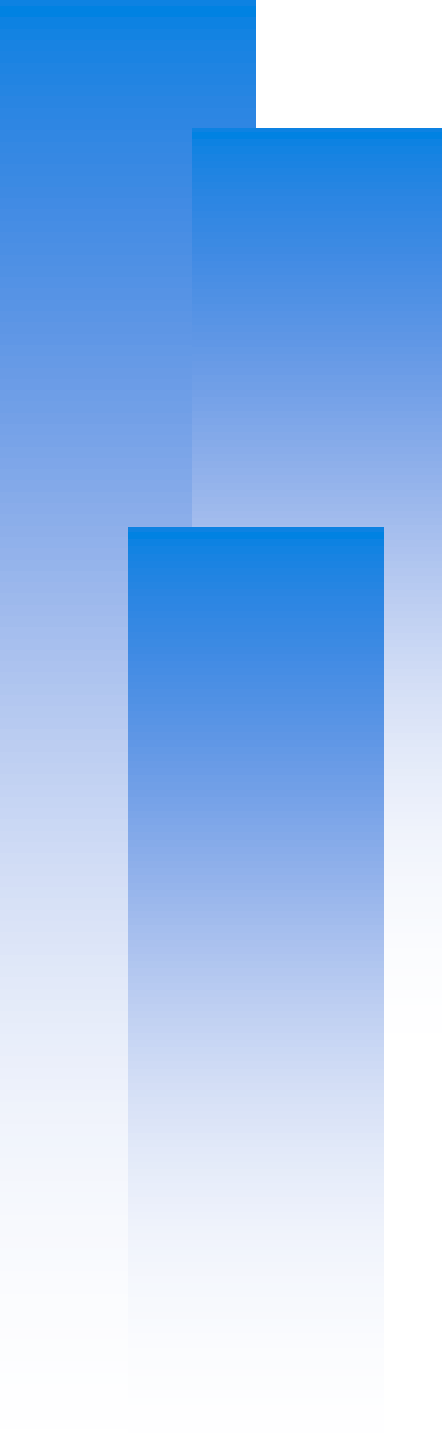




TCP/IP Protocols

TCP/IP and the Internet

- ❑ In 1969
 - ARPA funded and created the “ARPANET” network
 - 高等研究計劃署 (Advanced Research Project Agency)
 - NCP – Network Control Protocol
 - Allow an exchange of information between separated computers
- ❑ In 1973
 - How to connect ARPANET with SATNET and ALOHANET
 - TCP/IP begun to be developed
- ❑ In 1983
 - TCP/IP protocols replaced NCP as the ARPANET’s principal protocol
 - ARPANET → MILNET + ARPANET = Internet
- ❑ In 1985
 - The NSF created the NSFNET to connect to Internet
- ❑ In 1990
 - ARPANET passed out of existence, and in 1995, the NSFNET became the primary Internet backbone network

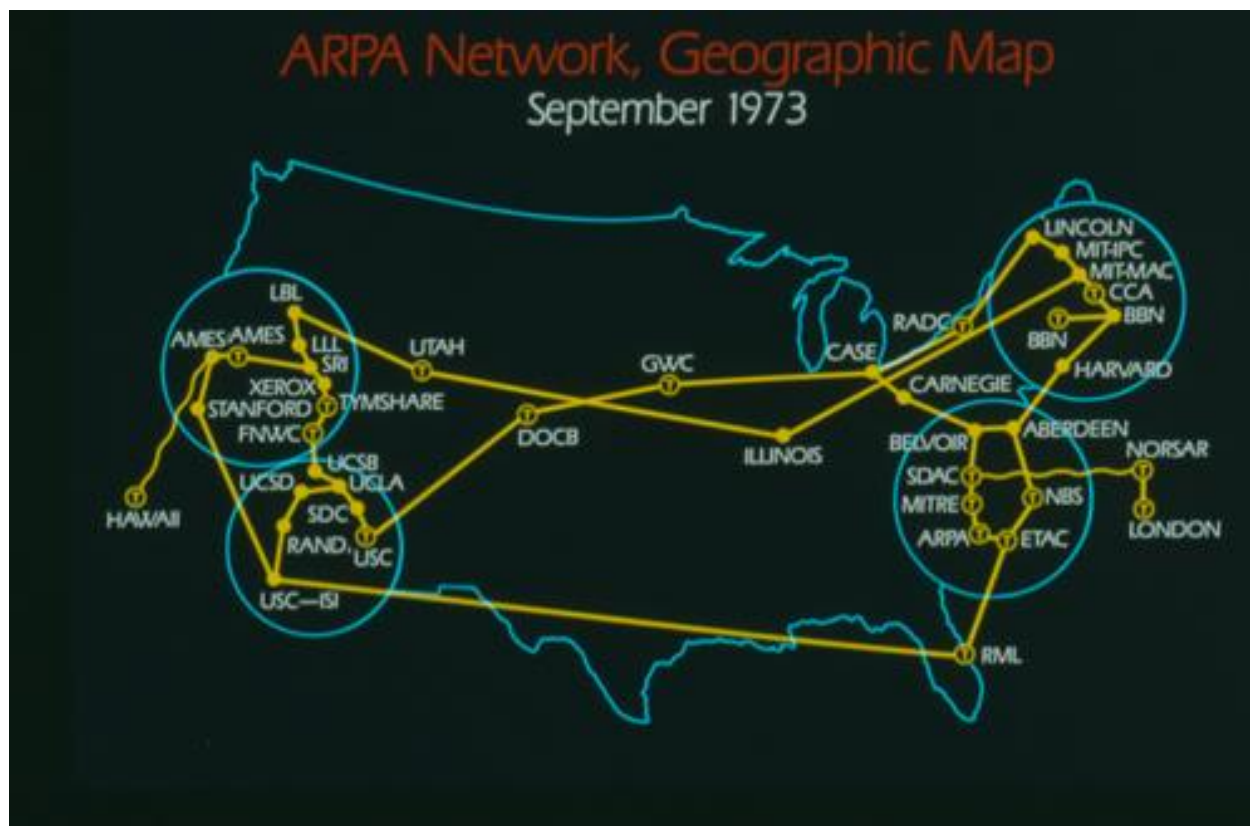
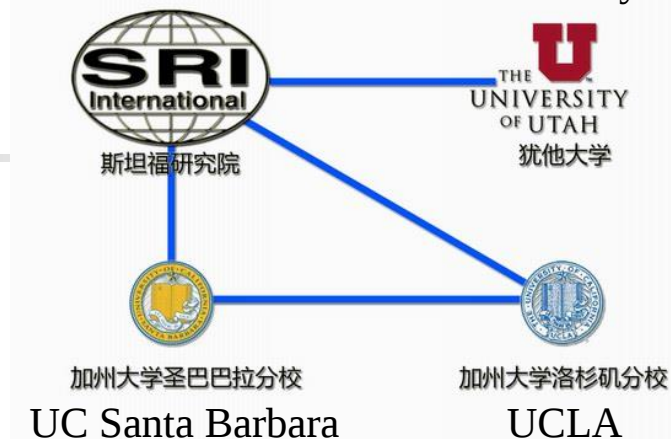
ARPA: Advanced Research Project Agency

NSF: National Science Foundation

Introduction

– APRANET

Stanford Research Institute University of Utah

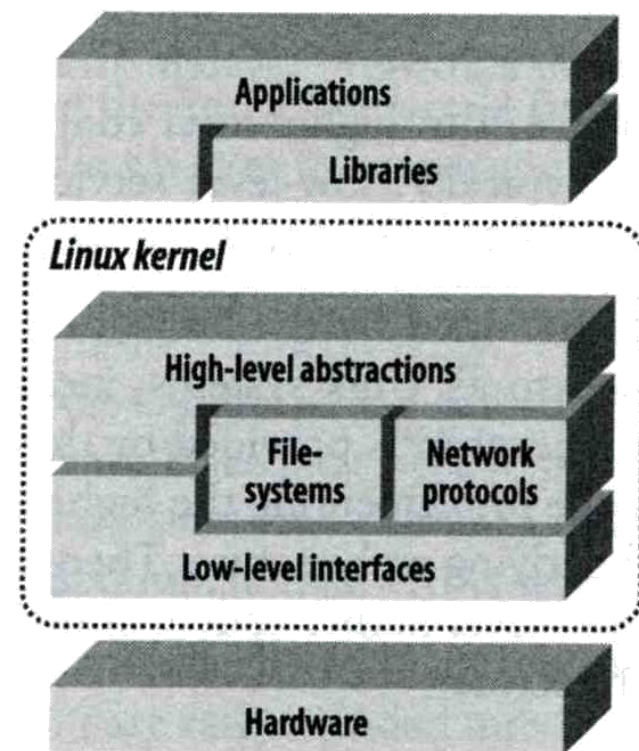


Introduction

– Why TCP/IP ?

❑ The gap between applications and Network

- Network
 - 802.3 Ethernet
 - 802.4 Token bus
 - 802.5 Token Ring
 - 802.11 Wireless
- Application
 - Reliable
 - Performance



We need something to do the translating work!
TCP/IP it is!!

Introduction

– Layers of TCP/IP (1)

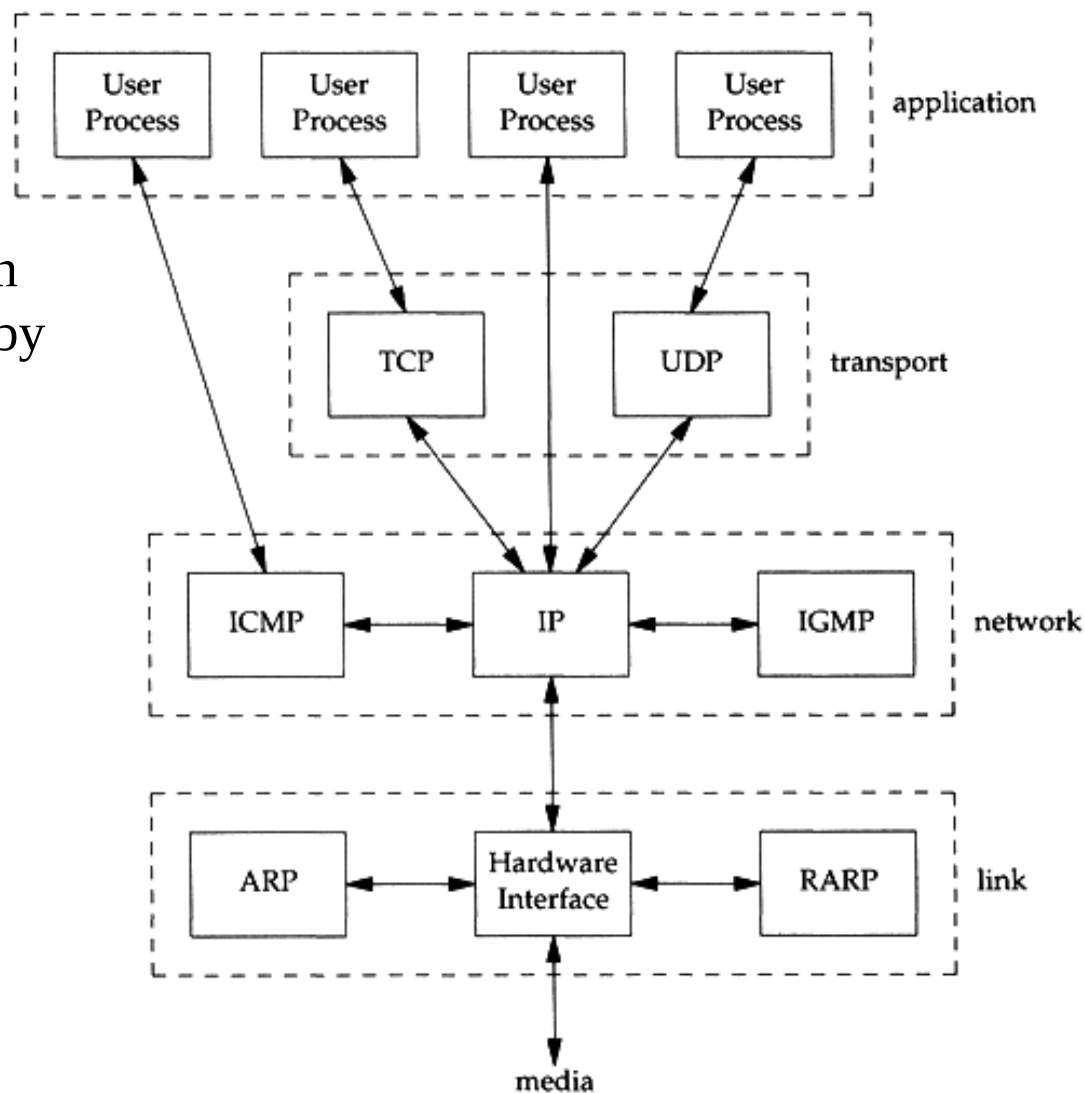
- ❑ TCP/IP is a suite of networking protocols
 - 4 layers Layering architecture
 - Link layer (data-link layer)
 - Include device drivers to handle hardware details
 - Network layer (IP)
 - Handle the movement of packets around the network
 - Transport layer (Port)
 - Handle flow of data between hosts
 - Application

Introduction

– Layers of TCP/IP (2)

❑ Each layer has several protocols

- A layer define a data communication function that may be performed by certain protocols
- A protocol provides a service suitable to the function of that layer

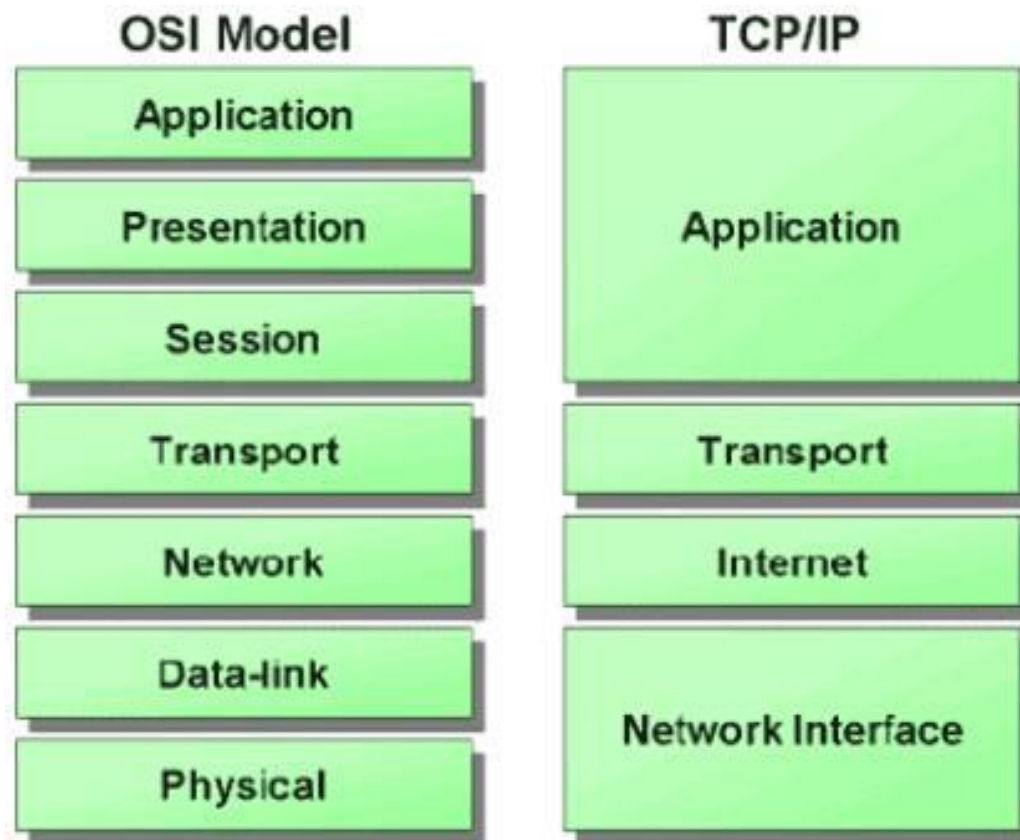


Introduction

– Layers of TCP/IP (3)

☐ ISO/OSI Model

☐ TCP/IP Model



TCP/IP and the OSI model

ISO: International Organization for Standardization

OSI: Open System Interconnection

Introduction

❑ TCP/IP

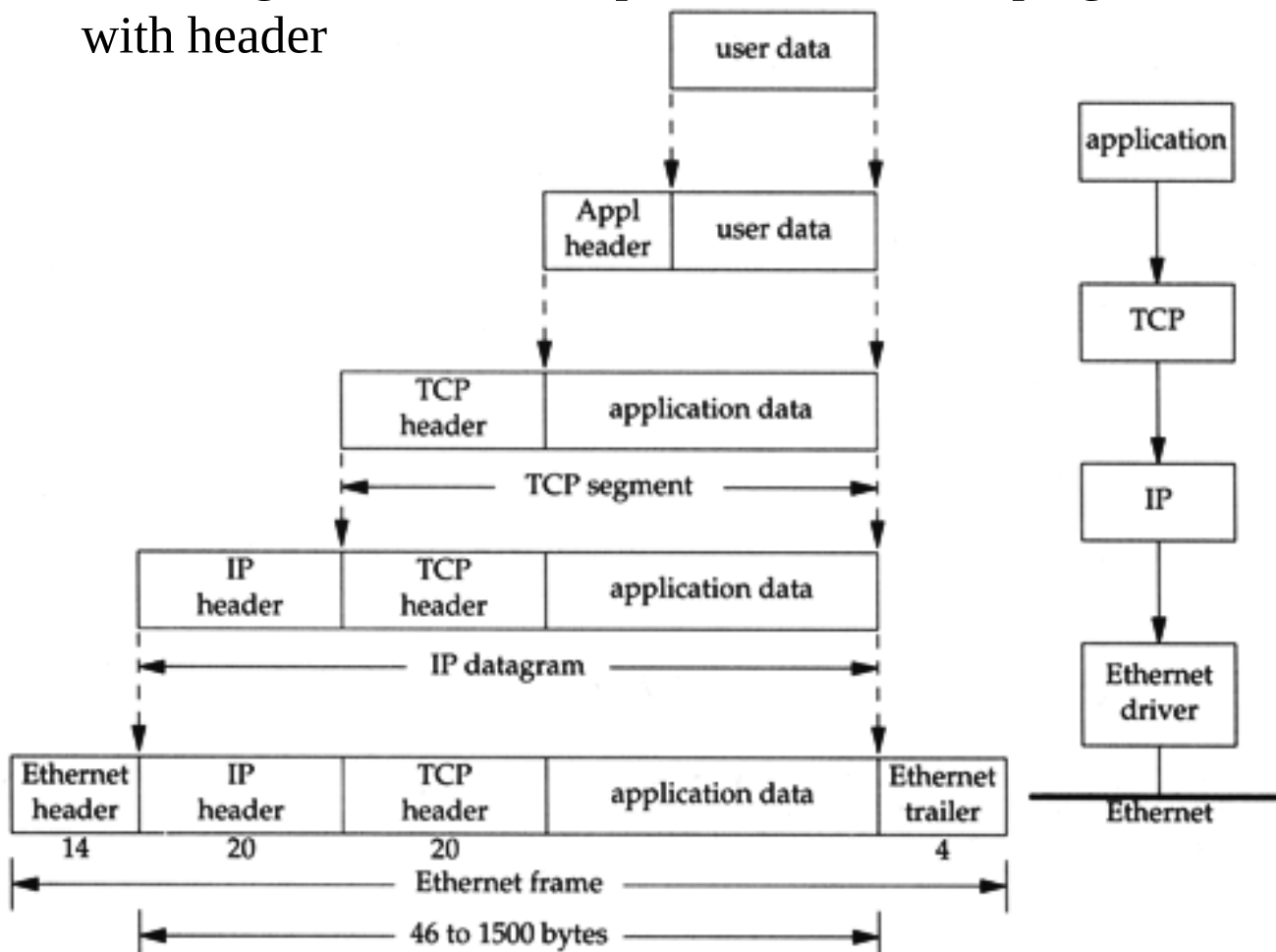
- Used to provide data communication between hosts
 - How to delivery data reliably
 - How to address remote host on the network
 - How to handle different type of hardware device

