

# Wireless Communication Systems

## @CS.NCTU

Lecture 12: Soft Information

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# PPR: Partial Packet Recovery for Wireless Networks

ACM SIGOCMM, 2017

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CSAIL, MIT

# What is Partial Packet Error?

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Lots of packets lost due to collisions and noise in wireless networks

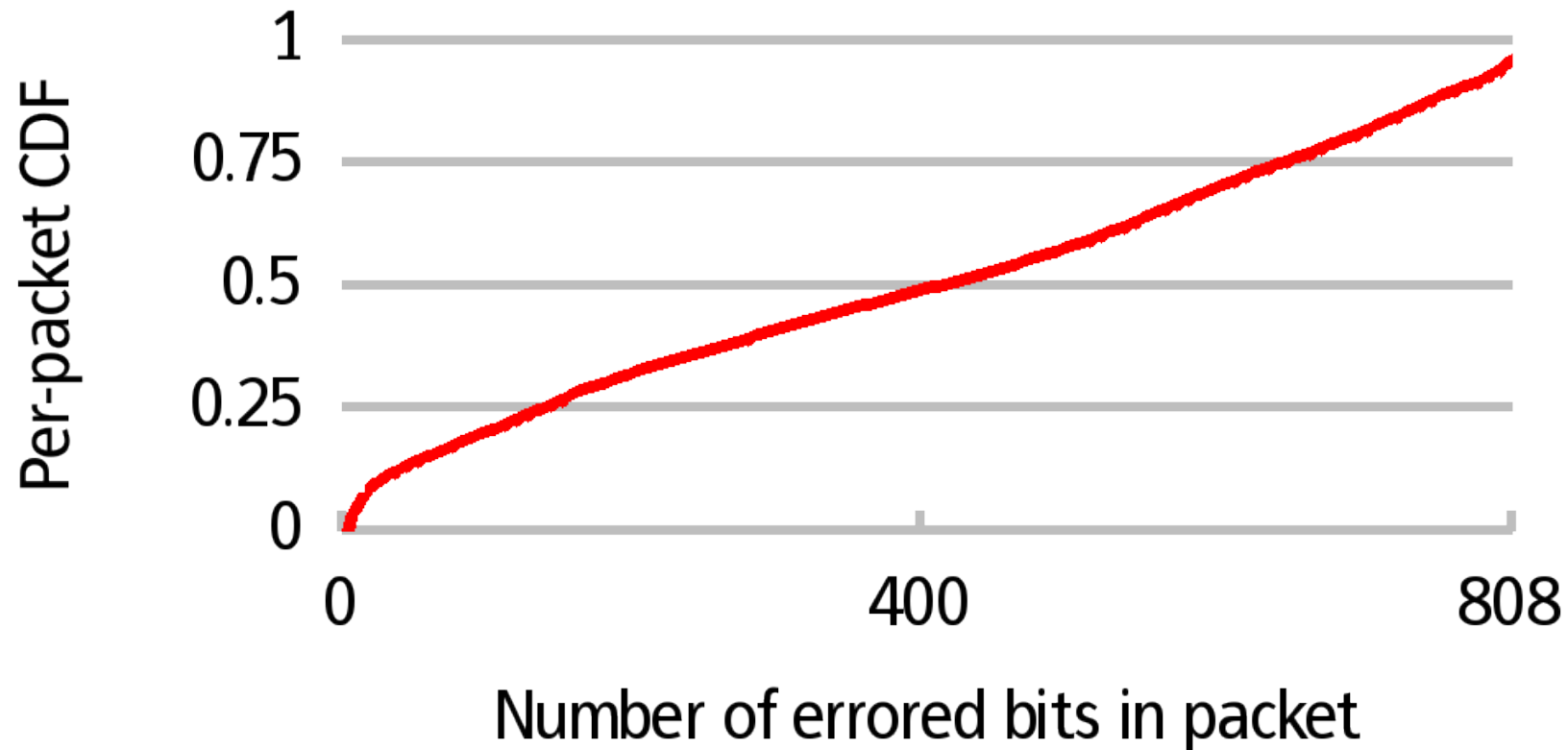


Can't receive non-colliding bits today!

# Bits in a packet don't share fate

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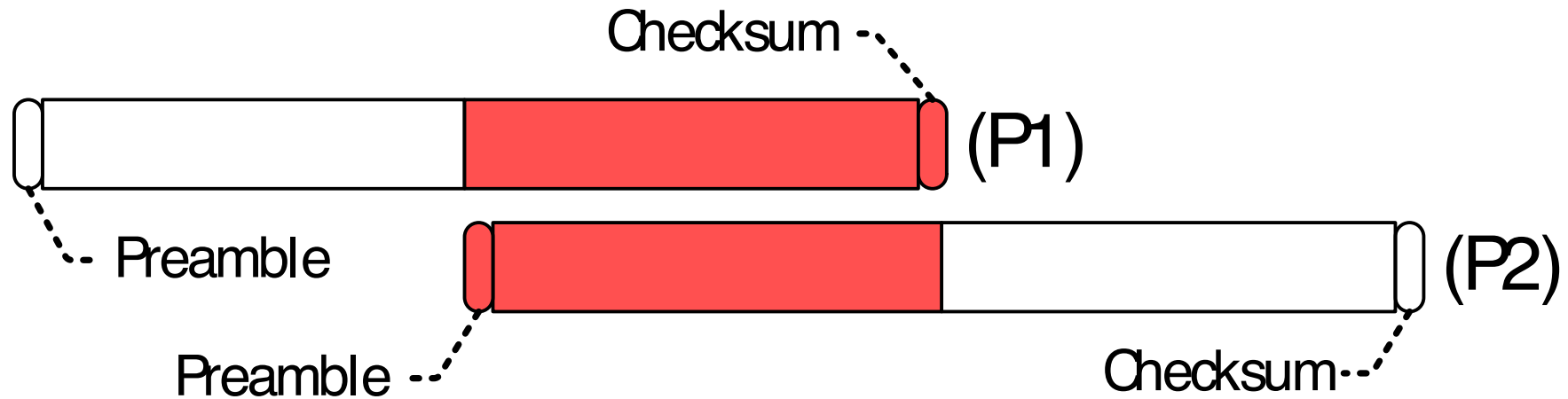
(30 node testbed, CSMA on)



Many bits from corrupted packets are correct,  
but status quo receivers don't know which!

