



# Lecture 4-2: C Memory Management

## **CS10014 Computer Organization**

Department of Computer Science

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Thursday: 1:20 pm– 3:10 pm

Classroom: EC-022



# Acknowledgements and Disclaimer

- Slides were developed in the reference with
  - CS 61C at UC Berkeley
    - <https://inst.eecs.berkeley.edu/~cs61c/sp23/>
  - CS 252 at UC Berkeley
    - <https://people.eecs.berkeley.edu/~culler/courses/cs252-s05/>
  - CSCE 513 at University of South Carolina
    - <https://passlab.github.io/CSCE513/>



# Outline

- C Memory Management



# C Memory Management(1/3)

- A program's **address space** contains 4 regions

- **Stack:**

- Local variables, grows downward

- **Heap:**

- Space requested for pointers via malloc()
- Resizes dynamically
- Grows upward

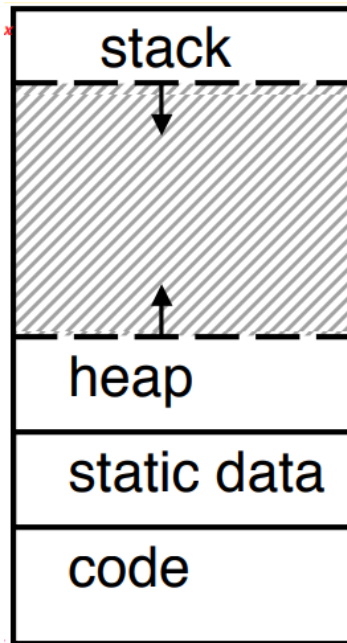
- **Static data:**

- Variables declared outside main
- Does not grow or shrink

- **Code**

- Load when the program starts, does not change

$\sim \text{FFFF FFFF}_{\text{hex}}$



$\sim 0_{\text{hex}}$



## C Memory Management(2/3)

- Variable declaration **does** allocate memory
  - If declare outside a procedure, allocated in static storage
  - If declare inside procedure allocated on the stack and **free when procedure returns**
    - main() is a procedure

```
int myGlobal;  
main() {  
    int myTemp;  
}
```



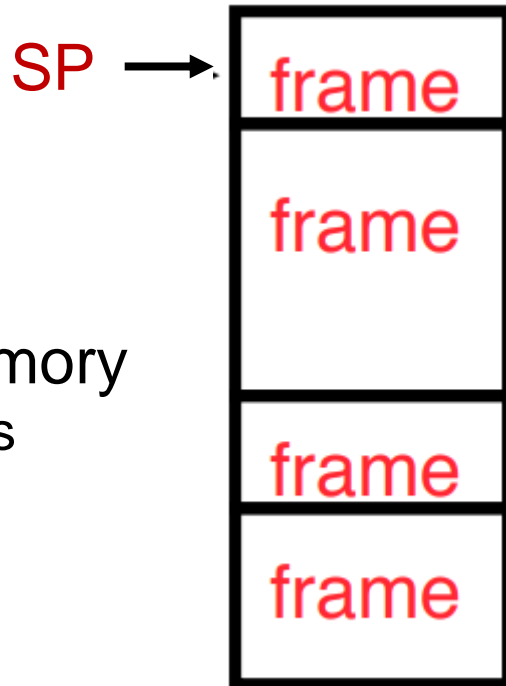
# C Memory Management(3/3)

- C has 3 pools of memory
  - **Static storage:** global variable storage
    - basically permanent, entire program run
  - **The stack:** local variable storage
    - Parameters, return address
    - Stack frame in C
  - **The Heap:** dynamic storage
    - Malloc()
    - Data lives until deallocated by the programmer
  - C requires knowing where objects are in memory otherwise things don't work as expected



# The Stack (1/1)

- Stack frame includes:
  - Return address
  - Parameters
  - Space for other local variables
- Stack frames contiguous blocks of memory
  - Stack pointer tells where top stack frame is
- When procedure ends
  - Stack frame is tossed off the stack
  - Free memory for future stack frame