Introduction

– Encapsulation

❑ Multiplexing

- Gathering data from multiple sockets, enveloping data with header



Introduction

– Decapsulation

❑ Demultiplexing

- Delivering received segments to correct socket

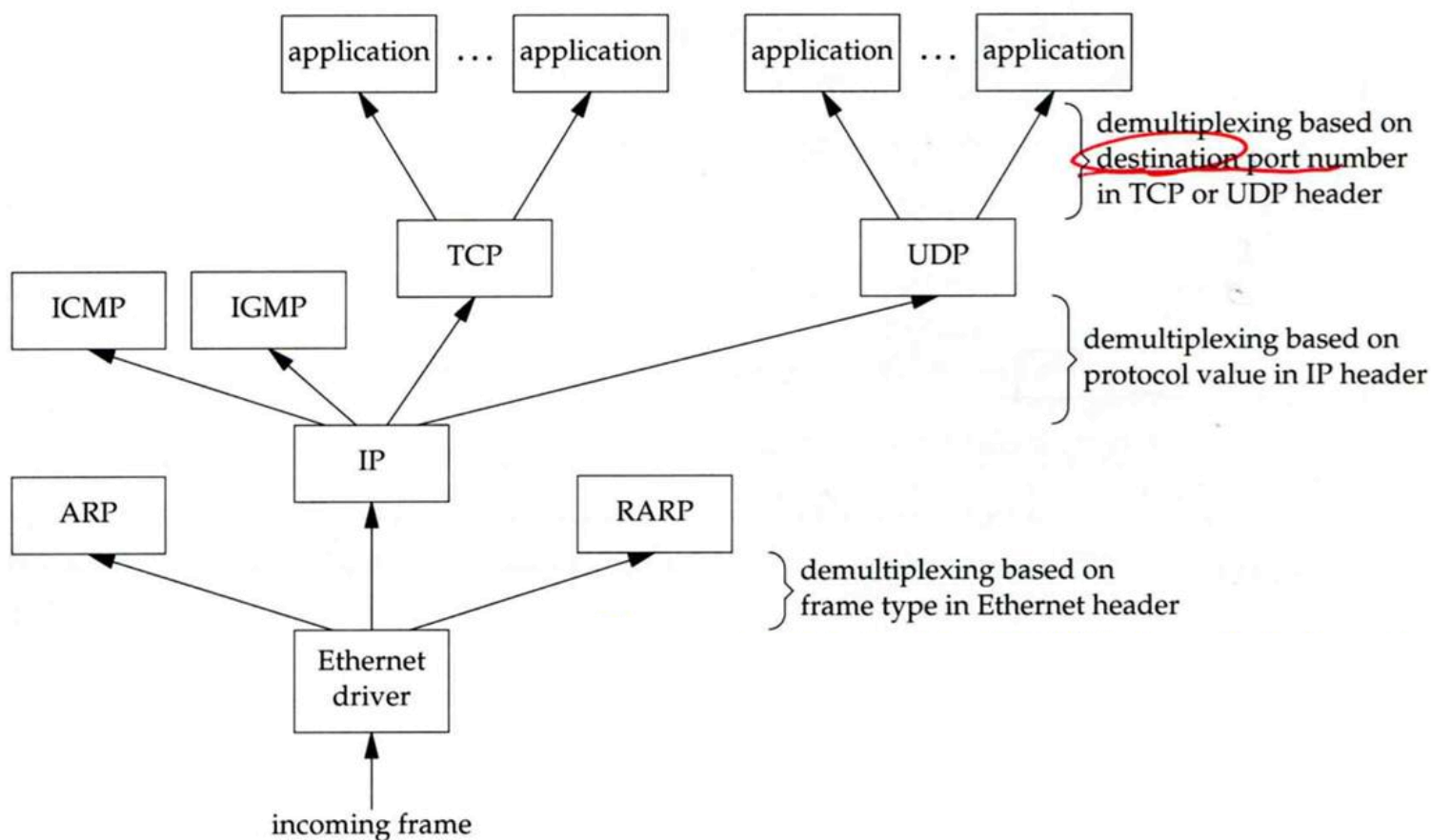


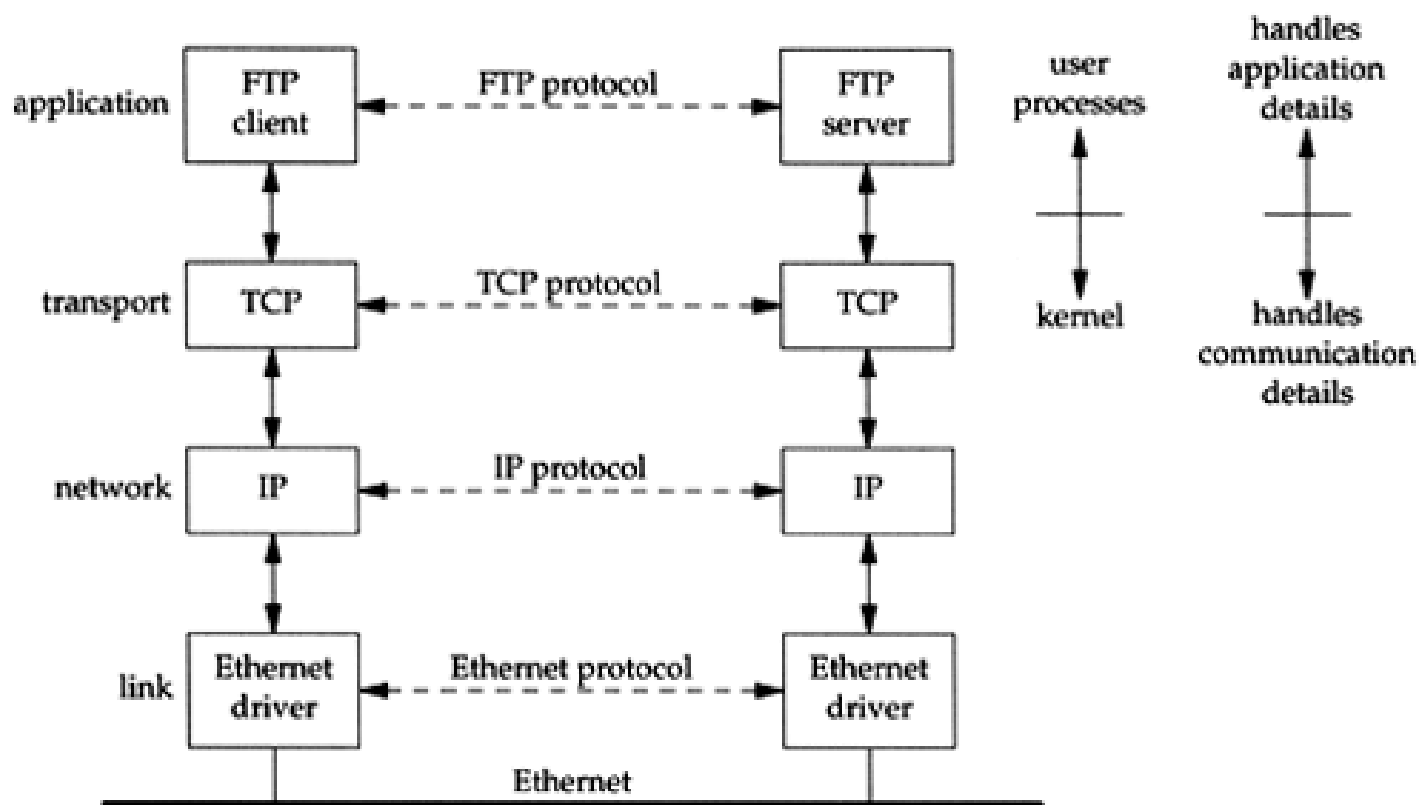
Figure 1.8 The demultiplexing of a received Ethernet frame.

Introduction

– Addressing

□ Addressing

- Nearby (same network)

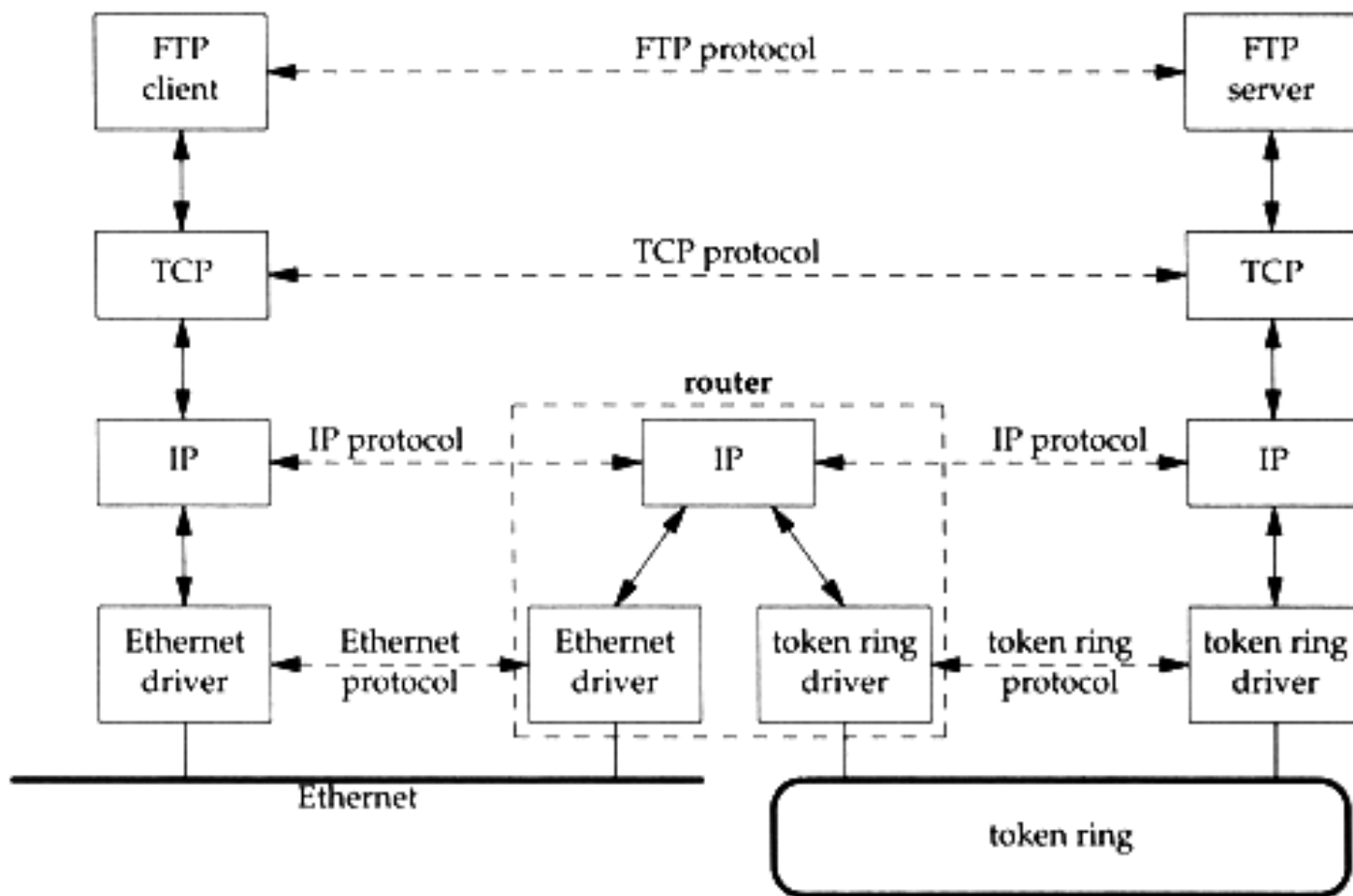


Introduction

– Addressing

□ Addressing

- Faraway (across network)



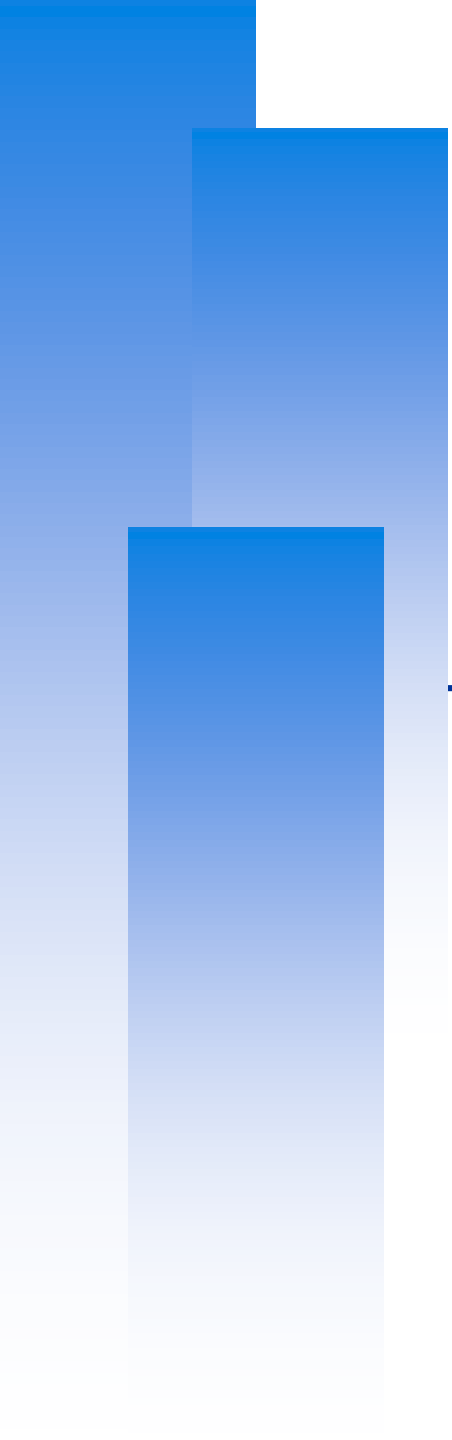
Introduction

– Addressing

□ Addressing

- MAC Address
 - Media Access Control Address
 - 48-bit Network Interface Card Hardware Address
 - 24-bit manufacture ID
 - 24-bit serial number
 - Ex:
 - 00:07:e9:10:e6:6b
- IP Address
 - 32-bit Internet Address (IPv4)
 - Ex:
 - 140.113.209.64
- Port
 - 16-bit uniquely identify application (1 ~ 65536)
 - Ex:
 - FTP port 21, SSH port 22, Telnet port 23

```
sabsd [/home/chwong] -chwong- ifconfig
sk0: flags=8843<UP,BROADCAST,RUNNING,SIMPLEX,MULTICAST> mtu 1500
options=b<RXCSUM,TXCSUM,VLAN_MTU>
inet 140.113.17.215 netmask 0xfffff00 broadcast 140.113.17.255
inet 140.113.17.221 netmask 0xffffffff broadcast 140.113.17.221
ether 00:11:d8:06:1e:81
media: Ethernet autoselect (100baseTX <full-duplex,flag0,flag1>)
status: active
lo0: flags=8049<UP,LOOPBACK,RUNNING,MULTICAST> mtu 16384
inet 127.0.0.1 netmask 0xff000000
```



Link Layer

Link Layer

– Introduction of Link Layer

- ❑ Purpose of the link layer
 - Send and receive IP datagram for IP module
 - ARP request and reply
 - RARP request and reply

- ❑ TCP/IP support various link layers, depending on the type of hardware used:
 - Ethernet
 - Teach in this class
 - Token Ring
 - FDDI (Fiber Distributed Data Interface)
 - Serial Line

Link Layer

– Ethernet

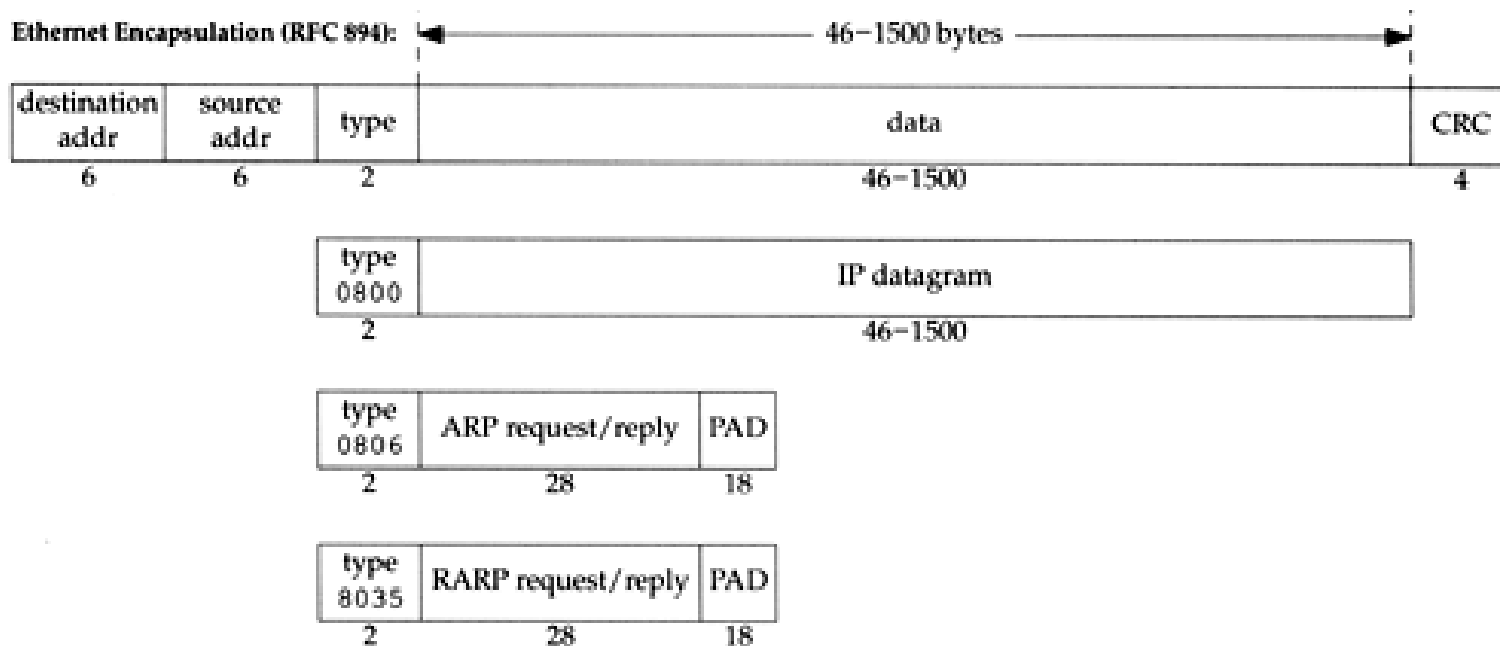
❑ Features

- Predominant form of local LAN technology used today
- Use CSMA/CD
 - Carrier Sense, Multiple Access with Collision Detection
- Use 48-bit MAC address
- Operate at 10 Mbps
 - Fast Ethernet at 100 Mbps
 - Gigabit Ethernet at 1000 Mbps
- Ethernet frame format is defined in RFC 894
 - This is the actually used format in reality

Link Layer

– Ethernet Frame Format

- ❑ 48-bit hardware address
 - For both destination and source address
- ❑ 16-bit type is used to specify the type of following data
 - 0800 → IP datagram
 - 0806 → ARP, 8035 → RARP

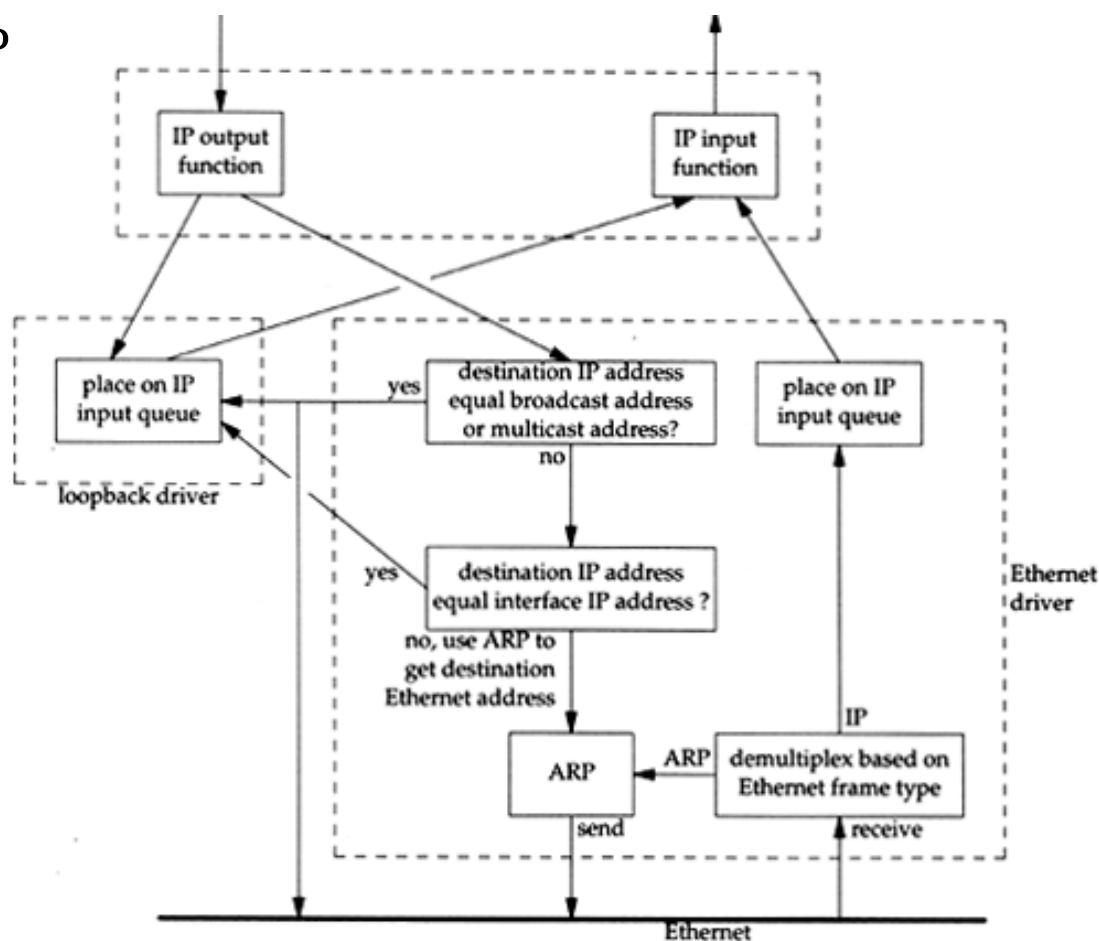


Link Layer

– Loopback Interface

❑ Pseudo NIC

- Allow client and server on the same host to communicate with each other using TCP/IP
- IP
 - 127.0.0.1
- Hostname
 - localhost