# Three Key Questions

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1. How does receiver know which bits are correct?
2. How does receiver know P2 is there at all?
3. How to design an efficient ARQ protocol?

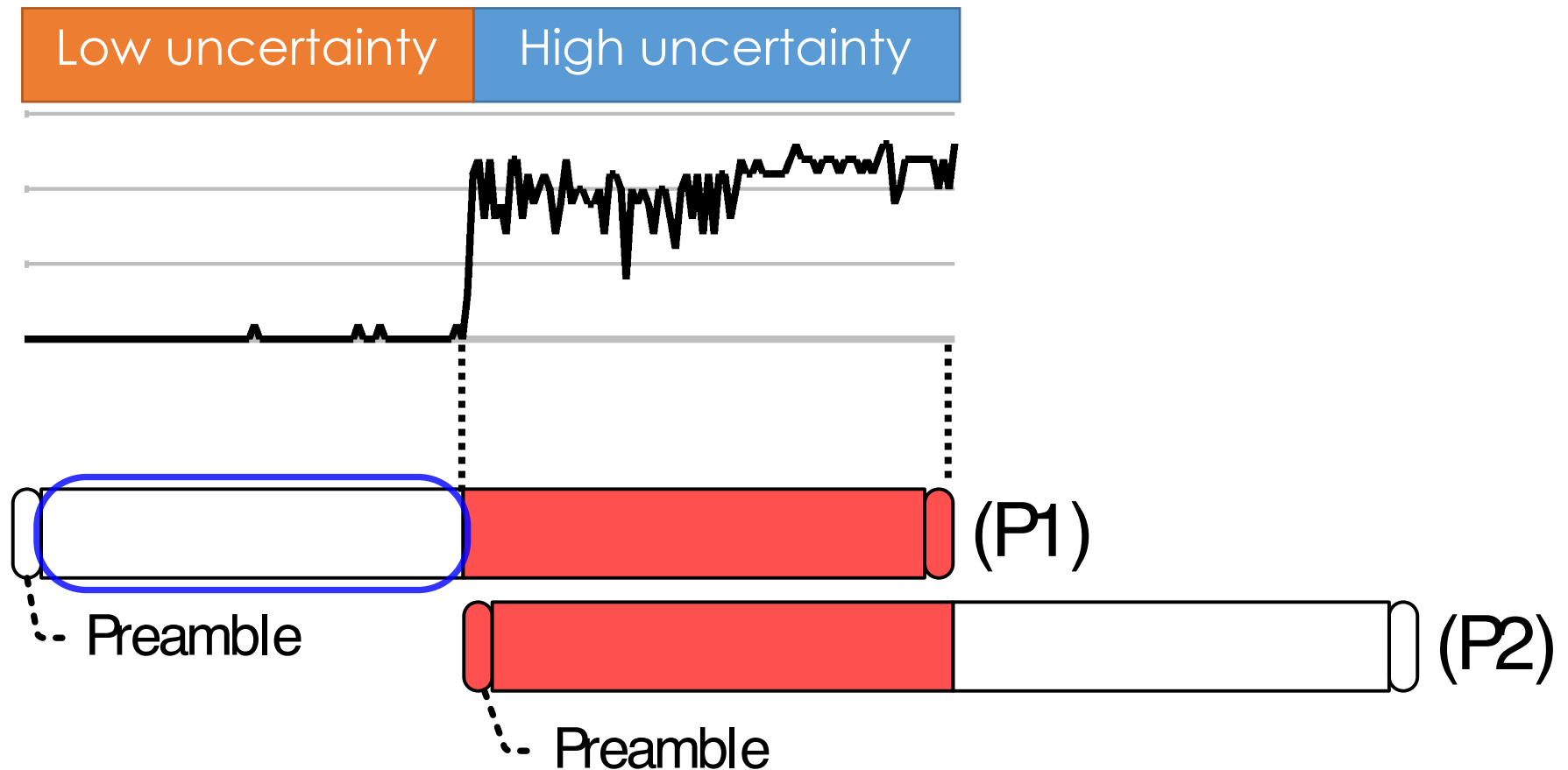
# Can Receiver Identify Correct Bits?

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- Use physical layer (PHY) hints: **SoftPHY**
  - Receiver PHY has the information!
  - Pass this **confidence information** to higher layer as a hint
- SoftPHY implementation is PHY-specific; interface is PHY-independent
- Implemented for direct sequence spread spectrum (DSSS) over MSK and other modulations

# Can We Leverage Soft Info?

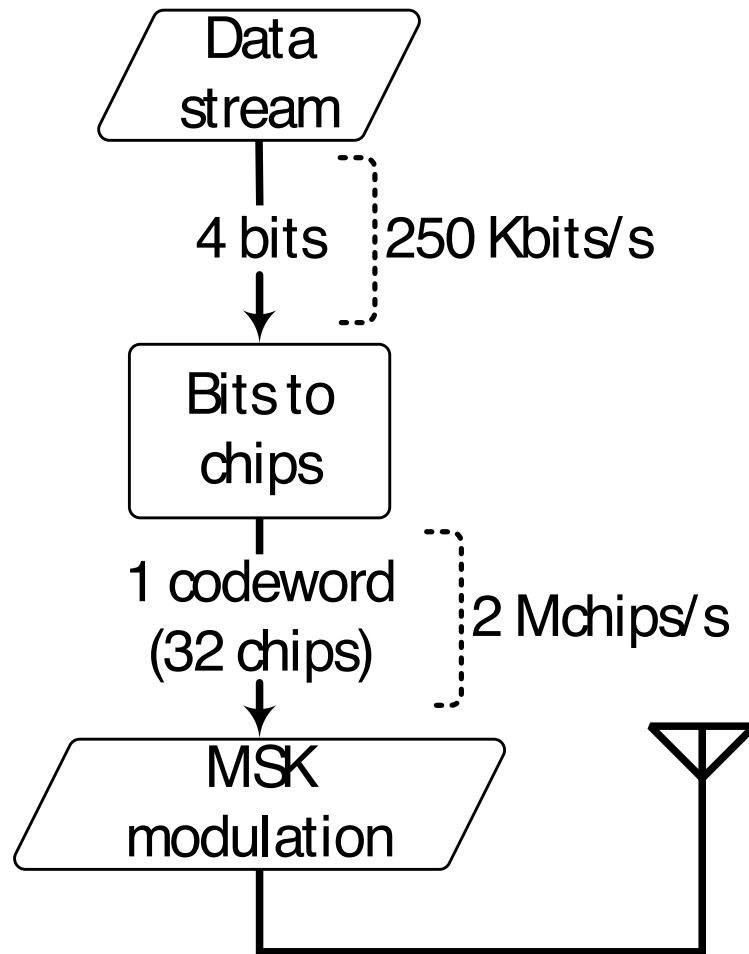
PHY conveys uncertainty in each bit it delivers up