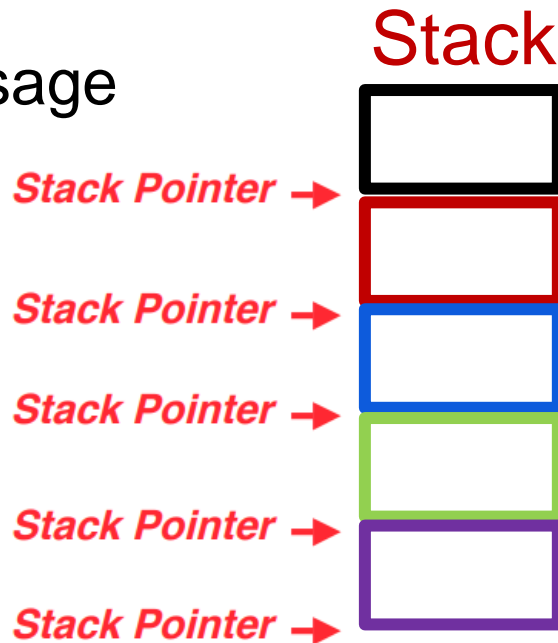




# The Stack (2/2)

- Last in, first out (LIFO) memory usage

```
main ()
{ a(0);
}
void a (int m)
{ b(1);
}
void b (int n)
{ c(2);
}
void c (int o)
{ d(3);
}
void d (int p)
{
}
```







# The Heap (Dynamic Memory)

- Large pool of memory, **not** allocated in contiguous order
- In C, specify the number of **bytes** of memory explicitly to allocate item

```
int *ptr;
ptr = (int *) malloc(sizeof(int));
/* malloc returns type (void *),
so need to cast to right type */
```

- **malloc ()**: Allocates raw, uninitialized memory from heap



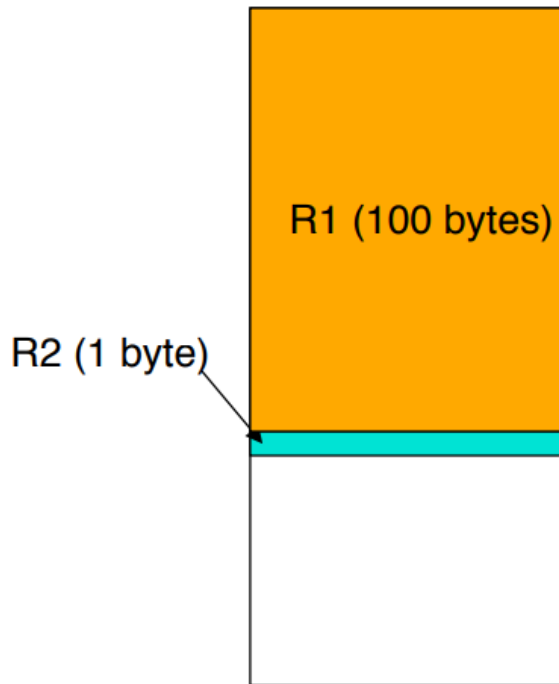
## So far ...

- How do we manage memory?
  - **Code, Static storage are easy:**
    - They never grow or shrink
  - **Stack space is also easy:**
    - Stack frames are created
    - Destroyed in last-in, first-out (LIFO) order
  - **Managing the heap is tricky:**
    - Memory can be allocated/deallocated at any time



# Heap Management (1/2)

- An example
  - Request R1 for 100 bytes
  - Request R2 for 1 byte
  - Memory from R1 is freed
  - Request R3 for 50 bytes





# Heap Management (2/2)

- An example
  - Request R1 for 100 bytes
  - Request R2 for 1 byte
  - Memory from R1 is freed
  - Request R3 for 50 bytes

Where should we place the R3?