Link Layer

– MTU

❑ Maximum Transmission Unit

- Limit size of payload part of Ethernet frame
 - 1500 bytes
- If the IP datagram is larger than MTU,
 - IP performs “fragmentation”

❑ MTU of various physical device

❑ Path MTU

- Smallest MTU of any data link MTU between the two hosts
- Depend on route

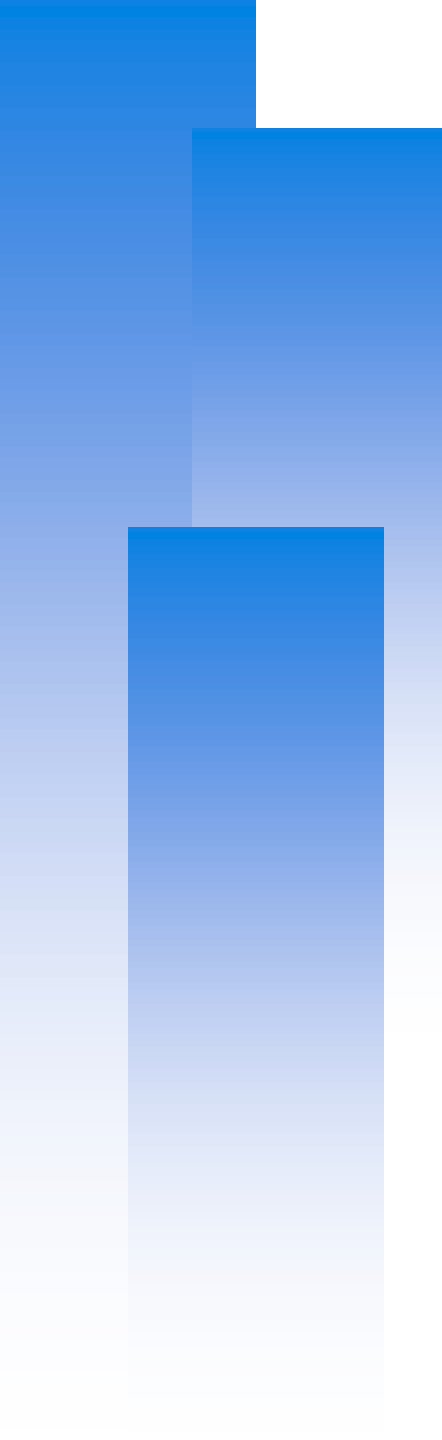
Network	MTU (bytes)
Hyperchannel	65535
16 Mbits/sec token ring (IBM)	17914
4 Mbits/sec token ring (IEEE 802.5)	4464
FDDI	4352
Ethernet	1500
IEEE 802.3/802.2	1492
X.25	576
Point-to-point (low delay)	296

Link Layer

– MTU

- ❑ To get MTU info

```
% ifconfig
em0: flags=8843<UP,BROADCAST,RUNNING,SIMPLEX,MULTICAST> mtu 9000
    options=b<RXCSUM,TXCSUM,VLAN_MTU>
    inet 192.168.7.1 netmask 0xffffffff broadcast 192.168.7.255
    ether 00:0e:0c:01:d7:c8
    media: Ethernet autoselect (1000baseTX <full-duplex>)
    status: active
fxp0: flags=8843<UP,BROADCAST,RUNNING,SIMPLEX,MULTICAST> mtu 1500
    options=b<RXCSUM,TXCSUM,VLAN_MTU>
    inet 140.113.17.24 netmask 0xffffffff broadcast 140.113.17.255
    ether 00:02:b3:99:3e:71
    media: Ethernet autoselect (100baseTX <full-duplex>)
    status: active
```



Network Layer

Network Layer

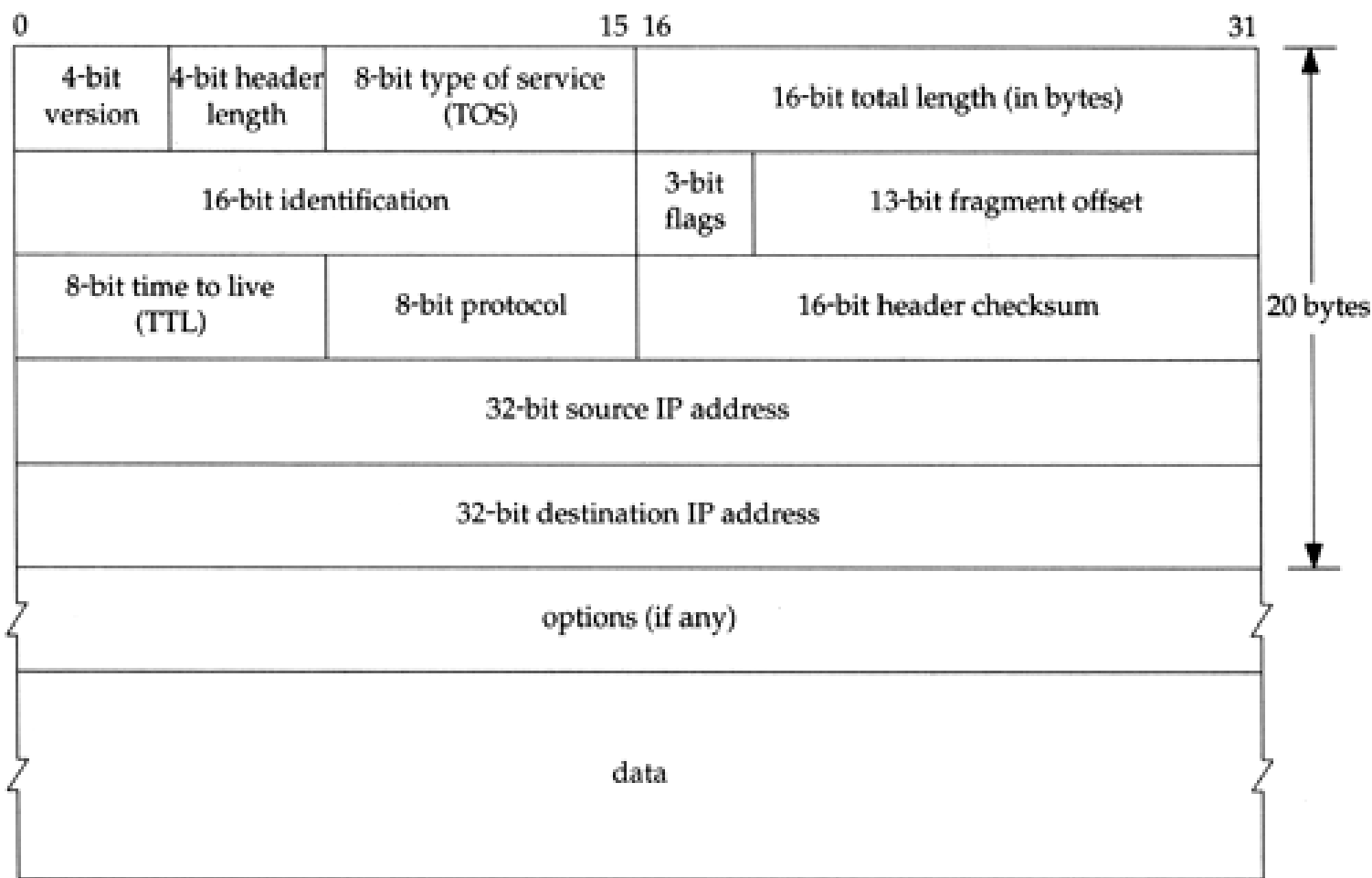
– Introduction to Network Layer

- ❑ Unreliable and connectionless datagram delivery service
 - IP Routing
 - IP provides best effort service (unreliable)
 - IP datagram can be delivered out of order (connectionless)
- ❑ Protocols using IP
 - TCP, UDP, ICMP, IGMP

Network Layer

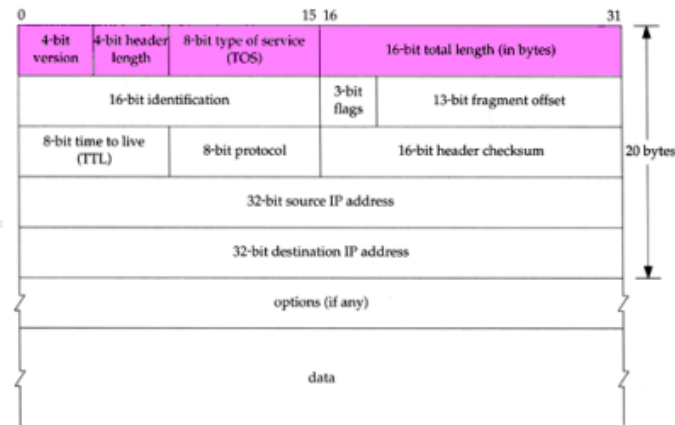
– IP Header (1)

- ❑ 20 bytes in total length, excepts options



Network Layer

– IP Header (2)



- ☐ Version (4-bit)
 - 4 for IPv4 and 6 for IPv6
- ☐ Header length (4-bit)
 - The number of 32-bit words in the header ($15 \times 4 = 60$ bytes)
 - Normally, the value is 5 (no option)
- ☐ TOS - Type of Service (8-bit)
 - IP Precedence: 3-bit precedence + 4-bit TOS + 1-bit unused
 - DSCP: 3-bit major class + 3-bit drop preference + 2-bit ECN
- ☐ Total length (16-bit)
 - Total length of the IP datagram in bytes

DSCP: Differentiated Services Code Point
ECN: Explicit Congestion Notification

Name	Binary Value	Application	Minimize delay	Maximize throughput	Maximize reliability	Minimize monetary cost	Hex value
Routine	000	Telnet/Rlogin FTP	1	0	0	0	0x10
Priority	001						
Immediate	010	control	1	0	0	0	0x10
Flash	011	data	0	1	0	0	0x08
Flash Override	100	any bulk data	0	1	0	0	0x08
Critic/Critical	101	TFTP	1	0	0	0	0x10
Internet Control	110	SMTP					
Network Control	111	command phase	1	0	0	0	0x10
		data phase	0	1	0	0	0x08

Network Layer

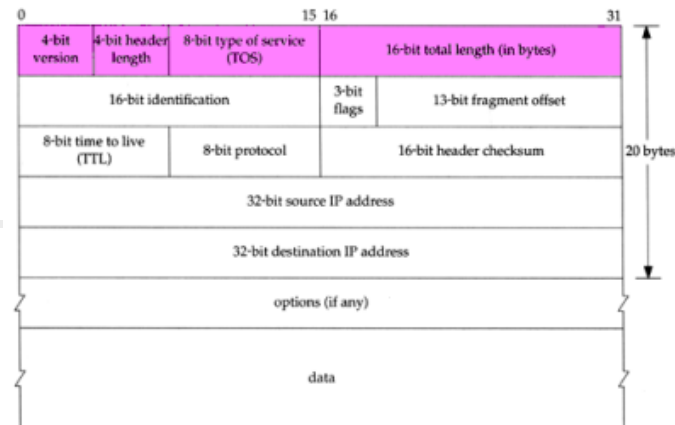
– IP Header (3)

❑ DSCP - Differentiated Services Code Point (6-bit)

- Supersede the ToS field in IPv4 to make per-hop behavior (PHB) decisions
 - Default
 - Best-effort traffic
 - Expedited Forwarding (EF)
 - Dedicated to low-loss, low-latency traffic
 - Class Selector
 - Backward compatibility with the IP Precedence field
 - Assured Forwarding (AF)
 - Give assurance of delivery under prescribed conditions

❑ ECN: Explicit Congestion Notification (2-bit)

- FreeBSD 8.0 implement ECN support for TCP
 - Enable ECN via sysctl(8)
 - net.inet.tcp.ecn.enable=1



DSCP Class Selector Names	Binary DSCP Values	IPP Binary Values	IPP Names
Default/CS0*	000000	000	Routine
CS1	001000	001	Priority
CS2	010000	010	Immediate
CS3	011000	011	Flash
CS4	100000	100	Flash Override
CS5	101000	101	Critic/Critical
CS6	110000	110	Internetwork Control
CS7	111000	111	Network Control

Queue Class	Low Drop Probability	Medium Drop Probability	High Drop Probability
	Name/Decimal/Binary	Name/Decimal/Binary	Name/Decimal/Binary
1	AF11 / 10 / 001010	AF12 / 12 / 001100	AF13 / 14 / 001110
2	AF21 / 18 / 010010	AF22 / 20 / 010100	AF23 / 22 / 010110
4	AF31 / 26 / 011010	AF32 / 28 / 011100	AF33 / 30 / 011110
5	AF41 / 34 / 100010	AF42 / 36 / 100100	AF43 / 38 / 100110

Binary Value	Description
00	Non ECN-Capable Transport, Non-ECT
10	ECN Capable Transport, ECT(0)
01	ECN Capable Transport, ECT(1)
11	Congestion Encountered, CE

Network Layer

– IP Header (3)

❑ Identification (16-bit)

- Identify the group of fragments of a single IP datagram

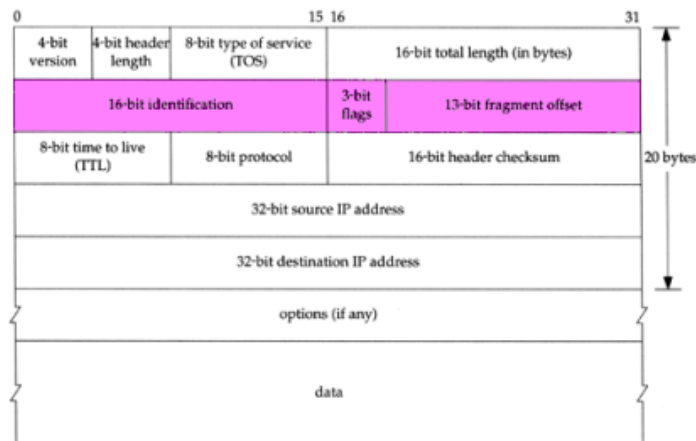
❑ Fragmentation offset (13-bit)

- Specify the offset of a particular fragment relative to the beginning of the original unfragmented IP datagram

❑ Flags (3-bit)

- All these three fields are used for fragmentation

Reserved	Don't Fragment (DF)	More Fragments (MF)
----------	---------------------	---------------------



Network Layer

– IP Header (4)

❑ TTL (8-bit)

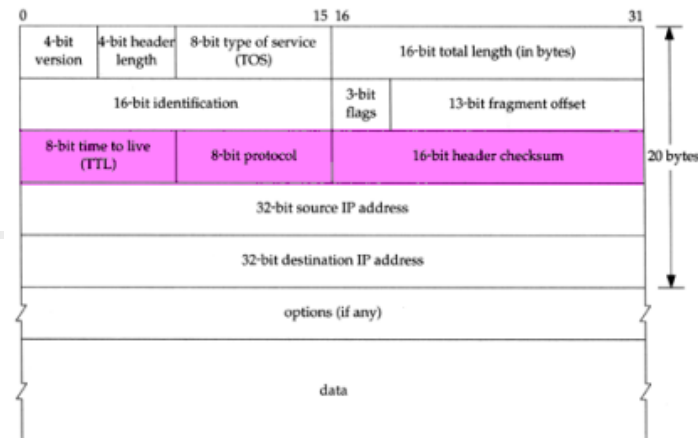
- Limit of next hop count of routers

❑ Protocol (8-bit)

- Used to demultiplex to other protocols
- TCP, UDP, ICMP, IGMP

❑ Header checksum (16-bit)

- Calculated over the IP header only
- If checksum error, IP discards the datagram and no error message is generated



Network Layer

– IP Address (1)

❑ 32-bit long

- Network part
 - Identify a logical network
- Host part
 - Identify a machine on certain network

❑ Ex:

- NCTU
 - Class B address: 140.113.0.0
 - Network ID: 140.113
 - Number of hosts: $256 \times 256 = 65536$

❑ IP address category

Class	1 st byte ^a	Format	Comments
A	1-126	N.H.H.H	Very early networks, or reserved for DOD
B	128-191	N.N.H.H	Large sites, usually subnetted, were hard to get
C	192-223	N.N.N.H	Easy to get, often obtained in sets
D	224-239	–	Multicast addresses, not permanently assigned
E	240-254	–	Experimental addresses

a. The values 0 and 255 are special and are not used as the first byte of regular IP addresses. 127 is reserved for the loopback address.

Network Layer

– Subnetting, CIDR, and Netmask (1)

❑ Problems of Class A or B network

- Number of hosts is enormous
- Hard to maintain and management
- Solution → Subnetting

❑ Problems of Class C network

- $255 \times 255 \times 255$ number of Class C network make the size of Internet routes huge
- Solution → Classless Inter-Domain Routing

Network Layer

– Subnetting, CIDR, and Netmask (2)

❑ Subnetting

- Borrow some bits from host ID to extend network ID
- Ex:
 - Class B address : 140.113.0.0
= 256 Class C-like IP addresses
in N.N.N.H subnetting method
 - 140.113.209.0 subnet
- Benefits of subnetting
 - Reduce the routing table size of Internet's routers
 - Ex:
 - All external routers have only one entry for 140.113 Class B network

Network Layer

– Subnetting, CIDR, and Netmask (3)

□ Netmask

- Specify how many bits of network-ID are used for network-ID
- Continuous 1 bits form the network part
- Ex:
 - 255.255.255.0 in NCTU-CS example
 - 256 hosts available
 - 255.255.255.248 in ADSL example
 - Only 8 hosts available
- Shorthand notation
 - Address/prefix-length
 - Ex: 140.113.209.8/24

Network Layer

– Subnetting, CIDR, and Netmask (4)

□ How to determine your network ID?

- Bitwise-AND IP and netmask
- Ex:
 - **140.113.214.37 & 255.255.255.0 → 140.113.214.0**
 - **140.113.209.37 & 255.255.255.0 → 140.113.209.0**

 - **140.113.214.37 & 255.255.0.0 → 140.113.0.0**
 - **140.113.209.37 & 255.255.0.0 → 140.113.0.0**

 - **211.23.188.78 & 255.255.255.248 → 211.23.188.72**
 - 78 = 01001110
 - 78 & 248 = 01001110 & 11111000 = 72

Network Layer

– Subnetting, CIDR, and Netmask (5)

❑ In a subnet, not all IP are available

- The first one IP → network ID
- The last one IP → broadcast address

- Ex:

Netmask 255.255.255.0
140.113.209.32/24

140.113.209.0 → network ID
140.113.209.255 → broadcast address
1 ~ 254, total 254 IPs are usable

Netmask 255.255.255.252
211.23.188.78/29

211.23.188.72 → network ID
211.23.188.79 → broadcast address
73 ~ 78, total 6 IPs are usable

Network Layer

– Subnetting, CIDR, and Netmask (6)

❑ The smallest subnetting

- Network portion : 30 bits
- Host portion : 2 bits
- ➔ 4 hosts, but only 2 IPs are available

❑ ipcalc

- /usr/ports/net-mgmt/ipcalc
- pkg install ipcalc

```
chbsd [/usr/ports/net-mgmt/ipcalc] -chwong- ipcalc 140.113.209.78/28
Address:    140.113.209.78      10001100.01110001.11010001.0100 1110
Netmask:    255.255.255.240 = 28 11111111.11111111.11111111.1111 0000
Wildcard:   0.0.0.15           00000000.00000000.00000000.0000 1111
=>
Network:    140.113.209.64/28   10001100.01110001.11010001.0100 0000
HostMin:    140.113.209.65      10001100.01110001.11010001.0100 0001
HostMax:    140.113.209.78      10001100.01110001.11010001.0100 1110
Broadcast:  140.113.209.79      10001100.01110001.11010001.0100 1111
Hosts/Net:  14                  Class B
```

Network Layer

– Subnetting, CIDR, and Netmask (7)

- Network configuration for various lengths of netmask

Length ^a	Host bits	Hosts/net ^b	Dec. netmask	Hex netmask
/20	12	4094	255.255.240.0	0xFFFFF000
/21	11	2046	255.255.248.0	0xFFFFF800
/22	10	1022	255.255.252.0	0xFFFFFC00
/23	9	510	255.255.254.0	0xFFFFFE00
/24	8	254	255.255.255.0	0xFFFFFFFF00
/25	7	126	255.255.255.128	0xFFFFFFFF80
/26	6	62	255.255.255.192	0xFFFFFFF00
/27	5	30	255.255.255.224	0xFFFFFFF80
/28	4	14	255.255.255.240	0xFFFFFFF0
/29	3	6	255.255.255.248	0xFFFFFFF8
/30	2	2	255.255.255.252	0xFFFFFFF4

Network Layer

– Subnetting, CIDR, and Netmask (8)

❑ CIDR (Classless Inter-Domain Routing)

- Use address mask instead of old address classes to determine the destination network
- CIDR requires modifications to routers and routing protocols
 - Need to transmit both destination address and mask
- Ex:
 - We can merge two Class C network:
203.19.68.0/24, 203.19.69.0/24 ➔ 203.19.68.0/23
- Benefit of CIDR
 - We can allocate continuous Class C network to organization
 - Reflect physical network topology
 - Reduce the size of routing table