# Direct Sequence Spread Spectrum

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## Transmitter:



## Receiver:

- Demodulate MSK signal
- Decide on closest codeword to received (Hamming distance)
- Many 32-bit chip sequences are not valid codewords
- Codewords separated by at least 11 in Hamming distance
- 802.11 similar



# SoftPHY Hint for Spread Spectrum

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**Hamming distance** between **received** chips and **decided-upon** codeword

Receive: 11101101000111000011010110100010  
 $C_1$ : 11101101100111000011010100100010  
→ SoftPHY hint is **2**

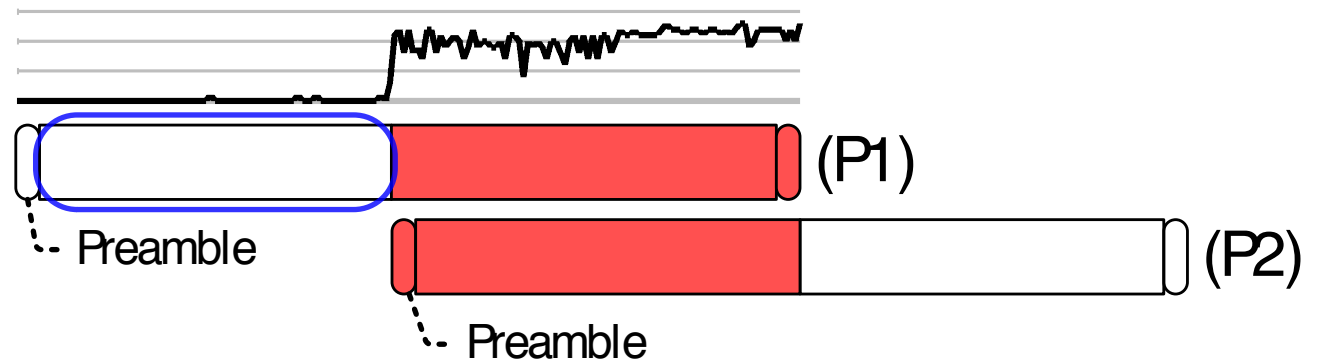
Receive: 11001101000111010111011110110111  
 $C_1$ : 11101101100111000011010100100010  
→ SoftPHY hint is **9**

# Three Key Questions

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1. How does receiver know which bits are correct?

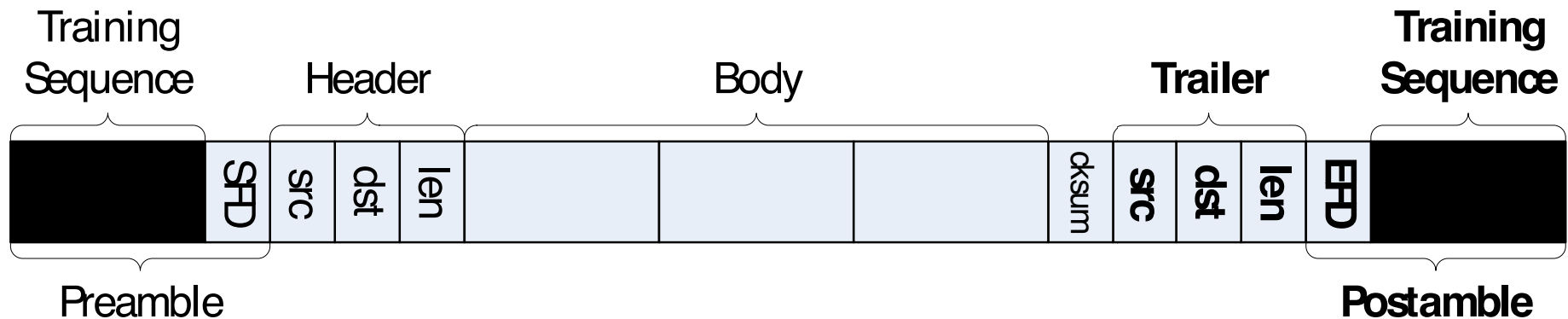
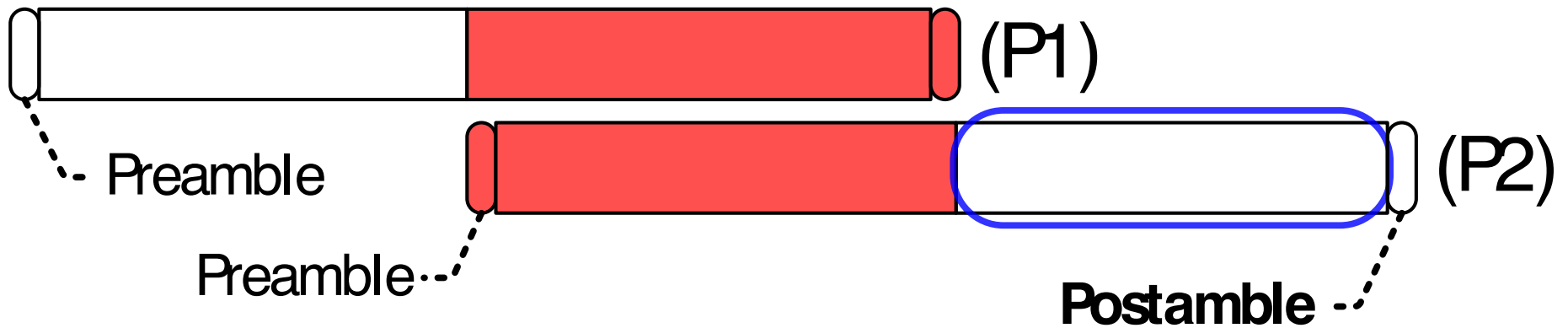
**A: SoftPHY:**



- 2. How does receiver know P2 is there at all?**
3. How to design an efficient ARQ protocol?

