing bit 0 to 1 or 1 to 0 in data embedding) to the length of the message data. Most

existing vertical data hiding methods yield distortion rates of about 50% for grayscale images because of the characteristic of randomness in bit flippings. In most existing vertical data hiding methods, when 200-byte secret message data are embedded into a 128×128 grayscale image, these data will be divided into pieces of 4 bits and each piece is embedded into the bits $b_3b_2b_1b_0$ of a pixel. Then the *average* grayscale change of the pixel, measured in terms of the number of flipped bits, may be computed to be $1 \times 50\% + 2 \times 50\% + 4 \times 50\% + 8 \times 50\% = 7.5$. Consequently, the corresponding mean square-error value *MSE* of the stego-image may be computed to be $MSE = [(200 \times 8) / 4] \times (7.5 \times 7.5) / (128 \times 128)$ where $(200 \times 8) / 4$ is the number pixels required for embedding the 200-byte message data, 7.5×7.5 is the square error incurred at each pixel, and 128×128 is the image size. Finally, the PSNR value of the stego-image may be computed to be $10 \times \log(255^2 / MSE) = 46.75$ dB. Contrastively, in Table 2 the PSNR value of our method is a much larger value 62.48 dB, which means the proposed method is superior in distortion minimization.

In summary, the experimental results above show that for grayscale images, the proposed method takes effectively into consideration both the effect of the HVS and the preservation of local edges, and yields stego-images that are less distorted than those yielded by most existing data hiding methods.

Table 2. Statistics of stego-images yielded by DPA using optimal encoding table.

Stego - image	Message data length (bytes)	PSNR (dB)	Table number	No. of used blocks	Cost value	Embedded bit number per block
Lena (128 × 128)	200	62.48	8	779	1774	2.054
	400	59.37	57	1636	3243	1.956
	600	55.57	57	2297	5680	2.09
	800	52.96	57	3101	8355	2.064
	1000	49.33	8	3826	11699	2.091
	1200	46.74	57	4295	16248	2.235

7. CONCLUSIONS

Two novel methods for hiding message data into grayscale images with distortion reduction effects have been proposed, with one being the DPA and the other being the use of multiple block pattern encoding tables. A cost function has been proposed to estimate the weight of each bit in each pixel to be replaced according to an HVS model. A horizontal data hiding method in which message data are embedded in a sequence of bit planes has also been proposed to decrease possible distortions in stego-images. An optimal block pattern encoding table is chosen from 128 alternative ones for use in the proposed data embedding process to minimize the distortion in the stego-image. The proposed method minimizes further the distortion using dynamic programming based on the proposed cost function. The proposed method can embed up to three bits in a 2×2 image block. As a result, not only the textures in the cover image are reserved, but also the distortion of the stego-image can be effectively reduced in an optimal way. Especially, the proposed DPA can predict the PSNR value of a stego-image before the embedding process starts according to the size of the data to be embedded.

The space and time complexities of the proposed DPA are both of the order of n^2 . It may cost more time to embed a long secret code. But in certain applications there is no need of real-time processing, and optimality in data embedding volumes or minimization in image distortion is the main concern. In such cases, our method is good to use. On the other hand, if time is really concerned, then one can alternatively use the other proposed method, the greedy search algorithm, that takes only linear computation time and still minimize distortion in the stego-image in a suboptimal way.

On the other hand, at least two approaches may be adopted to make the proposed method more robust. First, we can embed multiple copies of the secret data in the embedding process so that attacks will not entirely destroy the secret information. If attacks do occur, then after the data are extracted by the proposed method, we may apply a voting scheme to recover the secret. The second approach is to try to place secret data in the more significant bits of the cover image, for example, in bp_2 and bp_3 in the proposed method, assuming that most attacks to BMP images are conducted to the LSBs. Because the information encoded in these bit-planes cannot be removed in most applications (otherwise, the image will be seriously distorted or destructed), hopefully this method will work in real applications.

The proposed method processes 2×2 blocks in the data embedding process. It may be extended to process larger-sized blocks because when the block size is larger, the number of the block patterns which can be selected to encode a certain message value becomes larger as well, resulting possibly in greater reduction of image distortion. Other future works may be directed to embed multiple message data in a grayscale image for protecting the intellectual property right and authenticating multimedia data, to define more general cost functions for other HVS models, and to design better encoding tables to reduce stego-image distortion further.

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