Network Layer

– IP Routing (1)

❑ Difference between Host and Router

- Router forwards datagram from one of its interface to another, while host does not
- Almost every Unix system can be configured to act as a router or both
 - `net.inet.ip.forwarding=1`

❑ Router

- IP layer has a routing table, which is used to store the information for forwarding datagram
- When router receiving a datagram
 - If Dst. IP = my IP, demultiplex to other protocol
 - Other, forward the IP based on routing table

Network Layer

– IP Routing (2)

❑ Routing table information

- Destination IP
- IP address of next-hop router or IP address of a directly connected network
- Flags
- Next interface

❑ IP routing

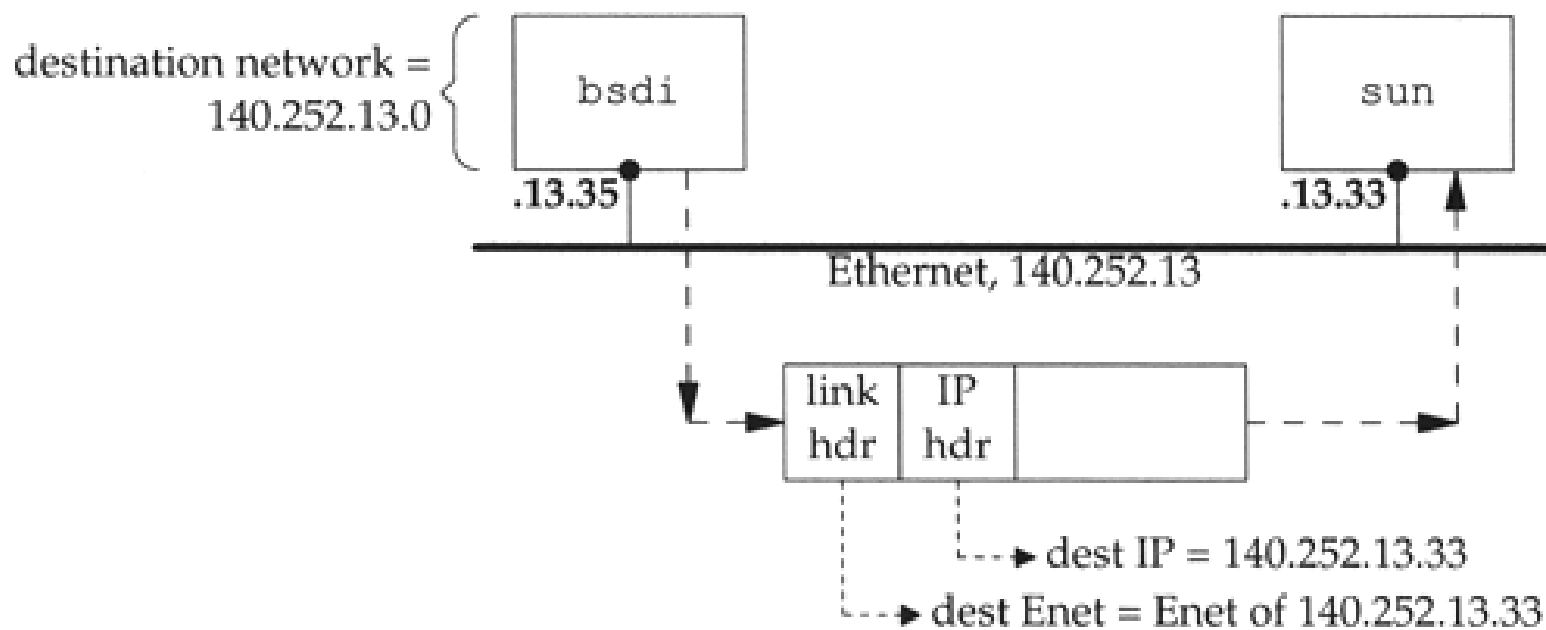
- Done on a hop-by-hop basis
- It assumes that the next-hop router is closer to the destination
- Steps:
 - Search routing table for complete matched IP address
 - Send to next-hop router or to the directly connected NIC
 - Search routing table for matched network ID
 - Send to next-hop router or to the directly connected NIC
 - Search routing table for default route
 - Send to this default next-hop router
 - host or network unreachable

Network Layer

– IP Routing (3)

❑ Ex1: routing in the same network

- bsd1: 140.252.13.35
- sun: 140.252.13.33



Ex Routing table:

140.252.13.33

00:d0:59:83:d9:16

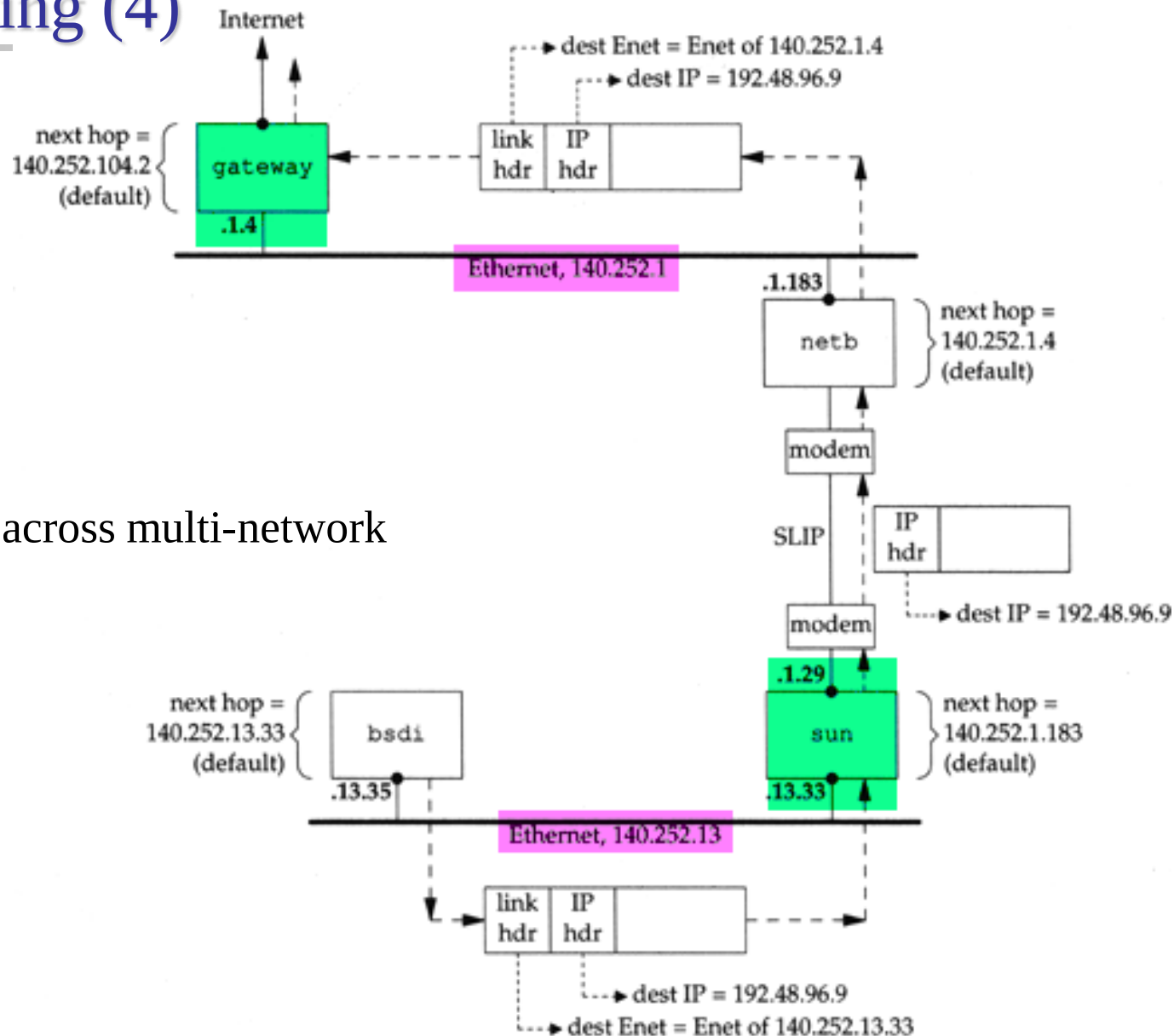
UHLW fxp1

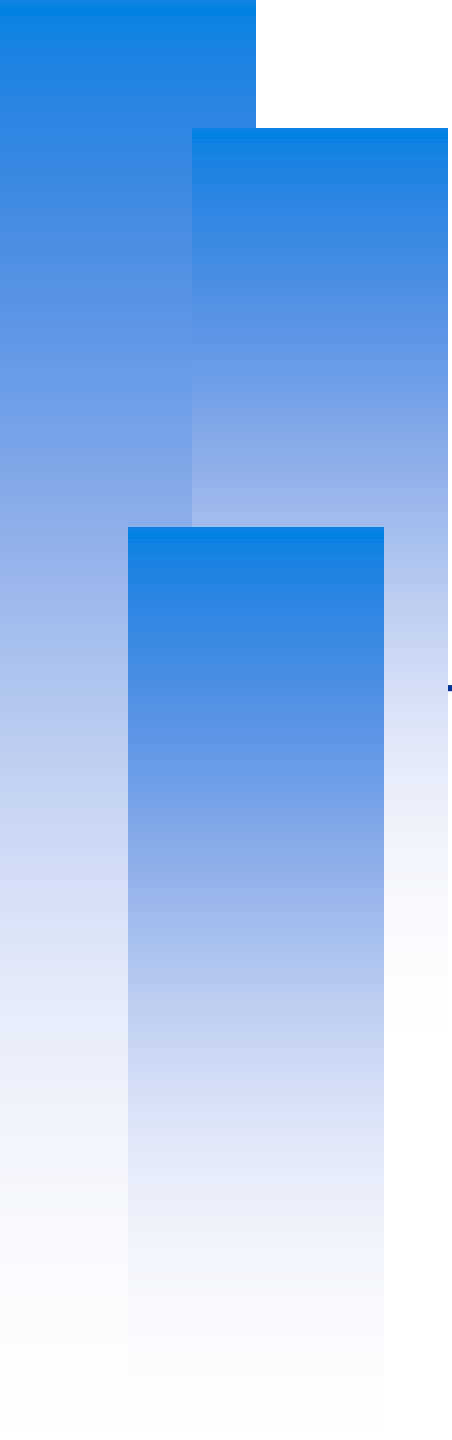
Network Layer

– IP Routing (4)

□ Ex2:

- routing across multi-network



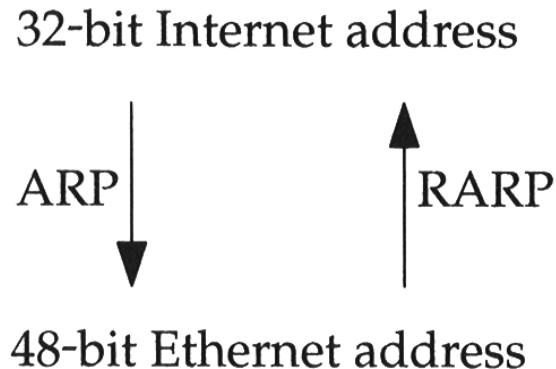


ARP and RARP

Something between
MAC (link layer)
And
IP (network layer)

ARP and RARP

- ❑ ARP – Address Resolution Protocol and RARP – Reverse ARP
 - Mapping between IP and Ethernet address



- ❑ When an Ethernet frame is sent on LAN from one host to another,
 - It is the 48-bit Ethernet address that determines for which interface the frame is destined

ARP and RARP

– ARP Example

□ Example

% ftp bsd1

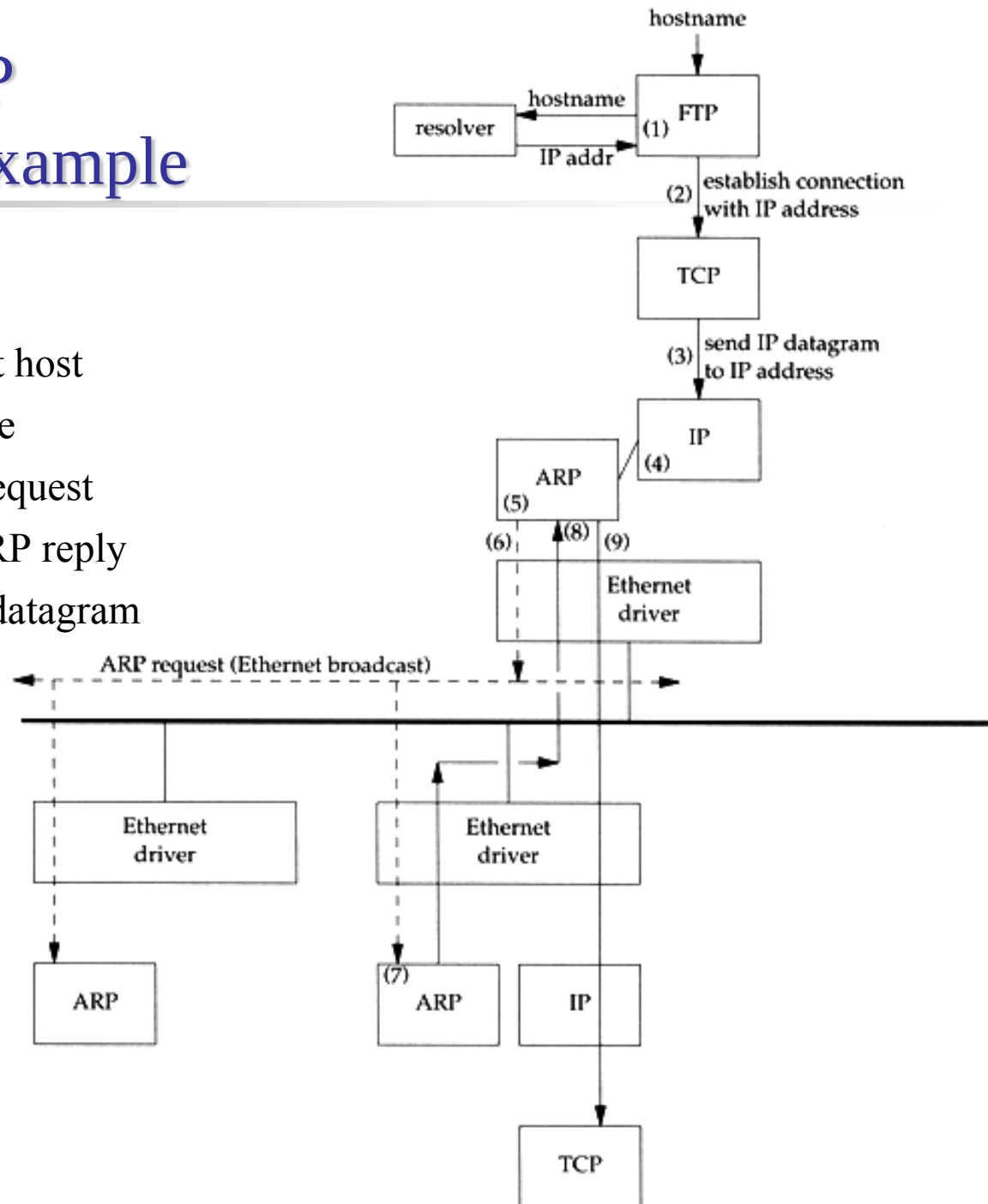
(4) next-hop or direct host

(5) Search ARP cache

(6) Broadcast ARP request

(7) bsd1 response ARP reply

(9) Send original IP datagram



ARP and RARP

– ARP Cache

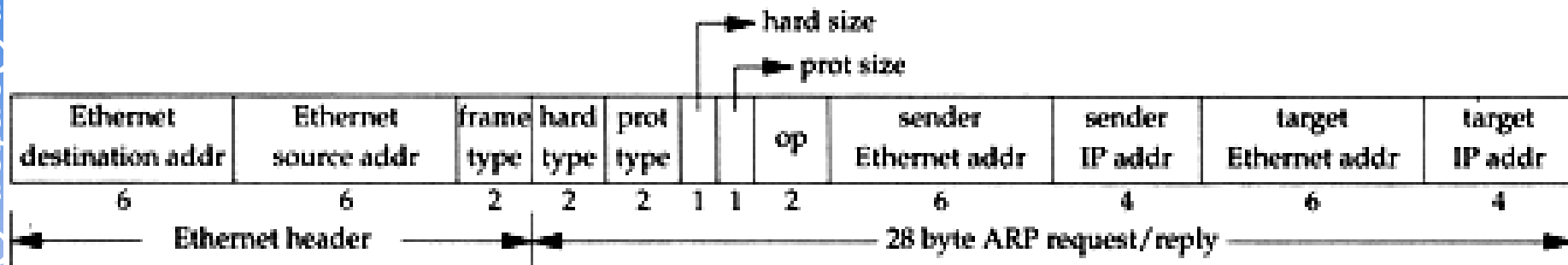
❑ Maintain recent ARP results

- come from both ARP request and reply
- expiration time
 - Complete entry = 20 minutes
 - Incomplete entry = 3 minutes
- Use arp command to see the cache
- Ex:
 - % arp -a
 - % arp -da
 - % arp -S 140.113.235.132 00:0e:a6:94:24:6e

```
csduty /home/chwong] -chwong- arp -a
cshome (140.113.235.101) at 00:0b:cd:9e:74:61 on em0 [ethernet]
bsd1 (140.113.235.131) at 00:11:09:a0:04:74 on em0 [ethernet]
? (140.113.235.160) at (incomplete) on em0 [ethernet]
```

ARP and RARP

– ARP/RARP Packet Format



- ☐ Ethernet destination addr: all 1's (broadcast)
- ☐ Known value for IP <-> Ethernet
 - Frame type: 0x0806 for ARP, 0x8035 for RARP
 - Hardware type: type of hardware address (1 for Ethernet)
 - Protocol type: type of upper layer address (0x0800 for IP)
 - Hard size: size in bytes of hardware address (6 for Ethernet)
 - Protocol size: size in bytes of upper layer address (4 for IP)
 - Op: 1, 2, 3, 4 for ARP request, reply, RARP request, reply

ARP and RARP

– Use tcpdump to see ARP

❑ Host 140.113.17.212 → 140.113.17.215

- Clear ARP cache of 140.113.17.212
 - % sudo arp -d 140.113.17.215
- Run tcpdump on 140.113.17.215 **(00:11:d8:06:1e:81)**
 - % sudo tcpdump -i sk0 -e arp
 - % sudo tcpdump -i sk0 -n -e arp
 - % sudo tcpdump -i sk0 -n -t -e arp
- On 140.113.17.212, ssh to 140.113.17.215

```
15:18:54.899779 00:90:96:23:8f:7d > Broadcast, ethertype ARP (0x0806), length 60:
arp who-has nabsd tell chbsd.csie.nctu.edu.tw
15:18:54.899792 00:11:d8:06:1e:81 > 00:90:96:23:8f:7d, ethertype ARP (0x0806), length 42:
arp reply nabsd is-at 00:11:d8:06:1e:81
```