# Postamble decoding

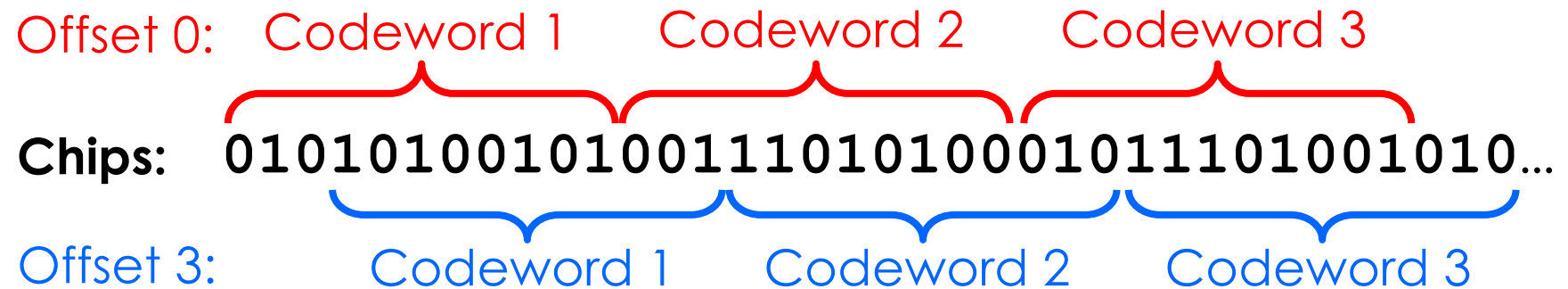
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# Receiver Design with Postamble

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- Codeword synchronization
  - Translate stream of chips to codewords
  - Search for **postamble** at all chip offsets

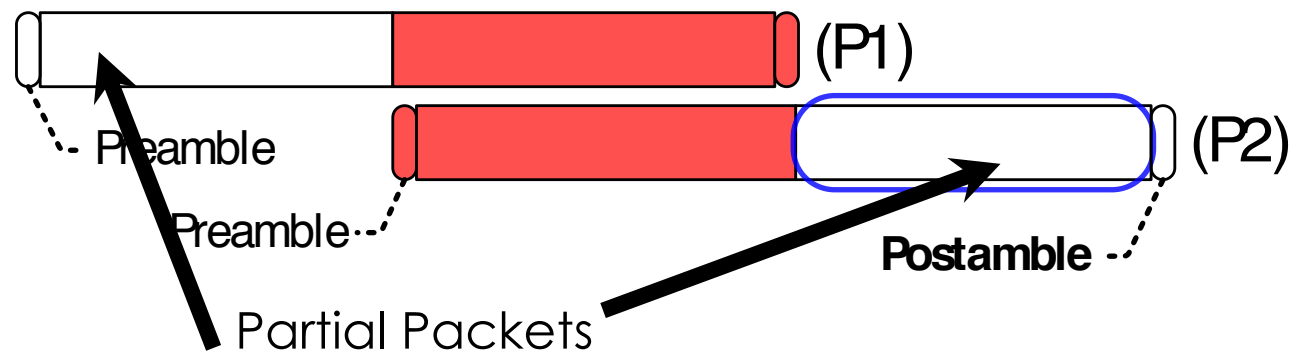


# Three Key Questions

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1. How does receiver know which bits are correct?
2. How does receiver know P2 is there at all?

**A: Postamble:**

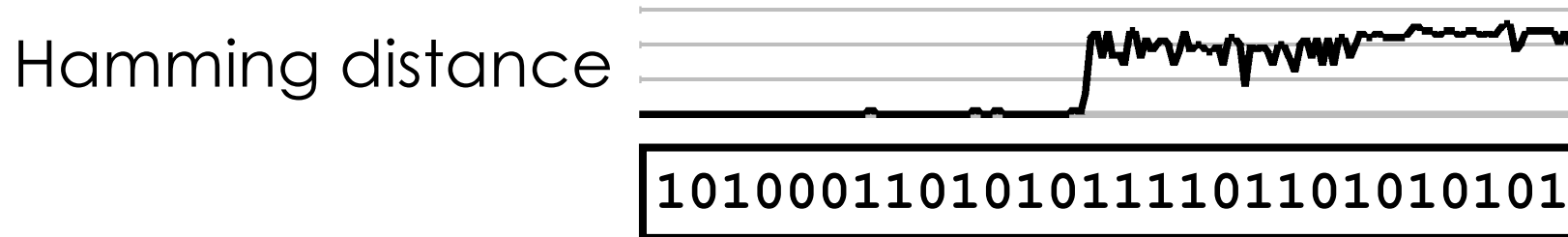


3. **How to design an efficient ARQ protocol?**

# ARQ with partial packets

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- ARQ today: correctly-received bits get resent
- **PP-ARQ** key idea: resend **only** incorrect bits

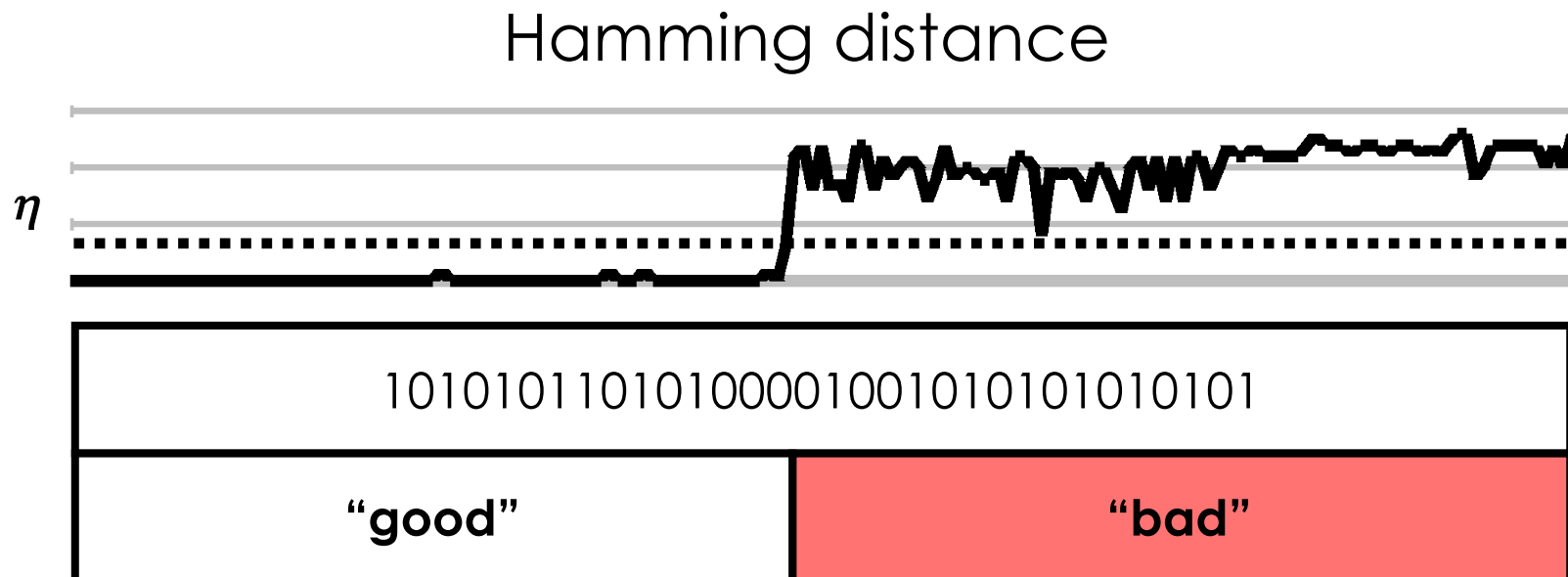


- **Efficiently** tell sender about what happened
  - Feedback packet

# Labeling Bits “good” or “bad”

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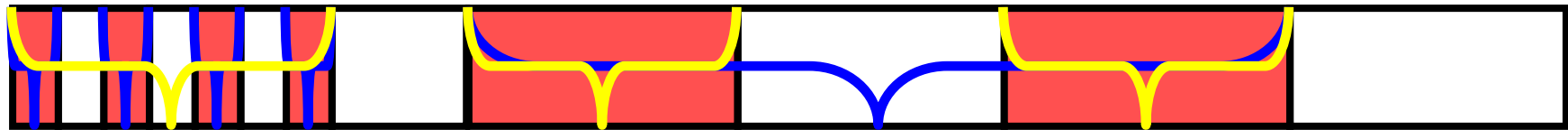
- **Threshold test:** pick a threshold  $\eta$ 
  - Label codewords with SoftPHY hint  $> \eta$  “**bad**”
  - Label codewords with SoftPHY hint  $\leq \eta$  “**good**”



# PP-ARQ protocol

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1. Assuming hints correct, which ranges to ask for?
  - Dynamic programming problem
  - Forward and feedback channels



□ "Good" bits  
■ "Bad" bits

2. Codewords are in fact **correct** or **incorrect**

- Two possibilities for mistakes
  - Labeling a correct codeword "bad"
  - Labeling an incorrect codeword "good"



# Implementation

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**Sender:** telos tmote sky sensor node

- Radio: CC2420 DSSS/MSK (Zigbee)
- Modified to send **postambles**



**Receiver:** USRP software radio with 2.4 GHz RFX 2400 daughterboard

- Despreading, postamble synchronization, demodulation
- **SoftPHY** implementation

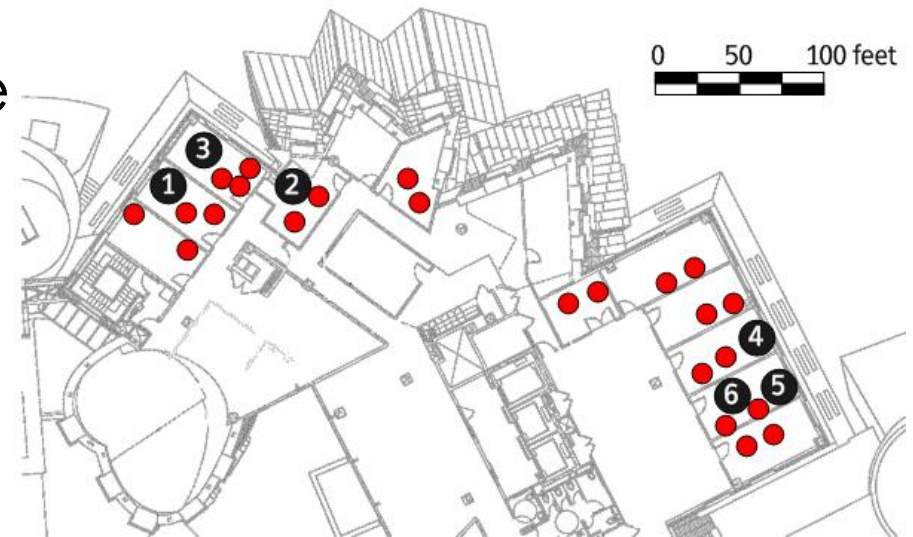
**PP-ARQ:** trace-driven simulation

# Experimental design

- Live wireless testbed experiments

- Senders transmit 101-byte packets, varying traffic rate
- Evaluate raw PPR throughput
- Evaluate SoftPHY and postamble improvements