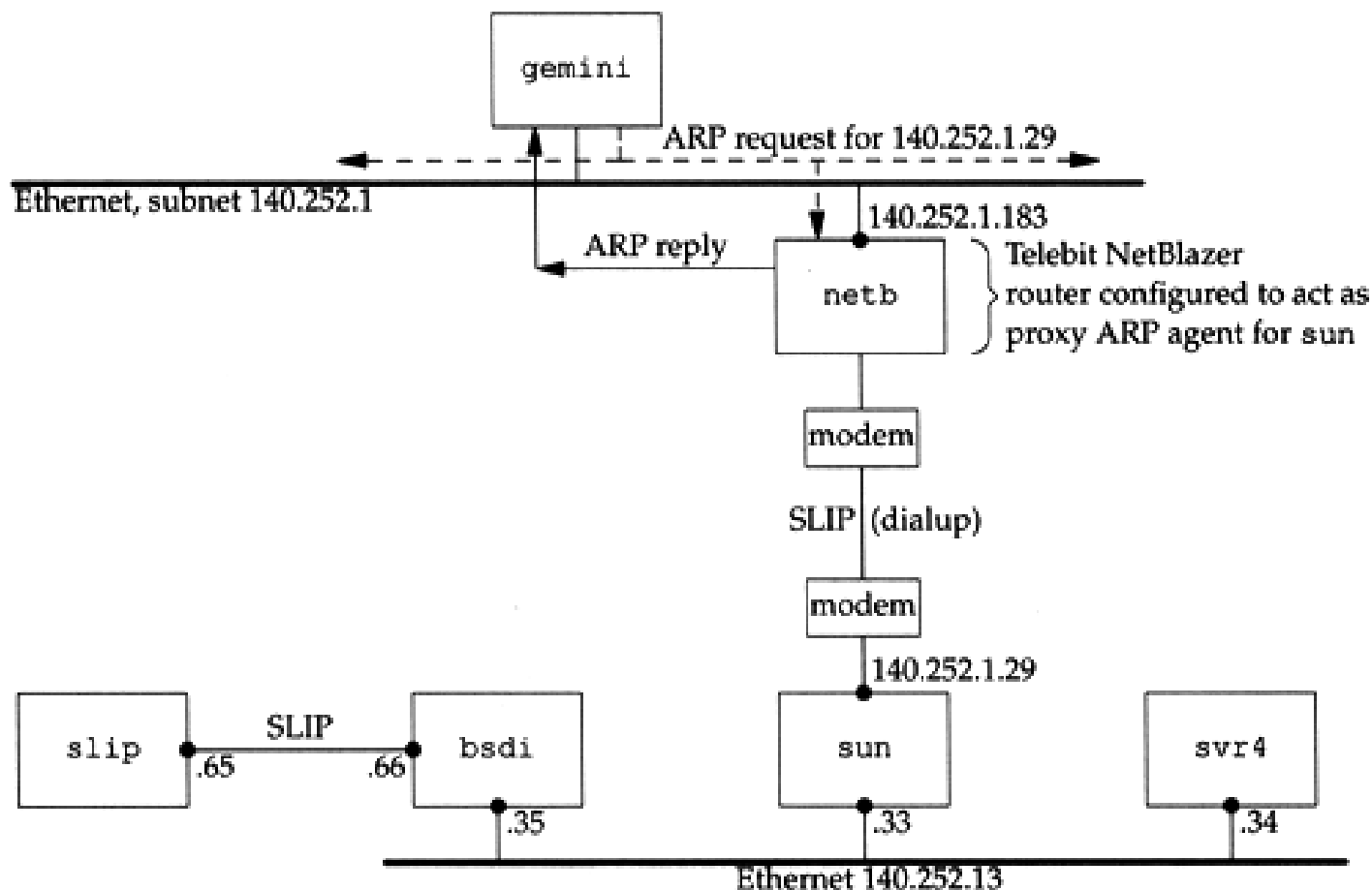
```
15:26:13.847417 00:90:96:23:8f:7d > ff:ff:ff:ff:ff:ff, ethertype ARP (0x0806), length 60:
arp who-has 140.113.17.215 tell 140.113.17.212
15:26:13.847434 00:11:d8:06:1e:81 > 00:90:96:23:8f:7d, ethertype ARP (0x0806), length 42:
arp reply 140.113.17.215 is-at 00:11:d8:06:1e:81
```

```
00:90:96:23:8f:7d > ff:ff:ff:ff:ff:ff, ethertype ARP (0x0806), length 60:
arp who-has 140.113.17.215 tell 140.113.17.212
00:11:d8:06:1e:81 > 00:90:96:23:8f:7d, ethertype ARP (0x0806), length 42:
arp reply 140.113.17.215 is-at 00:11:d8:06:1e:81
```

ARP and RARP

– Proxy ARP

- ❑ Let router answer ARP request on one of its networks for a host on another of its network



ARP and RARP

– Gratuitous ARP

❑ Gratuitous ARP

- The host sends an ARP request looking for its own IP
- Provide two features
 - Used to determine whether there is another host configured with the same IP
 - Used to cause any other host to update ARP cache when changing hardware address

ARP and RARP

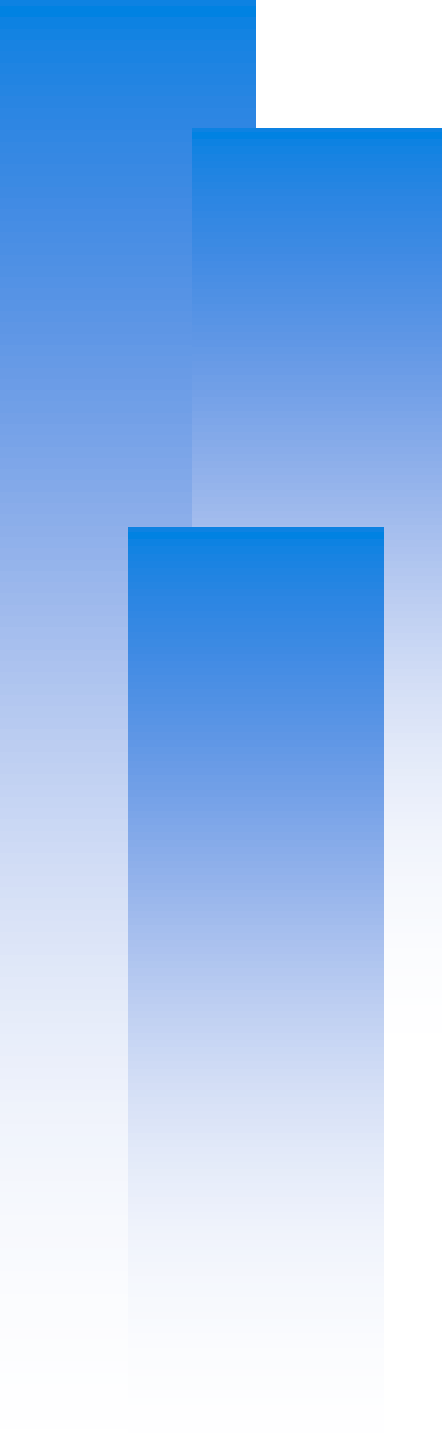
– RARP

❑ Principle

- Used for the diskless system to read its hardware address from the NIC and send an RARP request to gain its IP

❑ RARP Server Design

- RARP server must maintain the map from hardware address to an IP address for many host
- Link-layer broadcast
 - This prevent most routers from forwarding an RARP request



ICMP – Internet Control Message Protocol

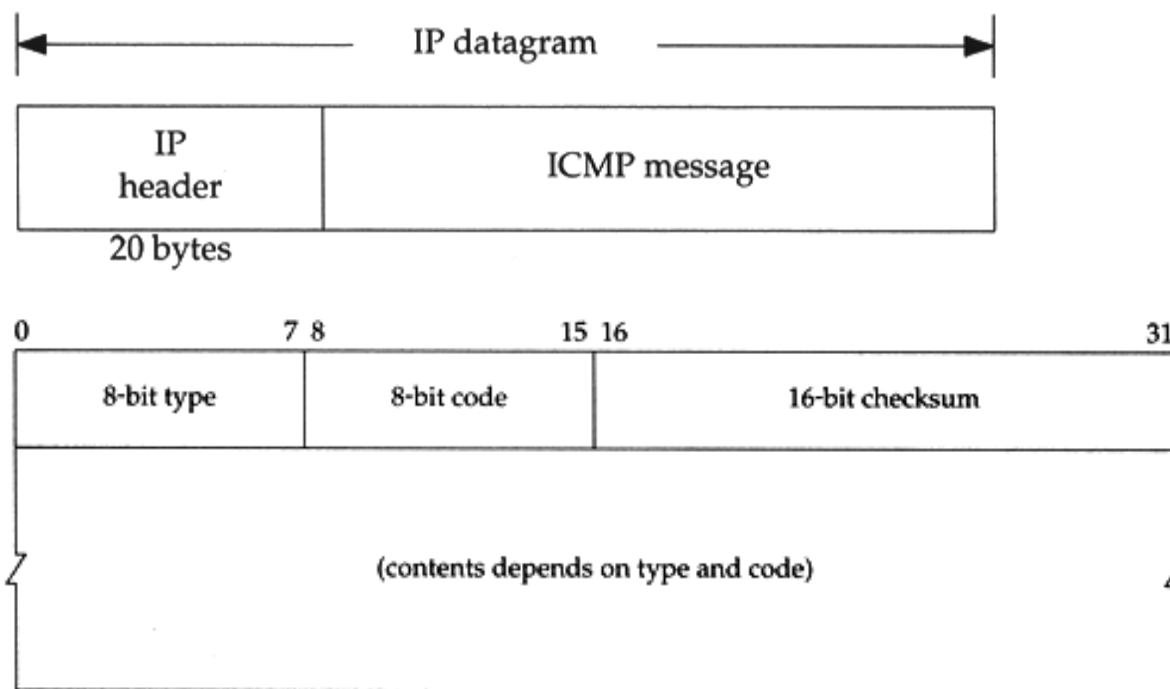
ICMP

– Introduction

❑ Part of the IP layer

- ICMP messages are transmitted within IP datagram
- ICMP communicates error messages and other conditions that require attention for other protocols

❑ ICMP message format



ICMP

– Message Type (1)

<i>type</i>	<i>code</i>	Description	Query	Error
0	0	echo reply (Ping reply, Chapter 7)	•	
3		destination unreachable:		
	0	network unreachable (Section 9.3)		•
	1	host unreachable (Section 9.3)		•
	2	protocol unreachable		•
	3	port unreachable (Section 6.5)		•
	4	fragmentation needed but don't-fragment bit set (Section 11.6)		•
	5	source route failed (Section 8.5)		•
	6	destination network unknown		•
	7	destination host unknown		•
	8	source host isolated (obsolete)		•
	9	destination network administratively prohibited		•
	10	destination host administratively prohibited		•
	11	network unreachable for TOS (Section 9.3)		•
	12	host unreachable for TOS (Section 9.3)		•
	13	communication administratively prohibited by filtering		•
	14	host precedence violation		•
	15	precedence cutoff in effect		•
4	0	source quench (elementary flow control, Section 11.11)		•

ICMP

– Message Type (2)

5		redirect (Section 9.5):		
	0	redirect for network		•
	1	redirect for host		•
	2	redirect for type-of-service and network		•
	3	redirect for type-of-service and host		•
8	0	echo request (Ping request, Chapter 7)	•	
9	0	router advertisement (Section 9.6)	•	
10	0	router solicitation (Section 9.6)	•	
11		time exceeded:		
	0	time-to-live equals 0 during transit (Traceroute, Chapter 8)		•
	1	time-to-live equals 0 during reassembly (Section 11.5)		•
12		parameter problem:		
	0	IP header bad (catchall error)		•
	1	required option missing		•
13	0	timestamp request (Section 6.4)	•	
14	0	timestamp reply (Section 6.4)	•	
15	0	information request (obsolete)	•	
16	0	information reply (obsolete)	•	
17	0	address mask request (Section 6.3)	•	
18	0	address mask reply (Section 6.3)	•	

ICMP – Query Message

– Address Mask Request/Reply (1)

❑ Address Mask Request and Reply

- Used for diskless system to obtain its subnet mask
- Identifier and sequence number
 - Can be set to anything for sender to match reply with request
- The receiver will response an ICMP reply with the subnet mask of the receiving NIC

0	8	16	31
TYPE (17 or 18)	CODE (0)	CHECKSUM	
IDENTIFIER		SEQUENCE NUMBER	
ADDRESS MASK			

ICMP – Query Message

– Address Mask Request/Reply (2)

❑ Ex:

```
chbsd [/home/chwong] -chwong- ping -M m sun1.cs.nctu.edu.tw
ICMP_MASKREQ
PING sun1.cs.nctu.edu.tw (140.113.235.171): 56 data bytes
68 bytes from 140.113.235.171: icmp_seq=0 ttl=251 time=0.663 ms mask=255.255.255.0
68 bytes from 140.113.235.171: icmp_seq=1 ttl=251 time=1.018 ms mask=255.255.255.0
68 bytes from 140.113.235.171: icmp_seq=2 ttl=251 time=1.028 ms mask=255.255.255.0
68 bytes from 140.113.235.171: icmp_seq=3 ttl=251 time=1.026 ms mask=255.255.255.0
^C
--- sun1.cs.nctu.edu.tw ping statistics ---
4 packets transmitted, 4 packets received, 0% packet loss
round-trip min/avg/max/stddev = 0.663/0.934/1.028/0.156 ms

chbsd [/home/chwong] -chwong- icmpquery -m sun1
sun1                               : 0xFFFFFFFF00
```

※ icmpquery can be found in /usr/ports/net-mgmt/icmpquery

ICMP – Query Message

– Timestamp Request/Reply (1)

❑ Timestamp request and reply

- Allow a system to query another for the current time
- Milliseconds resolution, since midnight UTC
- Requestor
 - Fill in the originate timestamp and send
- Reply system
 - Fill in the receive timestamp when it receives the request and the transmit time when it sends the reply

0	8	16	31
TYPE (13 or 14)	CODE (0)	CHECKSUM	
IDENTIFIER		SEQUENCE NUMBER	
ORIGINATE TIMESTAMP			
RECEIVE TIMESTAMP			
TRANSMIT TIMESTAMP			

ICMP – Query Message

– Timestamp Request/Reply (2)

❑ Ex:

```
chbsd [/home/chwong] -chwong- ping -M time nabsd
ICMP_TSTAMP
PING nabsd.cs.nctu.edu.tw (140.113.17.215): 56 data bytes
76 bytes from 140.113.17.215: icmp_seq=0 ttl=64 time=0.663 ms
    tso=06:47:46 tsr=06:48:24 tst=06:48:24
76 bytes from 140.113.17.215: icmp_seq=1 ttl=64 time=1.016 ms
    tso=06:47:47 tsr=06:48:25 tst=06:48:25

chbsd [/home/chwong] -chwong- icmpquery -t nabsd
nabsd                : 14:54:47
```

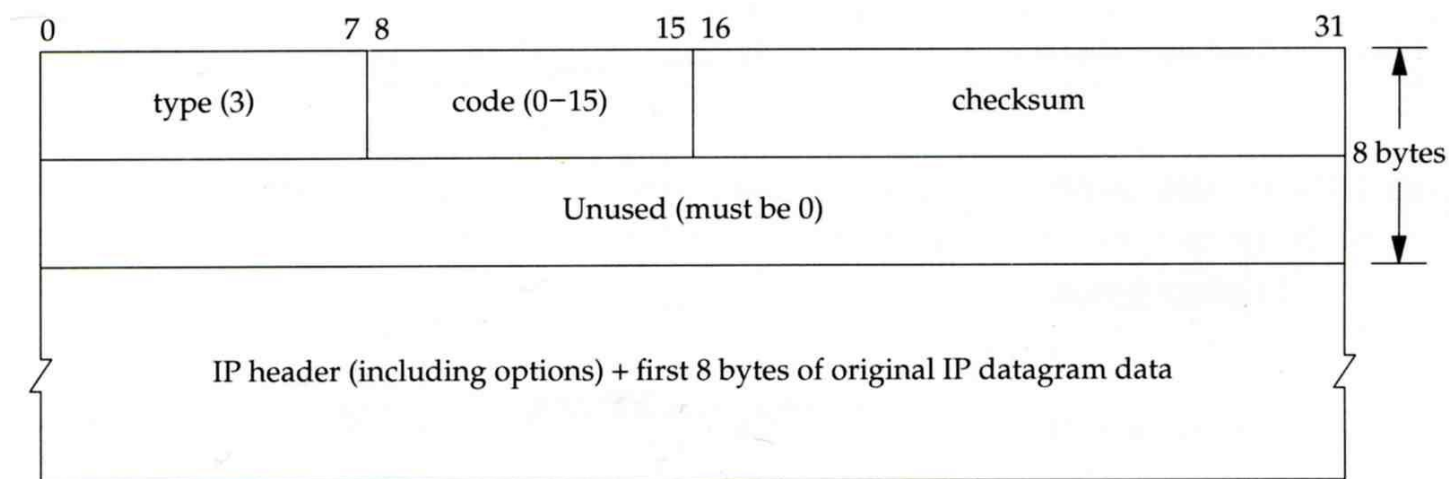
```
nabsd [/home/chwong] -chwong- sudo tcpdump -i sk0 -e icmp
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on sk0, link-type EN10MB (Ethernet), capture size 96 bytes
14:48:24.999106 00:90:96:23:8f:7d > 00:11:d8:06:1e:81, ethertype IPv4 (0x0800), length 110:
    chbsd.csie.nctu.edu.tw > nabsd: ICMP time stamp query id 18514 seq 0, length 76
14:48:24.999148 00:11:d8:06:1e:81 > 00:90:96:23:8f:7d, ethertype IPv4 (0x0800), length 110:
    nabsd > chbsd.csie.nctu.edu.tw: ICMP time stamp reply id 18514 seq 0: org 06:47:46.326,
    recv 06:48:24.998, xmit 06:48:24.998, length 76
14:48:26.000598 00:90:96:23:8f:7d > 00:11:d8:06:1e:81, ethertype IPv4 (0x0800), length 110:
    chbsd.csie.nctu.edu.tw > nabsd: ICMP time stamp query id 18514 seq 1, length 76
14:48:26.000618 00:11:d8:06:1e:81 > 00:90:96:23:8f:7d, ethertype IPv4 (0x0800), length 110:
    nabsd > chbsd.csie.nctu.edu.tw: ICMP time stamp reply id 18514 seq 1: org 06:47:47.327,
    recv 06:48:25.999, xmit 06:48:25.999, length 76
```

ICMP – Error Message

– Destination Unreachable Error Message

❑ Format

- 8 bytes ICMP Header
 - IP header
 - Let ICMP know how to interpret the 8 bytes that follow
 - first 8 bytes that followed this IP header
 - Information about who generates the error

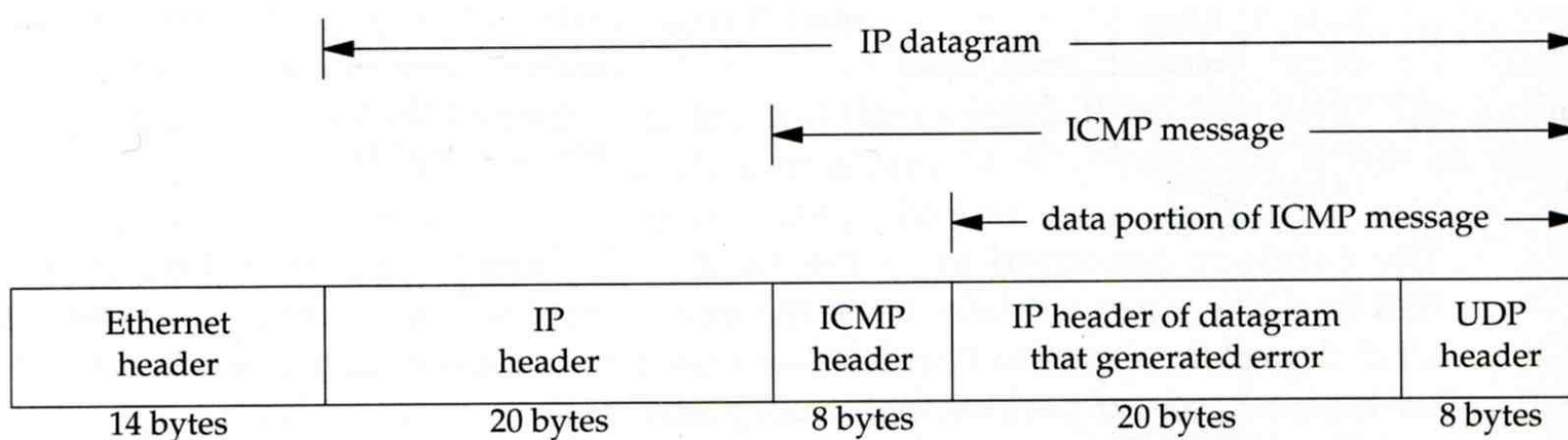


ICMP – Error Message

– Port Unreachable (1)

❑ ICMP port unreachable

- Type = 3 , code = 3
- Host receives a UDP datagram but the destination port does not correspond to a port that some process has in use



ICMP – Error Message

– Port Unreachable (2)

❑ Ex:

- Using TFTP (Trivial File Transfer Protocol)
 - Original port: 69

```
chbsd [/home/chwong] -chwong- tftp
tftp> connect localhost 8888
tftp> get temp.foo
Transfer timed out.

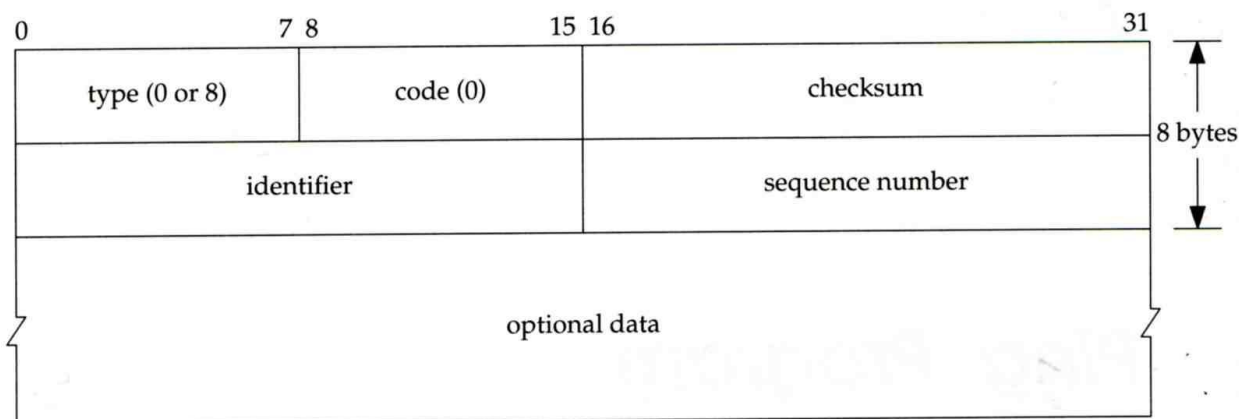
tftp>
```

```
chbsd [/home/chwong] -chwong- sudo tcpdump -i lo0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on lo0, link-type NULL (BSD loopback), capture size 96 bytes
15:01:24.788511 IP localhost.62089 > localhost.8888: UDP, length 16
15:01:24.788554 IP localhost > localhost:
    ICMP localhost udp port 8888 unreachable, length 36
15:01:29.788626 IP localhost.62089 > localhost.8888: UDP, length 16
15:01:29.788691 IP localhost > localhost:
    ICMP localhost udp port 8888 unreachable, length 36
```