- 25 senders
- 6 receivers



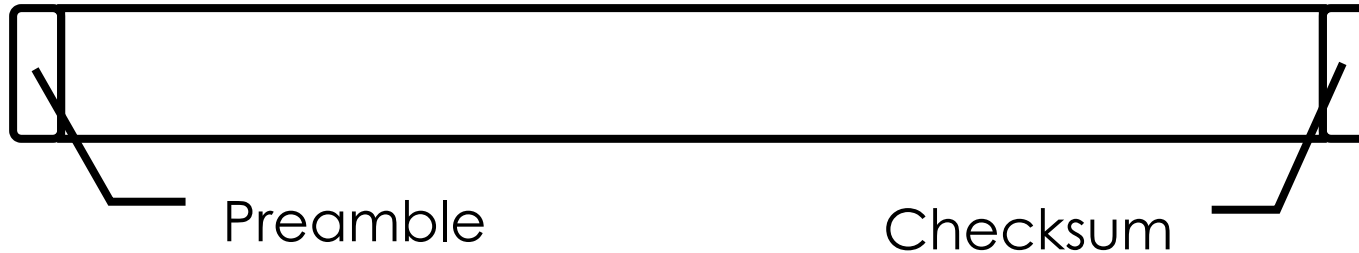
- Trace-driven experiments

- Evaluate end-to-end PP-ARQ performance
- Internet packet size distribution
- 802.11-size preambles

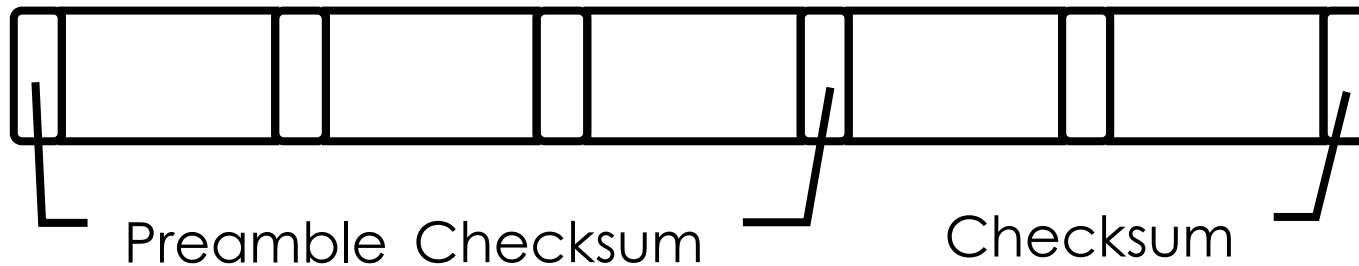
# PP-ARQ performance comparison

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- Packet CRC (no postamble)

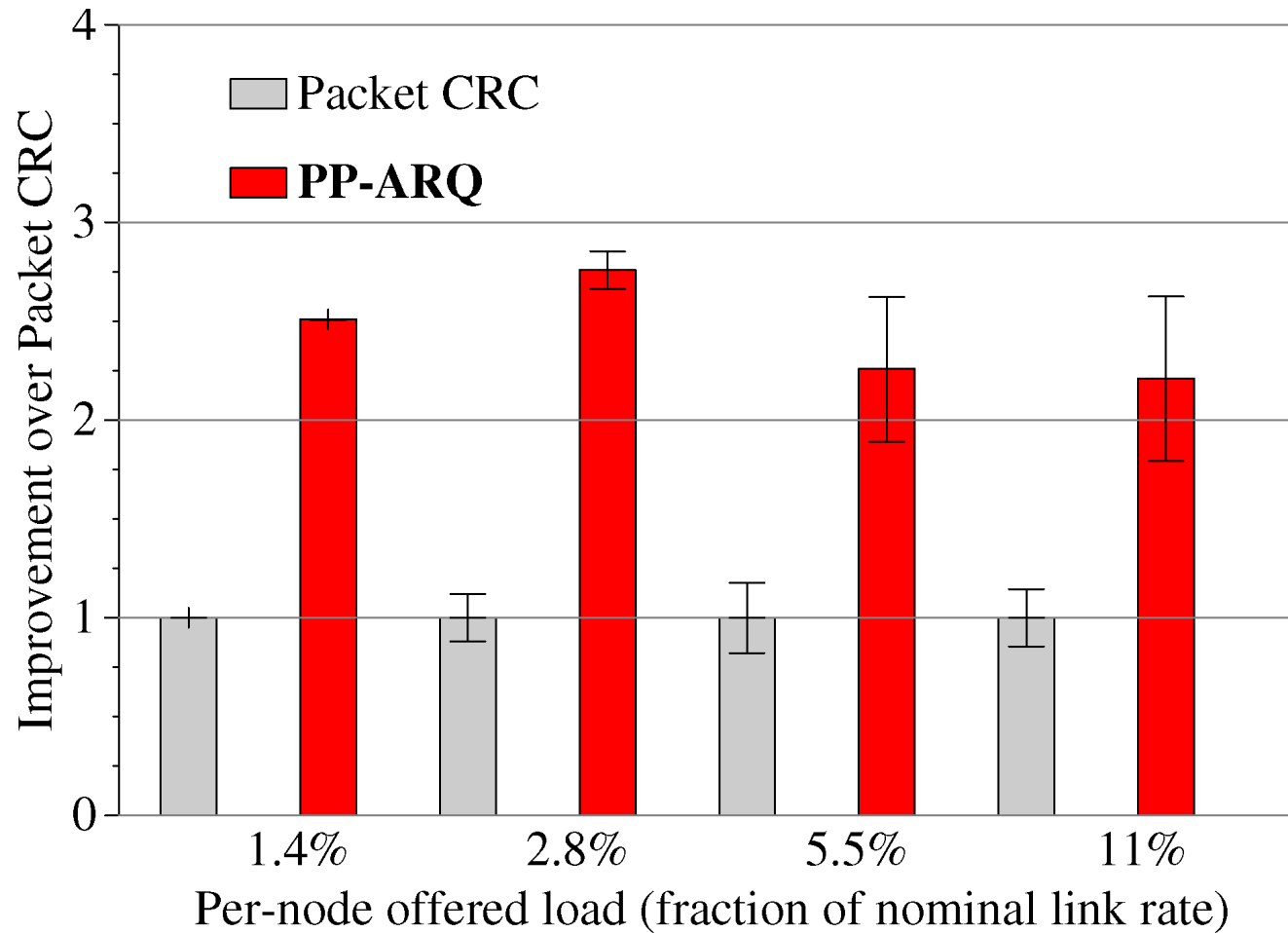


- Fragmented CRC (no postamble)
  - Tuned against traces for optimal fragment size