ICMP

– Ping Program (1)

- ❑ Use ICMP to test whether another host is reachable
 - Type 8, ICMP echo request
 - Type 0, ICMP echo reply
- ❑ ICMP echo request/reply format
 - Identifier: process ID of the sending process
 - Sequence number: start with 0
 - Optional data: any optional data sent must be echoed



ICMP

– Ping Program (2)

❑ Ex:

- chbsd ping nabsd
- execute “tcpdump -i sk0 -X -e icmp” on nabsd

```
chbsd [/home/chwong] -chwong- ping nabsd
PING nabsd.cs.nctu.edu.tw (140.113.17.215): 56 data bytes
64 bytes from 140.113.17.215: icmp_seq=0 ttl=64 time=0.520 ms
```

```
15:08:12.631925 00:90:96:23:8f:7d > 00:11:d8:06:1e:81, ethertype IPv4 (0x0800), length 98:
chbsd.csie.nctu.edu.tw > nabsd: ICMP echo request, id 56914, seq 0, length 64
0x0000: 4500 0054 f688 0000 4001 4793 8c71 11d4 E..T....@.G..q..
0x0010: 8c71 11d7 0800 a715 de52 0000 45f7 9f35 .q.....R..E..5
0x0020: 000d a25a 0809 0a0b 0c0d 0e0f 1011 1213 ...Z.....
0x0030: 1415 1617 1819 1a1b 1c1d 1e1f 2021 2223 .....!"#
0x0040: 2425 2627 2829 2a2b 2c2d 2e2f 3031 3233 $%&'()*+,-./0123
0x0050: 3435                                     45

15:08:12.631968 00:11:d8:06:1e:81 > 00:90:96:23:8f:7d, ethertype IPv4 (0x0800), length 98:
nabsd > chbsd.csie.nctu.edu.tw: ICMP echo reply, id 56914, seq 0, length 64
0x0000: 4500 0054 d97d 0000 4001 649e 8c71 11d7 E..T.}..@.d..q..
0x0010: 8c71 11d4 0000 af15 de52 0000 45f7 9f35 .q.....R..E..5
0x0020: 000d a25a 0809 0a0b 0c0d 0e0f 1011 1213 ...Z.....
0x0030: 1415 1617 1819 1a1b 1c1d 1e1f 2021 2223 .....!"#
0x0040: 2425 2627 2829 2a2b 2c2d 2e2f 3031 3233 $%&'()*+,-./0123
0x0050: 3435                                     45
```

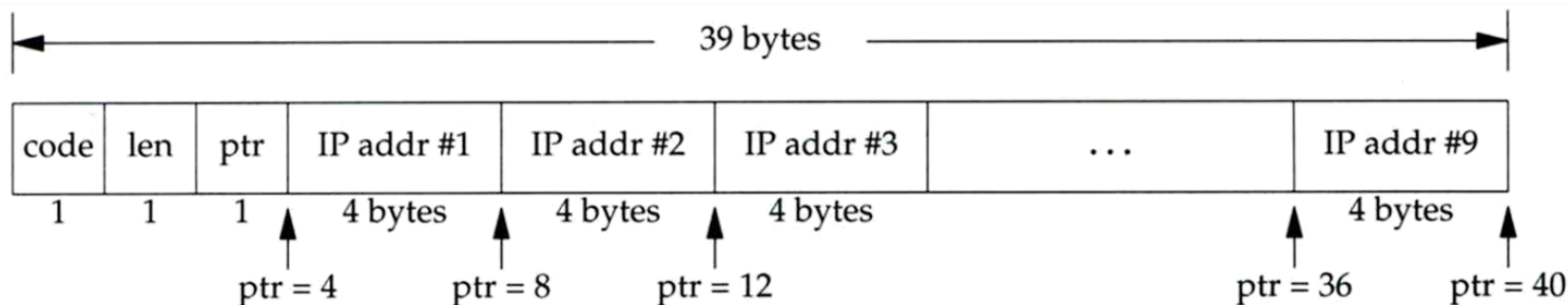
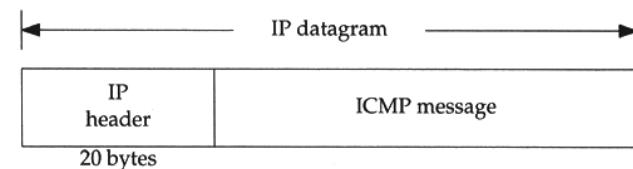
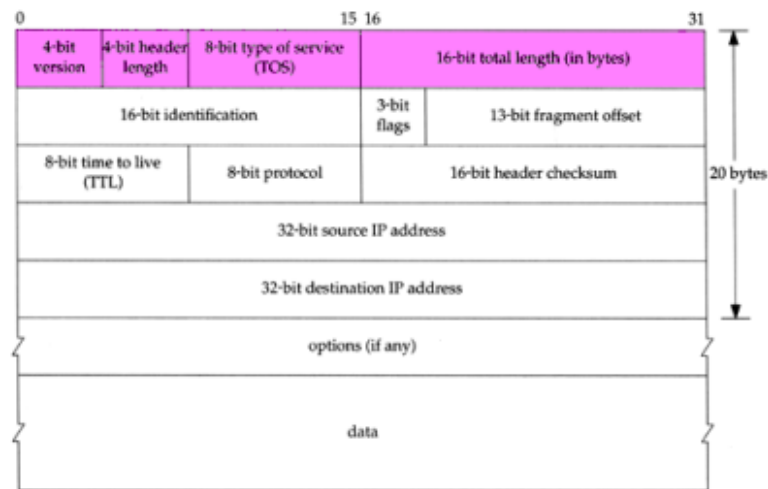
Type,
Code

id

ICMP

– Ping Program (3)

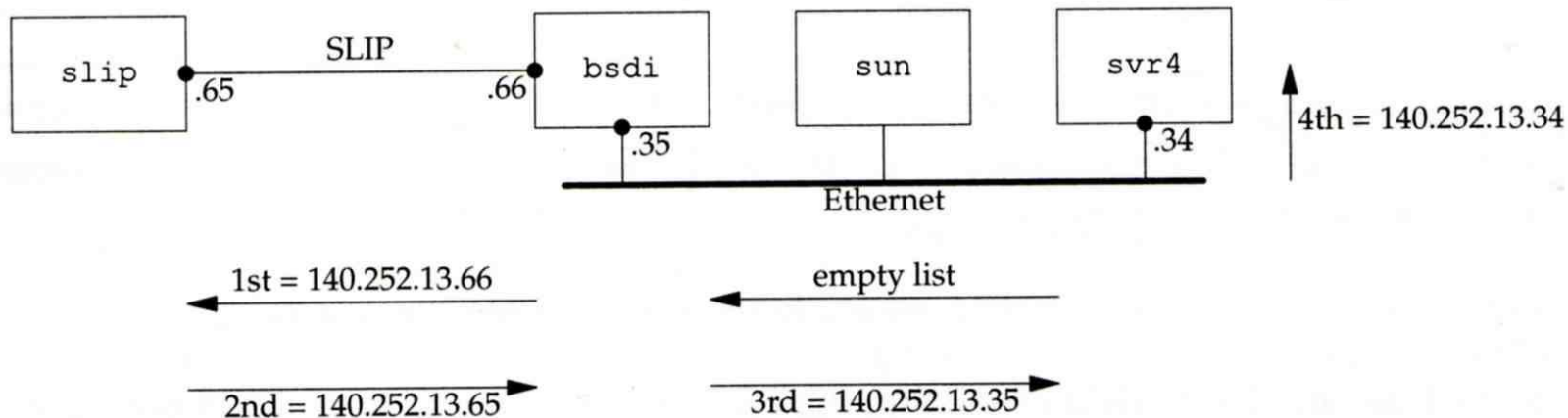
- ❑ To get the route that packets take to host
- Taking use of “IP Record Route Option”
 - Command: ping -R
 - Cause every router that handles the datagram to add its (**outgoing**) IP address to a list in the options field.
 - Format of Option field for IP RR Option
 - code: type of IP Option (7 for RR)
 - len: total number of bytes of the RR option
 - ptr: 4 ~ 40 used to point to the next IP address
 - Only **9** IP addresses can be stored
 - Limitation of IP header



ICMP

– Ping Program (4)

□ Example:



```
svr4 % ping -R slip
PING slip (140.252.13.65): 56 data bytes
64 bytes from 140.252.13.65: icmp_seq=0 ttl=254 time=280 ms
RR:      bsd (140.252.13.66)
         slip (140.252.13.65)
         bsd (140.252.13.35)
         svr4 (140.252.13.34)
64 bytes from 140.252.13.65: icmp_seq=1 ttl=254 time=280 ms (same route)
64 bytes from 140.252.13.65: icmp_seq=2 ttl=254 time=270 ms (same route)
^?
--- slip ping statistics ---
3 packets transmitted, 3 packets received, 0% packet loss
round-trip min/avg/max = 270/276/280 ms
```

```
chbsd [/home/chwong] -chwong- ping -R www.nctu.edu.tw
PING www.nctu.edu.tw (140.113.250.5): 56 data bytes
64 bytes from 140.113.250.5: icmp_seq=0 ttl=61 time=2.361 ms
RR:   ProjE27-253.NCTU.edu.tw (140.113.27.253)
      140.113.0.57
      CC250-gw.NCTU.edu.tw (140.113.250.253)
      www.NCTU.edu.tw (140.113.250.5)
      www.NCTU.edu.tw (140.113.250.5)
      140.113.0.58
      ProjE27-254.NCTU.edu.tw (140.113.27.254)
      e3rtn.csie.nctu.edu.tw (140.113.17.254)
      chbsd.csie.nctu.edu.tw (140.113.17.212)
64 bytes from 140.113.250.5: icmp_seq=1 ttl=61 time=3.018 ms   (same route)
```

65

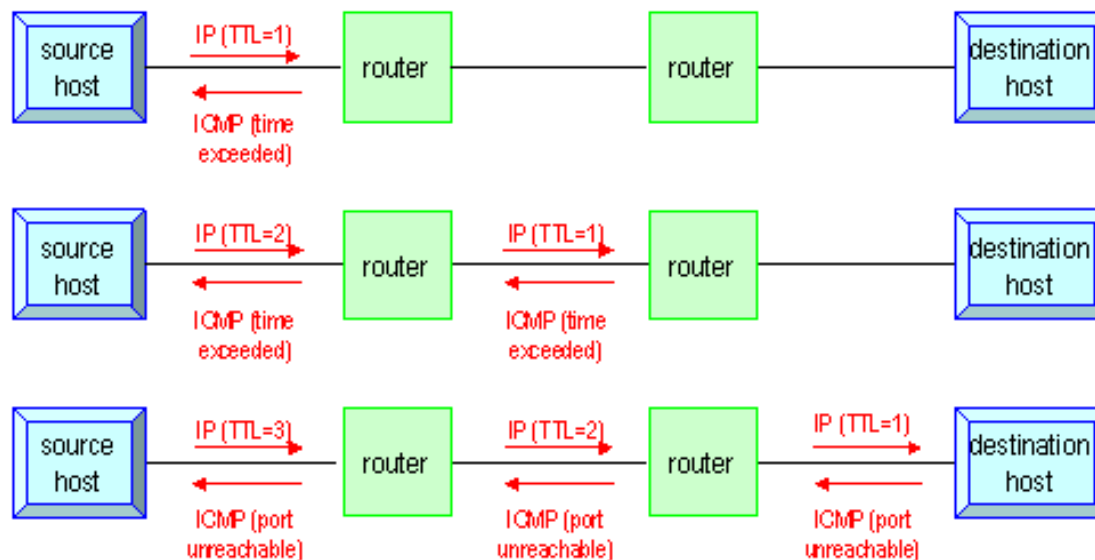
Traceroute Program (1)

- ❑ To print the route packets take to network host
- ❑ Drawbacks of IP RR options (ping -R)
 - Not all routers have supported the IP RR option
 - Limitation of IP header length
- ❑ Background knowledge of traceroute
 - When a router receive a datagram, , it will decrement the TTL by one
 - When a router receive a datagram with TTL = 0 or 1,
 - it will through away the datagram and
 - sends back a "Time exceeded" ICMP message
 - Unused UDP port will generate a "port unreachable" ICMP message

Traceroute Program (2)

❑ Operation of traceroute

- Send UDP with port > 30000, encapsulated with IP header with TTL = 1, 2, 3, ... continuously
- When router receives the datagram and TTL = 1, it returns a “Time exceed” ICMP message
- When destination host receives the datagram and TTL = 1, it returns a “Port unreachable” ICMP message



Traceroute Program (3)

❑ Time exceed ICMP message

- Type = 11, code = 0 or 1
 - Code = 0 means TTL=0 during transit
 - Code = 1 means TTL=0 during reassembly
- First 8 bytes of datagram
 - UDP header

0	8	16	31
TYPE (11)	CODE (0 or 1)	CHECKSUM	
UNUSED (MUST BE ZERO)			
INTERNET HEADER + FIRST 64 BITS OF DATAGRAM			
...			

Traceroute Program (4)

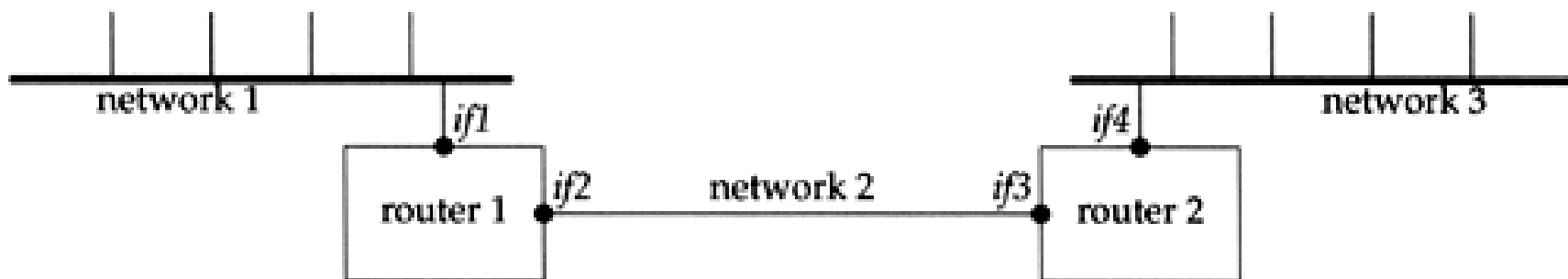
□ Ex:

```
nabsd [/home/chwong] -chwong- traceroute bsd1.cs.nctu.edu.tw
traceroute to bsd1.cs.nctu.edu.tw (140.113.235.131), 64 hops max, 40 byte packets
 1 e3rtn.csie.nctu.edu.tw (140.113.17.254) 0.377 ms 0.365 ms 0.293 ms
 2 ProjE27-254.NCTU.edu.tw (140.113.27.254) 0.390 ms 0.284 ms 0.391 ms
 3 140.113.0.58 (140.113.0.58) 0.292 ms 0.282 ms 0.293 ms
 4 140.113.0.165 (140.113.0.165) 0.492 ms 0.385 ms 0.294 ms
 5 bsd1.cs.nctu.edu.tw (140.113.235.131) 0.393 ms 0.281 ms 0.393 ms
```

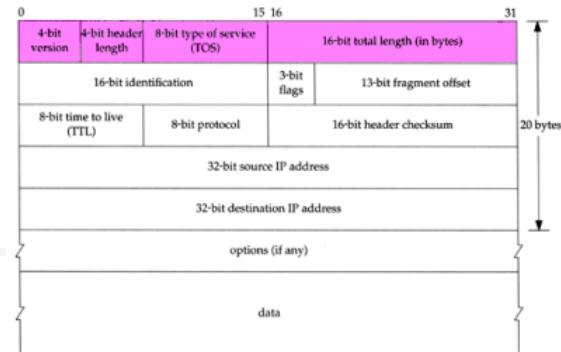
```
nabsd [/home/chwong] -chwong- sudo tcpdump -i sk0 -t icmp
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on sk0, link-type EN10MB (Ethernet), capture size 96 bytes
IP e3rtn.csie.nctu.edu.tw > nabsd: ICMP time exceeded in-transit, length 36
IP e3rtn.csie.nctu.edu.tw > nabsd: ICMP time exceeded in-transit, length 36
IP e3rtn.csie.nctu.edu.tw > nabsd: ICMP time exceeded in-transit, length 36
IP ProjE27-254.NCTU.edu.tw > nabsd: ICMP time exceeded in-transit, length 36
IP ProjE27-254.NCTU.edu.tw > nabsd: ICMP time exceeded in-transit, length 36
IP ProjE27-254.NCTU.edu.tw > nabsd: ICMP time exceeded in-transit, length 36
IP 140.113.0.58 > nabsd: ICMP time exceeded in-transit, length 36
IP 140.113.0.58 > nabsd: ICMP time exceeded in-transit, length 36
IP 140.113.0.58 > nabsd: ICMP time exceeded in-transit, length 36
IP 140.113.0.165 > nabsd: ICMP time exceeded in-transit, length 36
IP 140.113.0.165 > nabsd: ICMP time exceeded in-transit, length 36
IP 140.113.0.165 > nabsd: ICMP time exceeded in-transit, length 36
IP bsd1.cs.nctu.edu.tw > nabsd: ICMP bsd1.cs.nctu.edu.tw udp port 33447 unreachable, length 36
IP bsd1.cs.nctu.edu.tw > nabsd: ICMP bsd1.cs.nctu.edu.tw udp port 33448 unreachable, length 36
IP bsd1.cs.nctu.edu.tw > nabsd: ICMP bsd1.cs.nctu.edu.tw udp port 33449 unreachable, length 36
```

Traceroute Program (5)

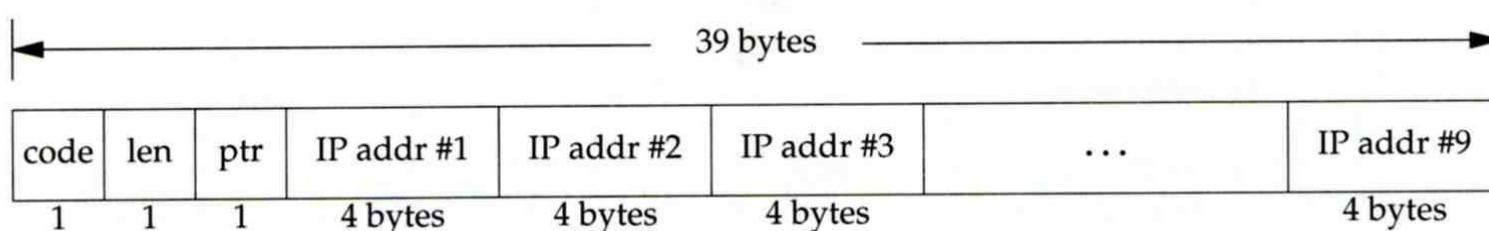
- ❑ The router IP in traceroute is the interface that receives the datagram. (incoming IP)
 - Traceroute from left host to right host
 - if1, if3
 - Traceroute from right host to left host
 - if4, if2



Traceroute Program – IP Source Routing Option (1)



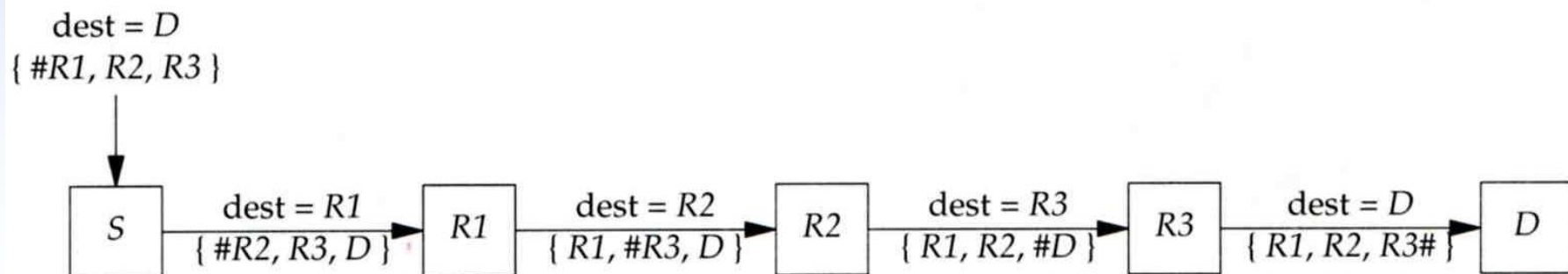
- ❑ Source Routing
 - Sender specifies the route
- ❑ Two forms of source routing
 - Strict source routing
 - Sender specifies the **exact path** that the IP datagram must follow
 - Loose source routing
 - As strict source routing, but the datagram can pass through other routers between any two addresses in the list
- ❑ Format of IP header option field
 - Code = 0x89 for strict and code = 0x83 for loose SR option



Traceroute Program – IP Source Routing Option (2)