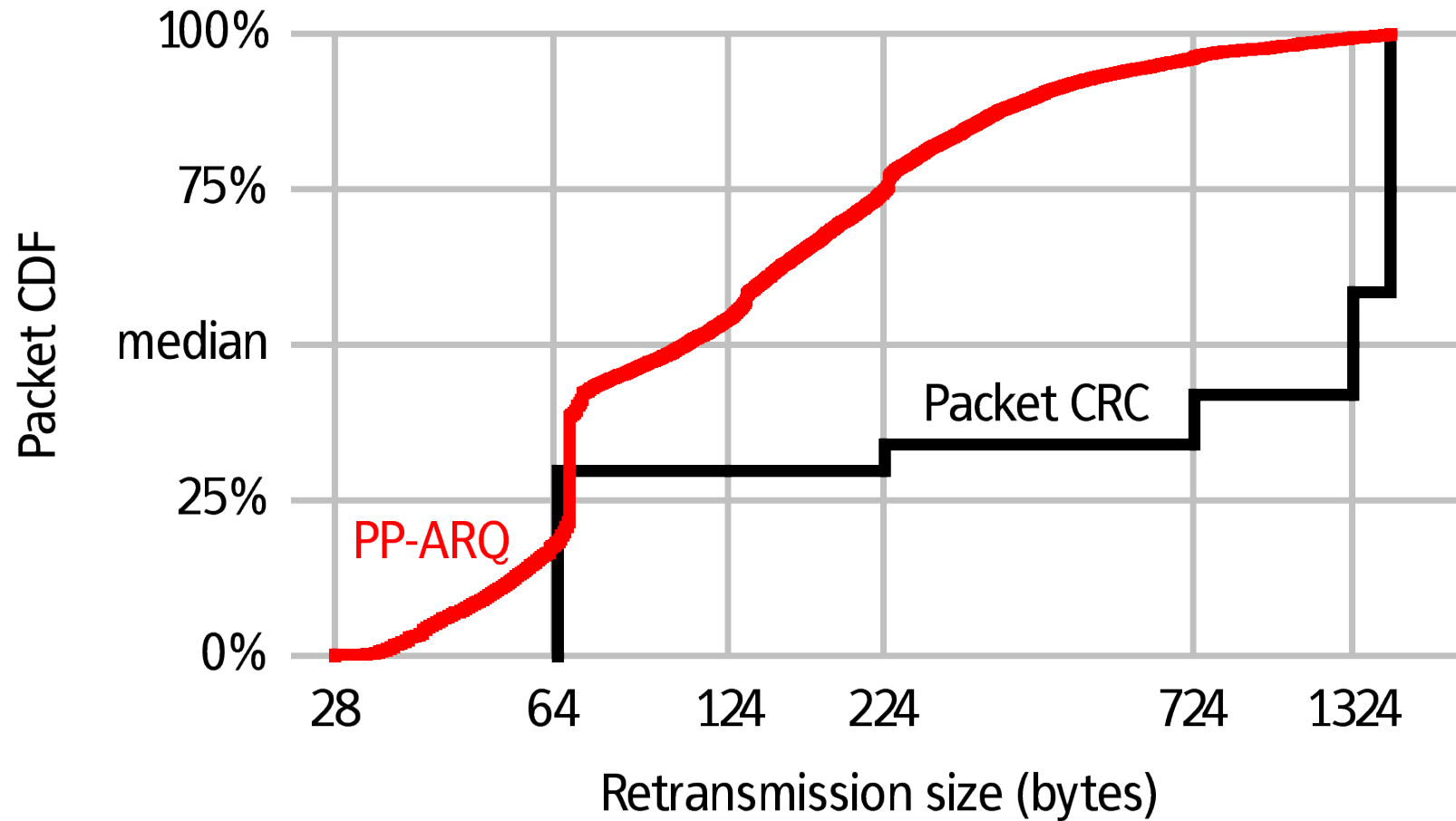


# Throughput Gain: 2.3-2.8x

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