❑ Scenario of source routing

- Sending host
 - Remove first entry and append destination address in the final entry of the list
- Receiving router $\neq$ destination
 - Loose source route, forward it as normal
- Receiving router = destination
 - Next address in the list becomes the destination
 - Change source address
 - Increment the pointer



Traceroute Program – IP Source Routing Option (3)

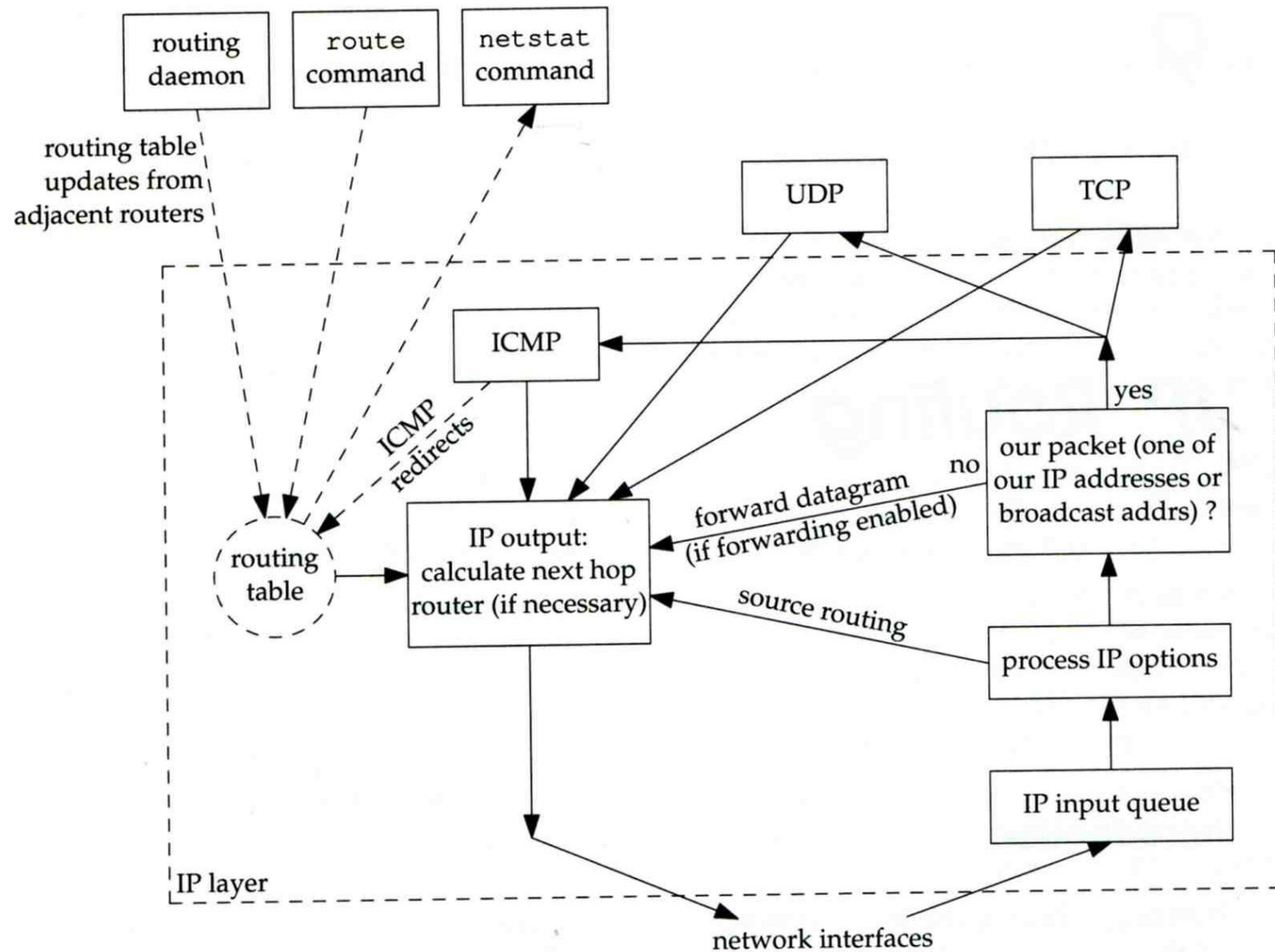
❑ Traceroute using IP loose SR option

❑ Ex:

```
nabsd [/home/chwong] -chwong- traceroute u2.nctu.edu.tw
traceroute to u2.nctu.edu.tw (211.76.240.193), 64 hops max, 40 byte packets
 1 e3rtn-235 (140.113.235.254) 0.549 ms 0.434 ms 0.337 ms
 2 140.113.0.166 (140.113.0.166) 108.726 ms 4.469 ms 0.362 ms
 3 v255-194.NTCU.net (211.76.255.194) 0.529 ms 3.446 ms 5.464 ms
 4 v255-229.NTCU.net (211.76.255.229) 1.406 ms 2.017 ms 0.560 ms
 5 h240-193.NTCU.net (211.76.240.193) 0.520 ms 0.456 ms 0.315 ms
nabsd [/home/chwong] -chwong- traceroute -g 140.113.0.149 u2.nctu.edu.tw
traceroute to u2.nctu.edu.tw (211.76.240.193), 64 hops max, 48 byte packets
 1 e3rtn-235 (140.113.235.254) 0.543 ms 0.392 ms 0.365 ms
 2 140.113.0.166 (140.113.0.166) 0.562 ms 9.506 ms 0.624 ms
 3 140.113.0.149 (140.113.0.149) 7.002 ms 1.047 ms 1.107 ms
 4 140.113.0.150 (140.113.0.150) 1.497 ms 6.653 ms 1.595 ms
 5 v255-194.NTCU.net (211.76.255.194) 1.639 ms 7.214 ms 1.586 ms
 6 v255-229.NTCU.net (211.76.255.229) 1.831 ms 9.244 ms 1.877 ms
 7 h240-193.NTCU.net (211.76.240.193) 1.440 ms !S 2.249 ms !S 1.737 ms !S
```

IP Routing

– Processing in IP Layer



IP Routing

– Routing Table (1)

□ Routing Table

- Command to list: `netstat -rn`
- Flag
 - U: the route is up
 - G: the route is to a router (indirect route)
 - Indirect route: IP is the dest. IP, MAC is the router's MAC
 - H: the route is to a host (Not to a network)
 - The dest. filed is either an IP address or network address
 - S: the route is static
- Expire: expiration time for each route

```
nasa [/home/wangth] -wangth- netstat -rn
Routing tables
```

Internet:

Destination	Gateway	Flags	Netif	Expire
Default	140.113.17.254	UGS	em0	
127.0.0.1	link#2	UH	lo0	
140.113.17.0/24	link#1	U	em0	
140.113.17.225	link#1	UHS	lo0	

IP Routing

– Routing Table (2)

❑ Ex:

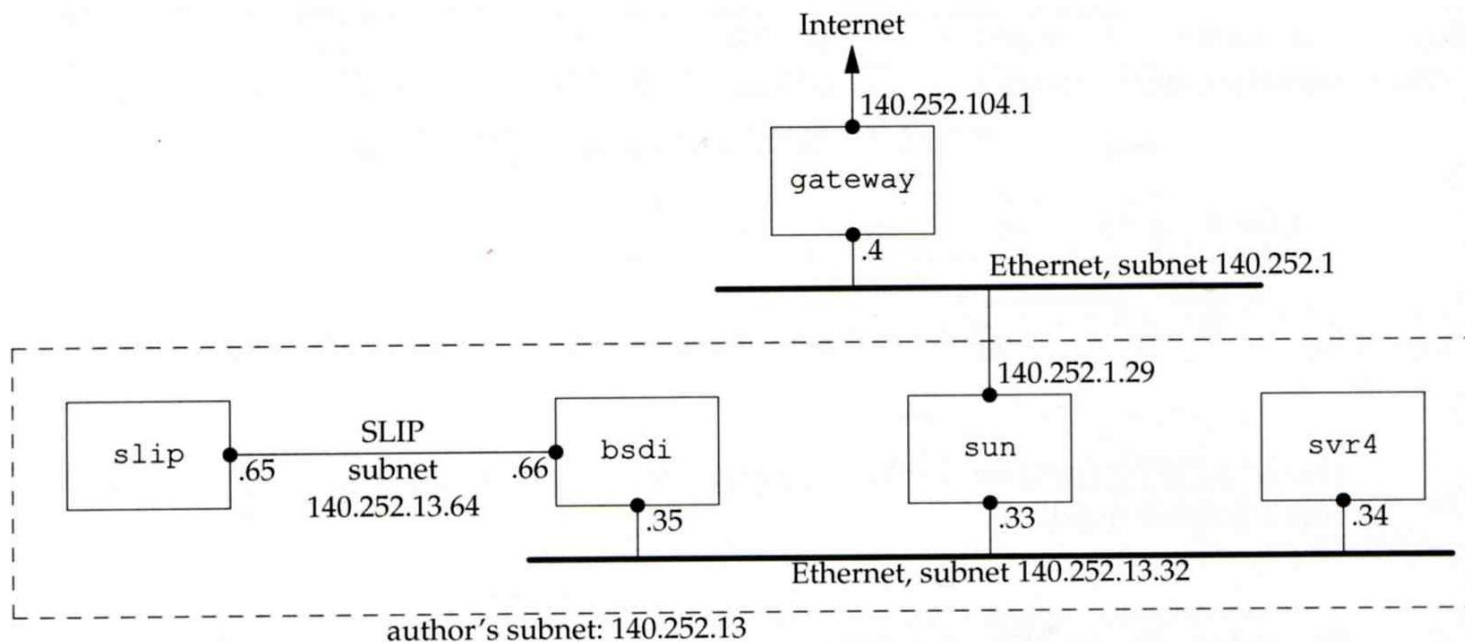
1. dst. = sun
2. dst. = slip
3. dst. = 192.207.117.2
4. dst. = svr4 or 140.252.13.34
5. dst. = 127.0.0.1

```
svr4 % netstat -rn
```

Routing tables

Destination	Gateway	Flags	Refcnt	Use	Interface
140.252.13.65	140.252.13.35	UGH	0	0	emd0
127.0.0.1	127.0.0.1	UH	1	0	lo0
default	140.252.13.33	UG	0	0	emd0
140.252.13.32	140.252.13.34	U	4	25043	emd0

loopback

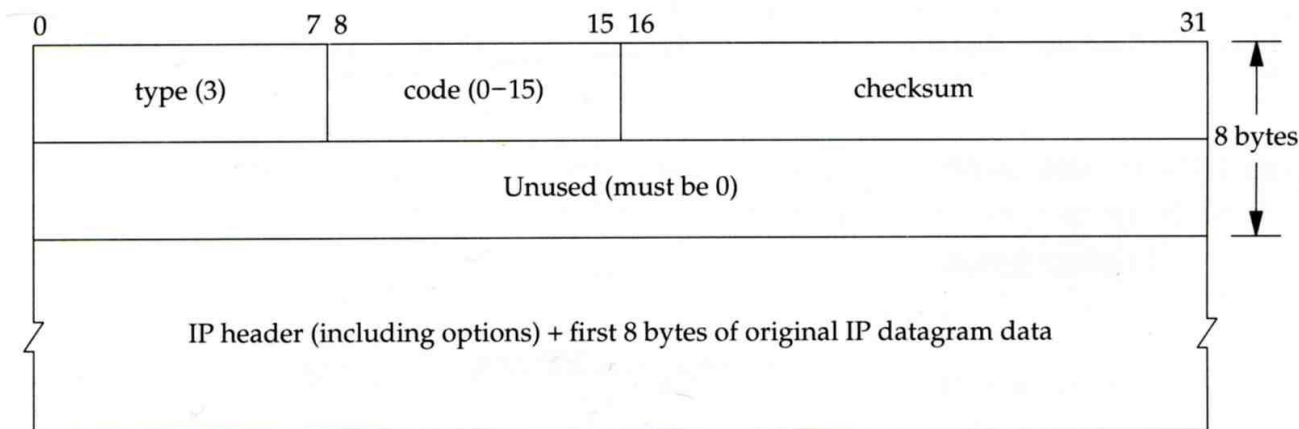


ICMP

– No Route to Destination

❑ If there is no match in routing table

- If the IP datagram is generated on the host
 - “host unreachable” or “network unreachable”
- If the IP datagram is being forwarded
 - ICMP “host unreachable” error message is generated and sends back to sending host
 - ICMP message
 - Type = 3, code = 0 for host unreachable
 - Type = 3, code = 1 for network unreachable



ICMP

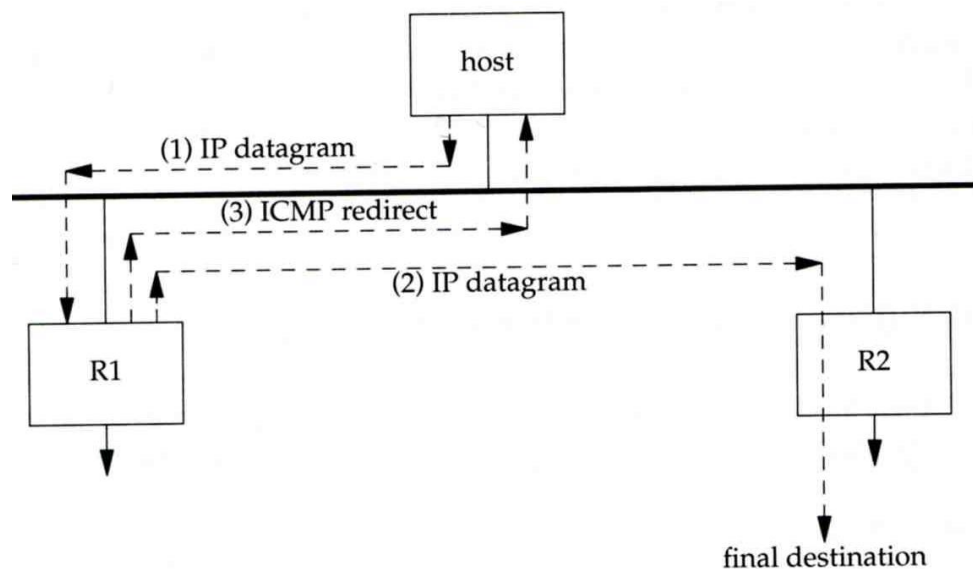
– Redirect Error Message (1)

❑ Concept

- Used by router to inform the sender that the datagram should be sent to a different router
- This will happen if the host has a choice of routers to send the packet to

➤ Ex:

- R1 found sending and receiving interface are the same

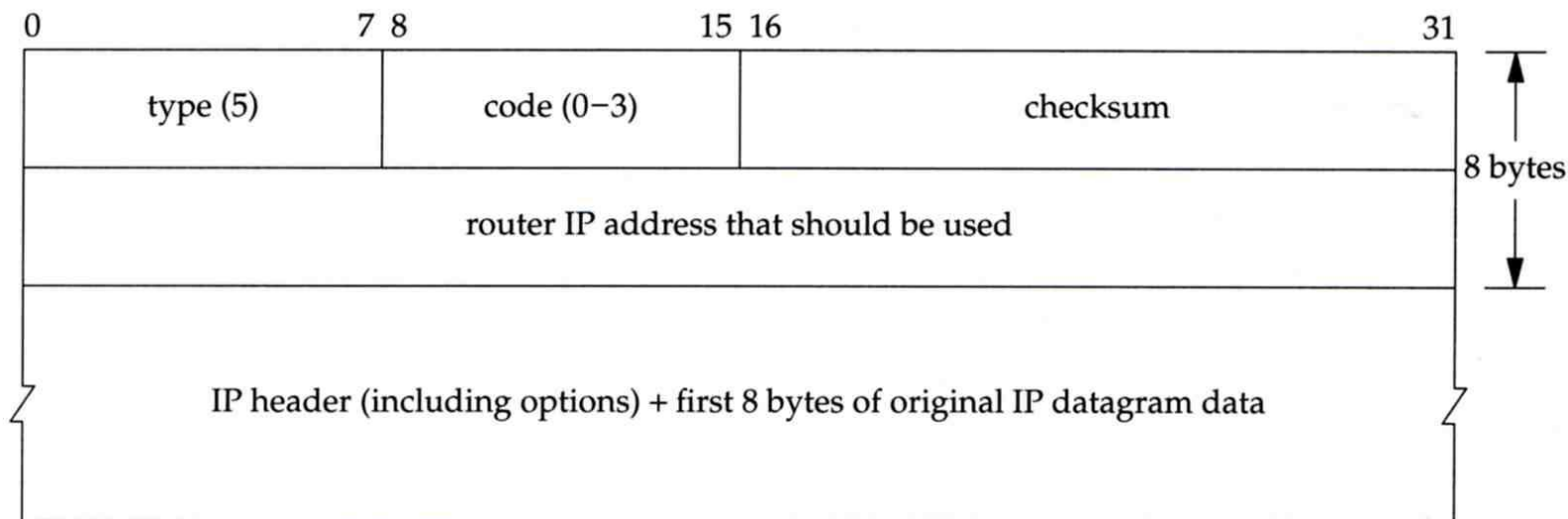


ICMP

– Redirect Error Message (2)

❑ ICMP redirect message format

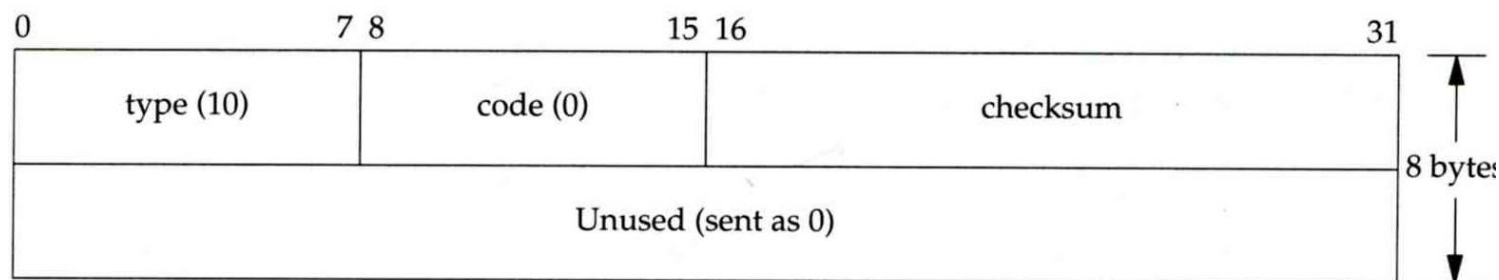
- Code 0: redirect for network
- Code 1: redirect for host
- Code 2: redirect for TOS and network (RFC 1349)
- Code 3: redirect for TOS and hosts (RFC 1349)



ICMP

– Router Discovery Messages (1)

- ❑ Dynamic update host's routing table
 - ICMP router solicitation message (懇求)
 - Host broadcast or multicast after bootstrapping
 - ICMP router advertisement message
 - Router response
 - Router periodically broadcast or multicast
- ❑ Format of ICMP router solicitation message

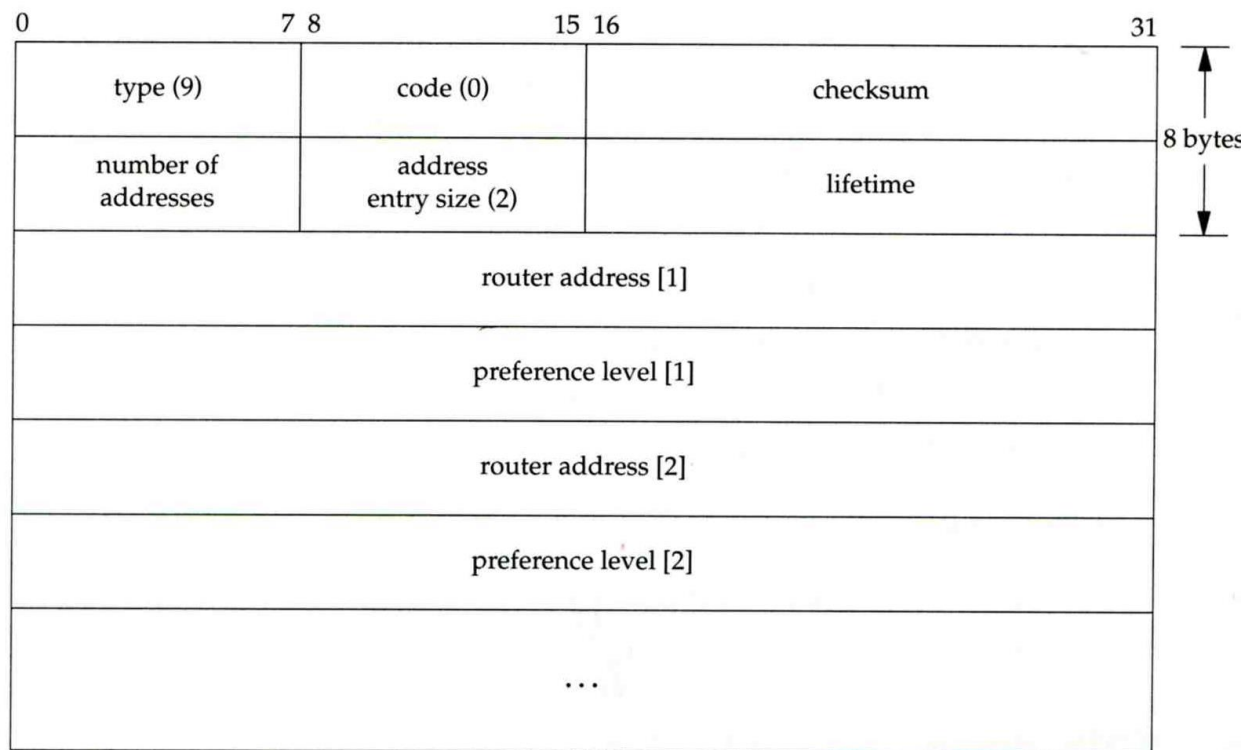


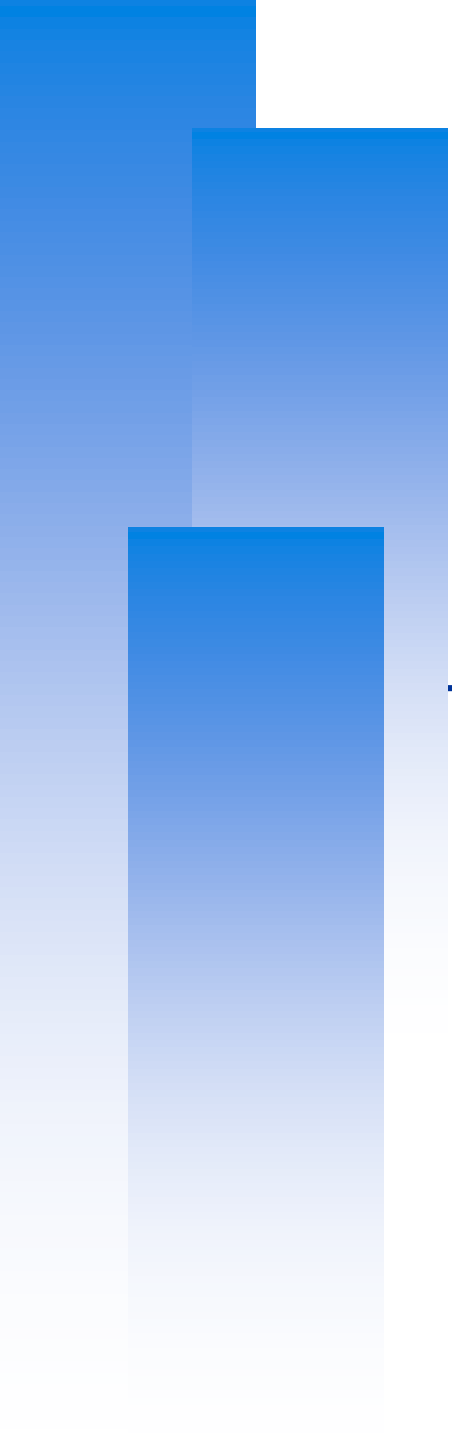
ICMP

– Router Discovery Messages (2)

□ Format of ICMP router advertisement message

- Router address
 - Must be one of the router's IP address
- Preference level
 - Preference as a default router address





UDP – User Datagram Protocol

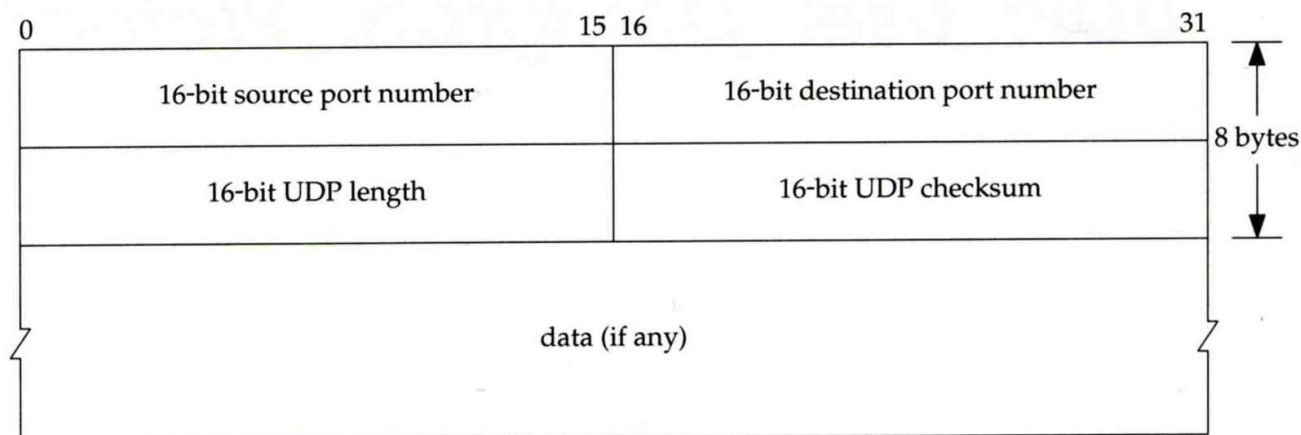
UDP

❑ No reliability

- Datagram-oriented, not stream-oriented protocol

❑ UDP header

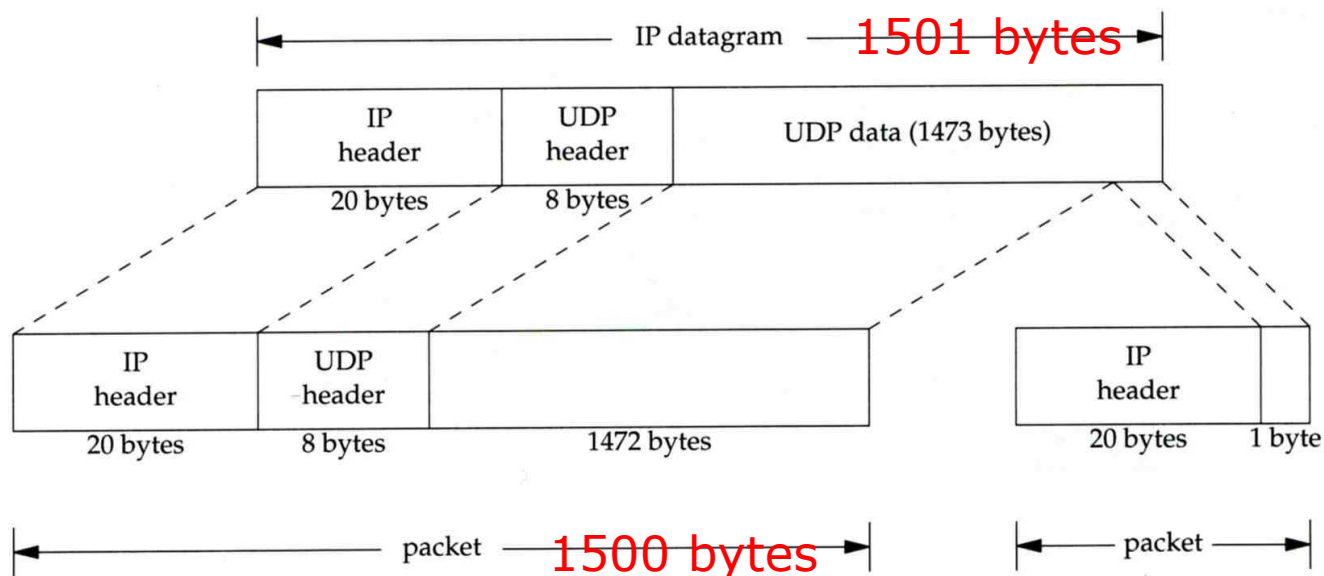
- 8 bytes
 - Source port and destination port
 - Identify sending and receiving process
 - UDP length: ≥ 8



IP Fragmentation (1)

❑ MTU limitation

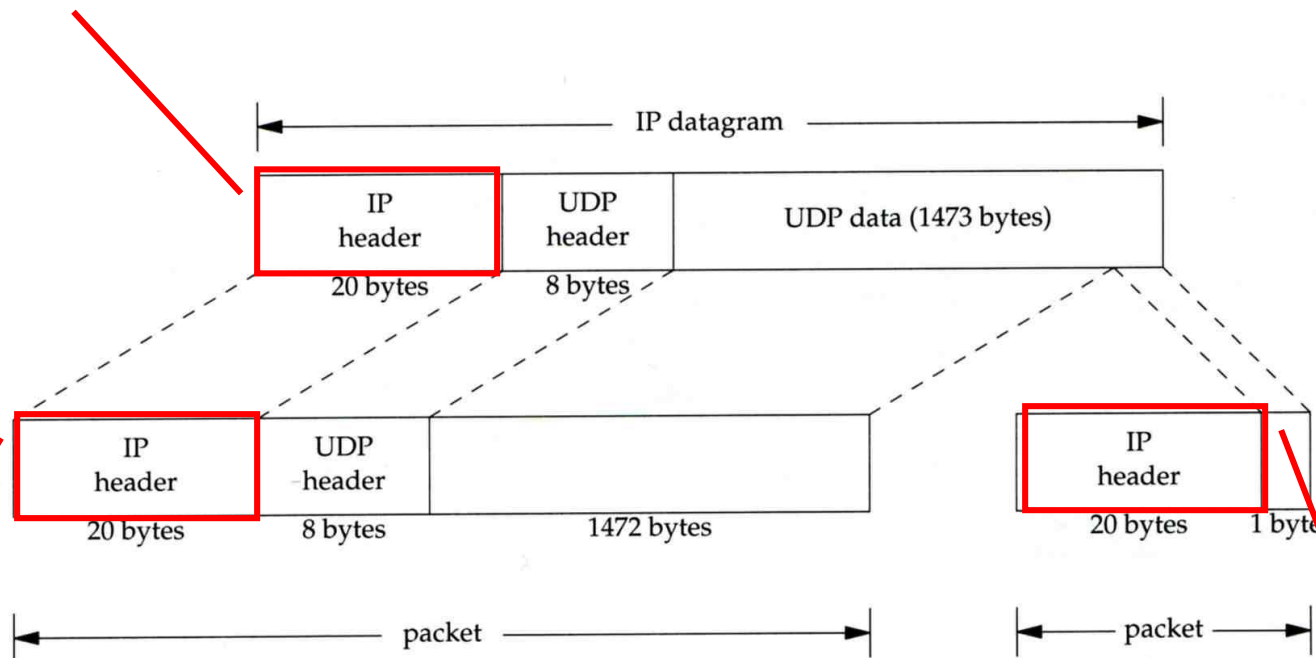
- Before network-layer to link-layer
 - IP will check the size and link-layer MTU
 - Do fragmentation if necessary
- Fragmentation may be done at sending host or routers
- Reassembly is done only in receiving host



IP Fragmentation (2)

identification:
flags:
fragment offset

which unique IP datagram
more fragments?
offset of this datagram from the beginning of original datagram



identification:
flags:
fragment offset

the same
more fragments
0

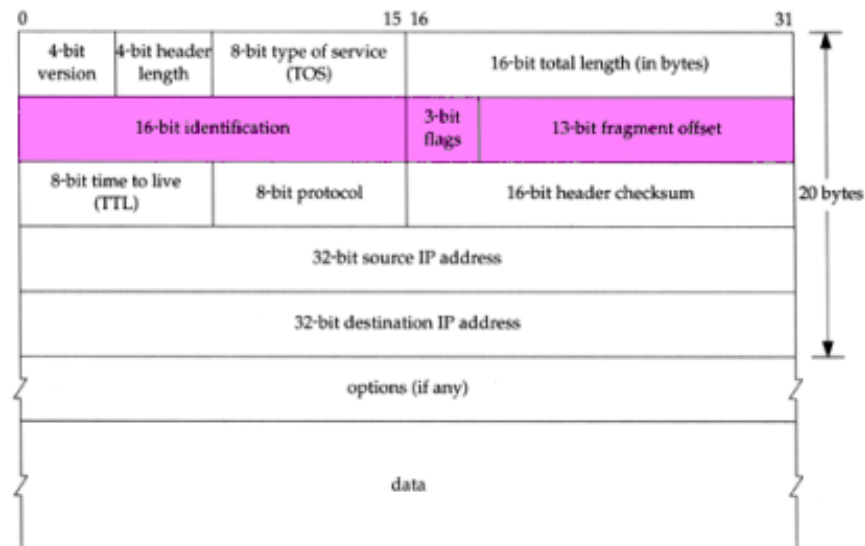
identification:
flags:
fragment offset

the same
end of fragments
1480

IP Fragmentation (3)

❑ Issues of fragmentation

- One fragment lost, entire datagram must be retransmitted
- If the fragmentation is performed by intermediate router, there is no way for sending host how fragmentation did
- Fragmentation is often avoided
 - There is a “don’t fragment” bit in flags of IP header

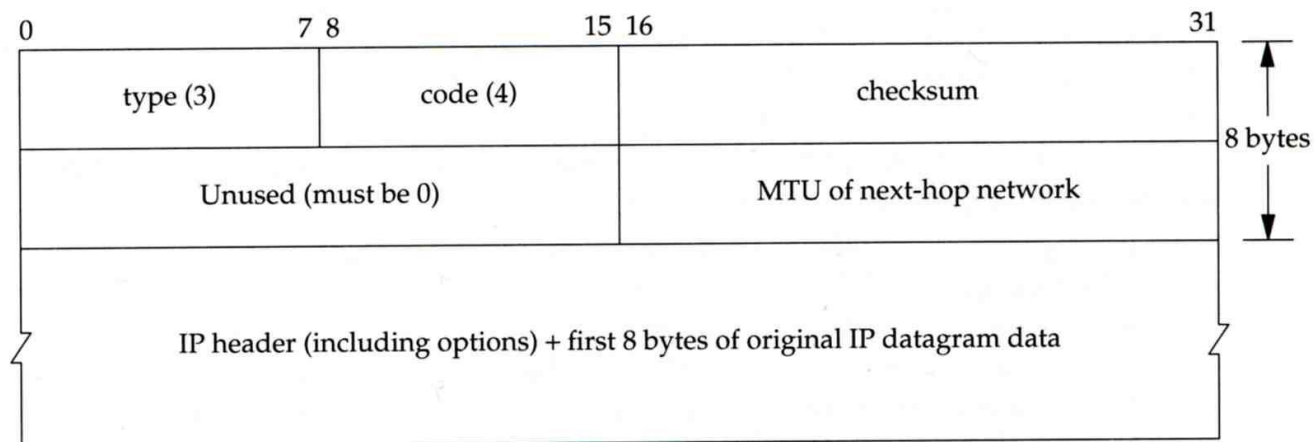


ICMP Unreachable Error – Fragmentation Required

❑ Type=3, code=4

- Router will generate this error message if the datagram needs to be fragmented, but the “don’t fragment” bit is turn on in IP header

❑ Message format

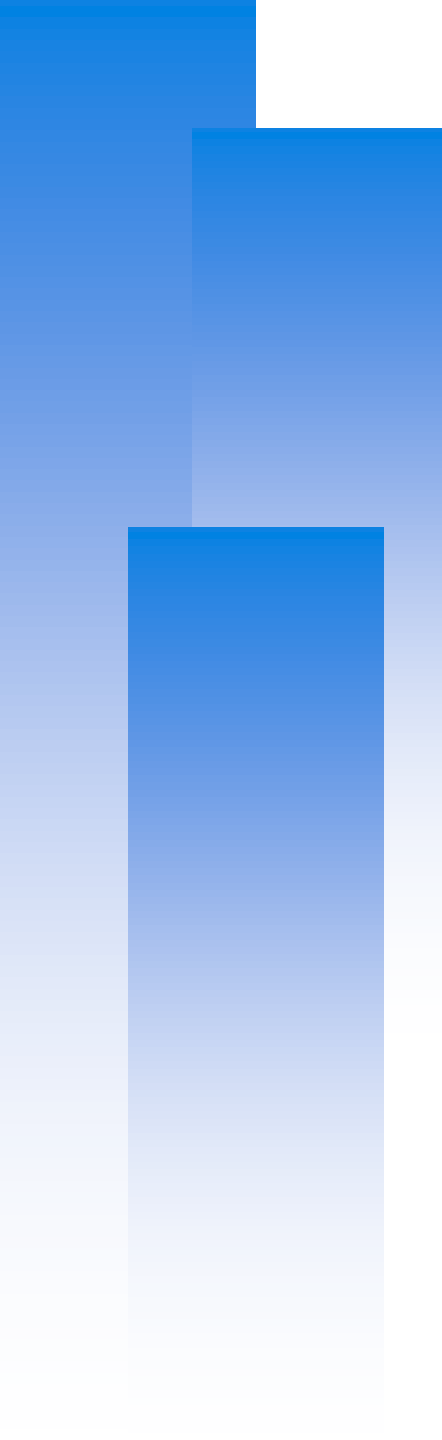


ICMP

– Source Quench Error

❑ Type=4, code=0

- May be generated by system when it receives datagram at a rate that is too fast to be processed
- Host receiving more than it can handle datagram
 - Send ICMP source quench or
 - Throw it away
- Host receiving UDP source quench message
 - Ignore it or
 - Notify application



TCP – Transmission Control Protocol

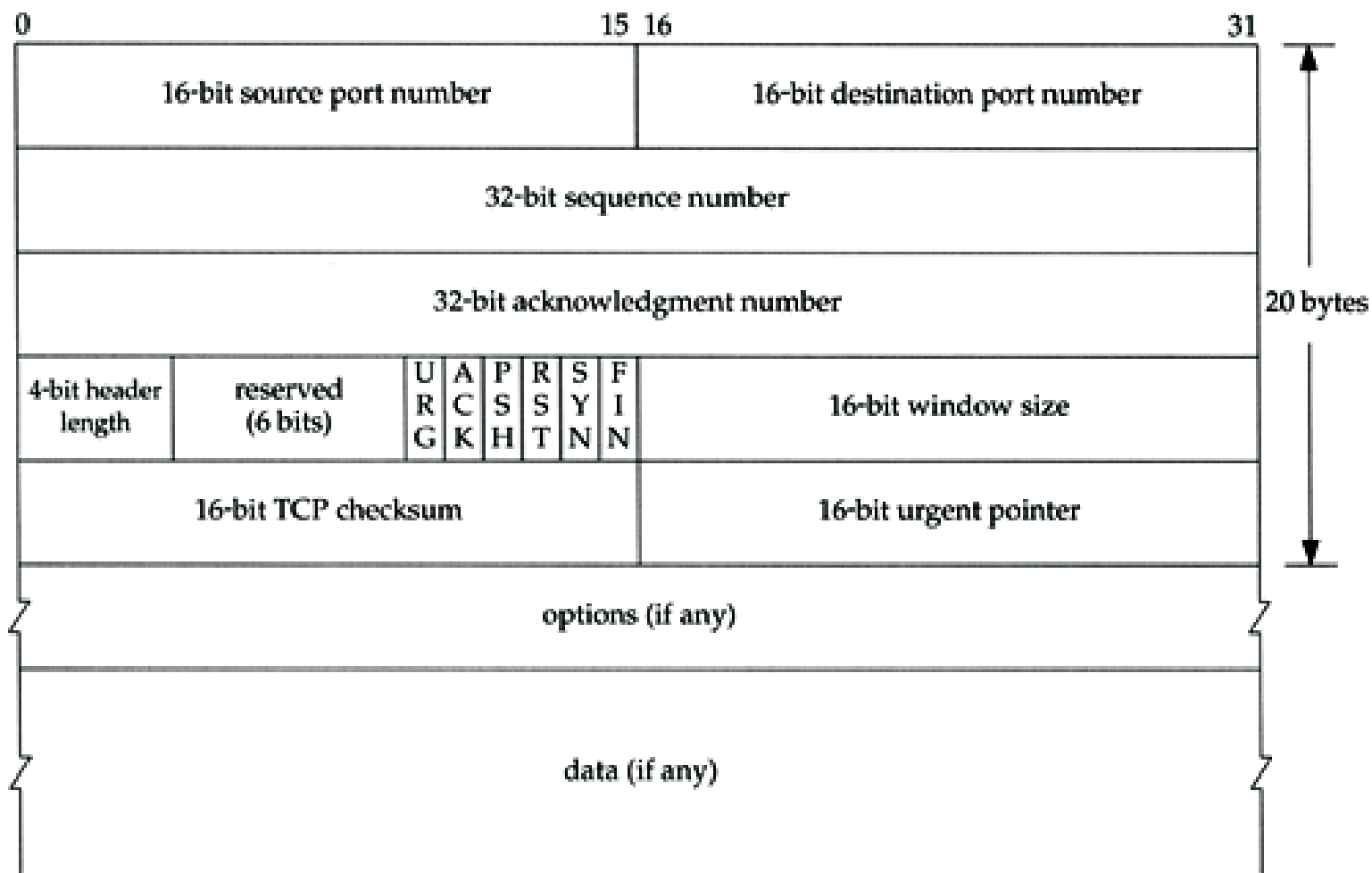
TCP

❑ Services

- Connection-oriented
 - Establish TCP connection before exchanging data
- Reliability
 - Acknowledgement when receiving data
 - Retransmission when timeout
 - Ordering
 - Discard duplicated data
 - Flow control

TCP

– Header (1)



TCP

– Header (2)

❑ Flags

- SYN
 - Establish new connection
- ACK
 - Acknowledgement number is valid
 - Used to ack previous data that host has received
- RST
 - Reset connection
- FIN
 - The sender is finished sending data

TCP connection establishment and termination

↕
Three-way handshake

TCP's half close