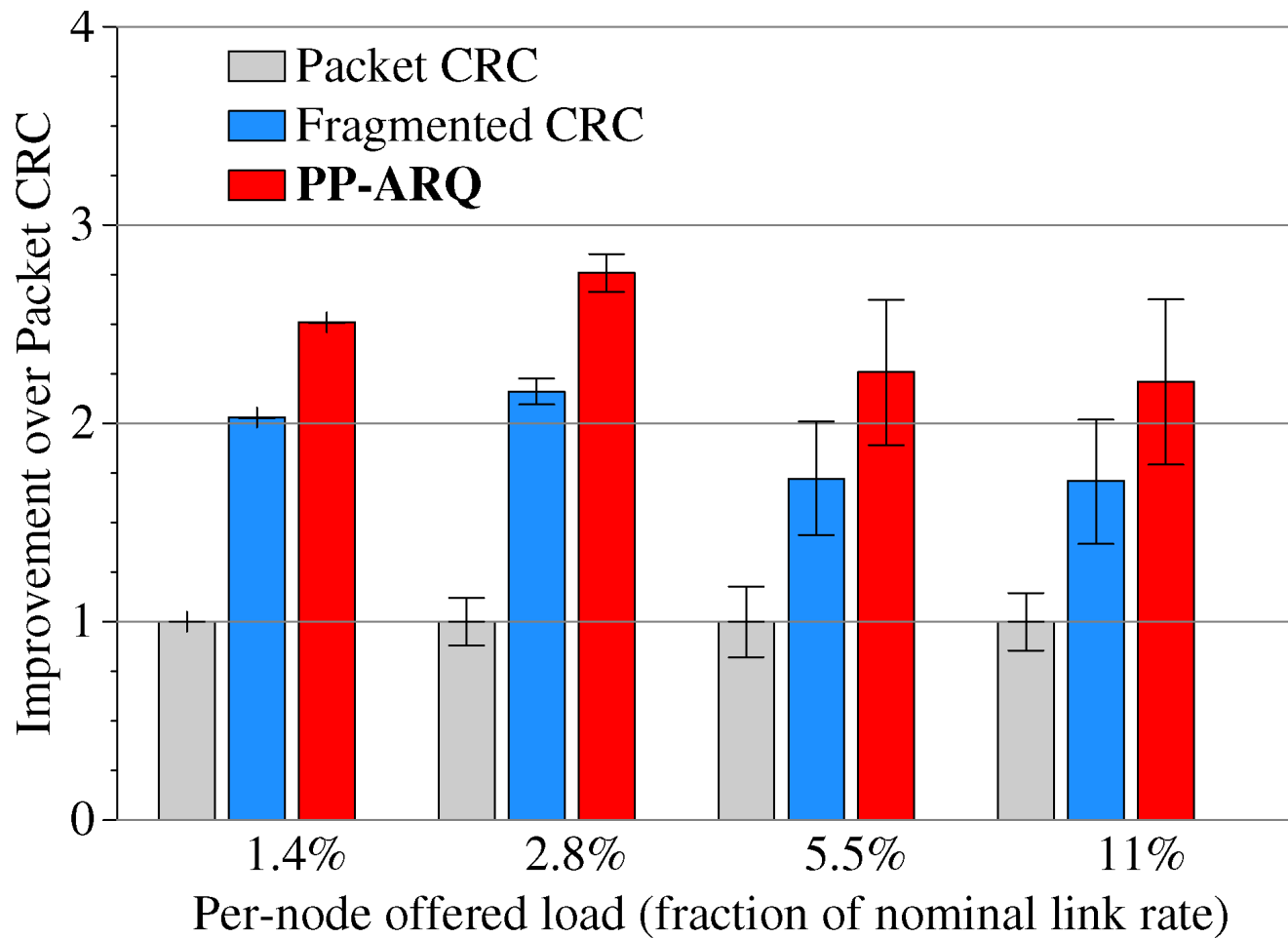


# PP-ARQ Retransmissions are Short

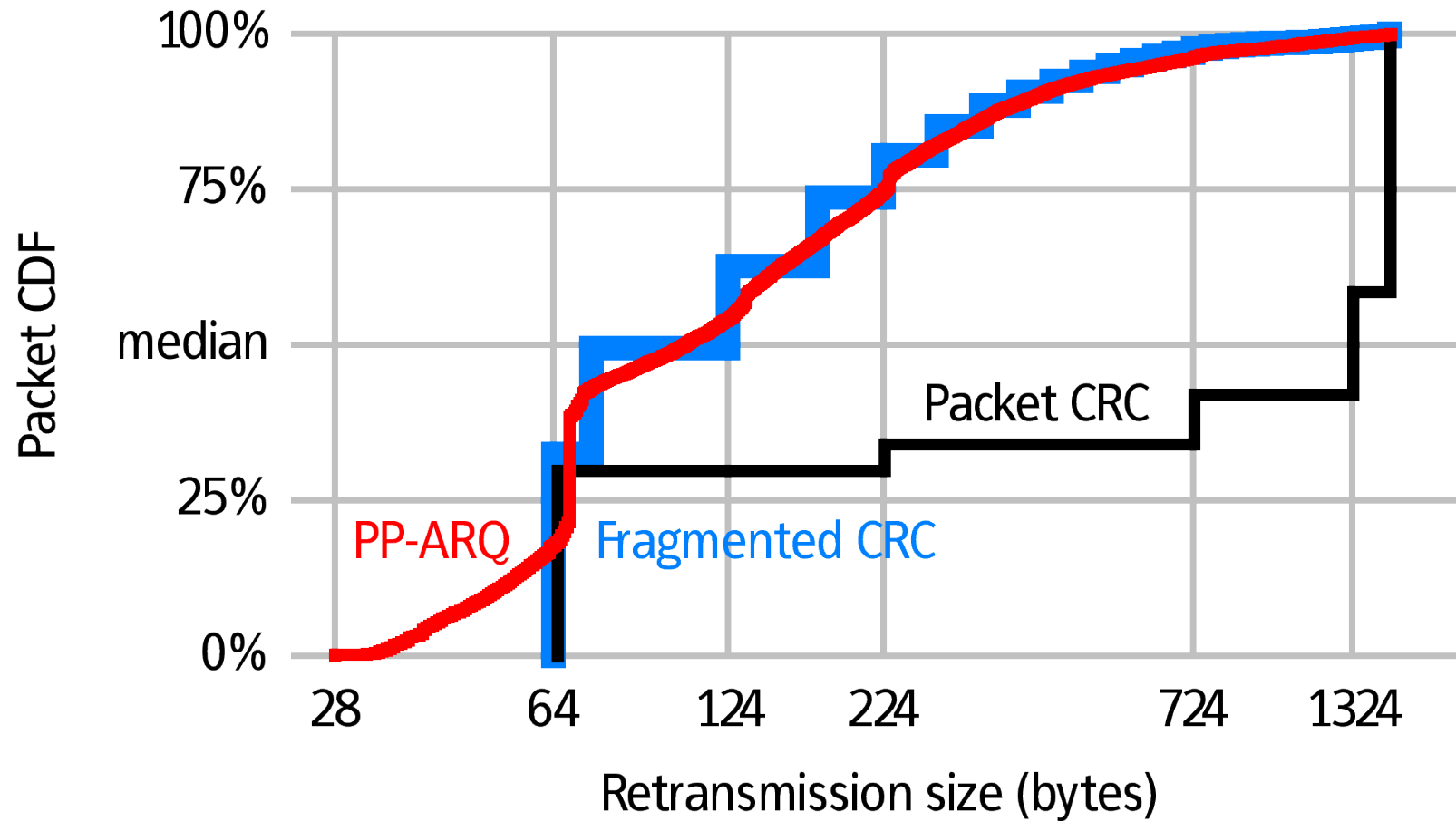


# 25% Gain over Fragmented

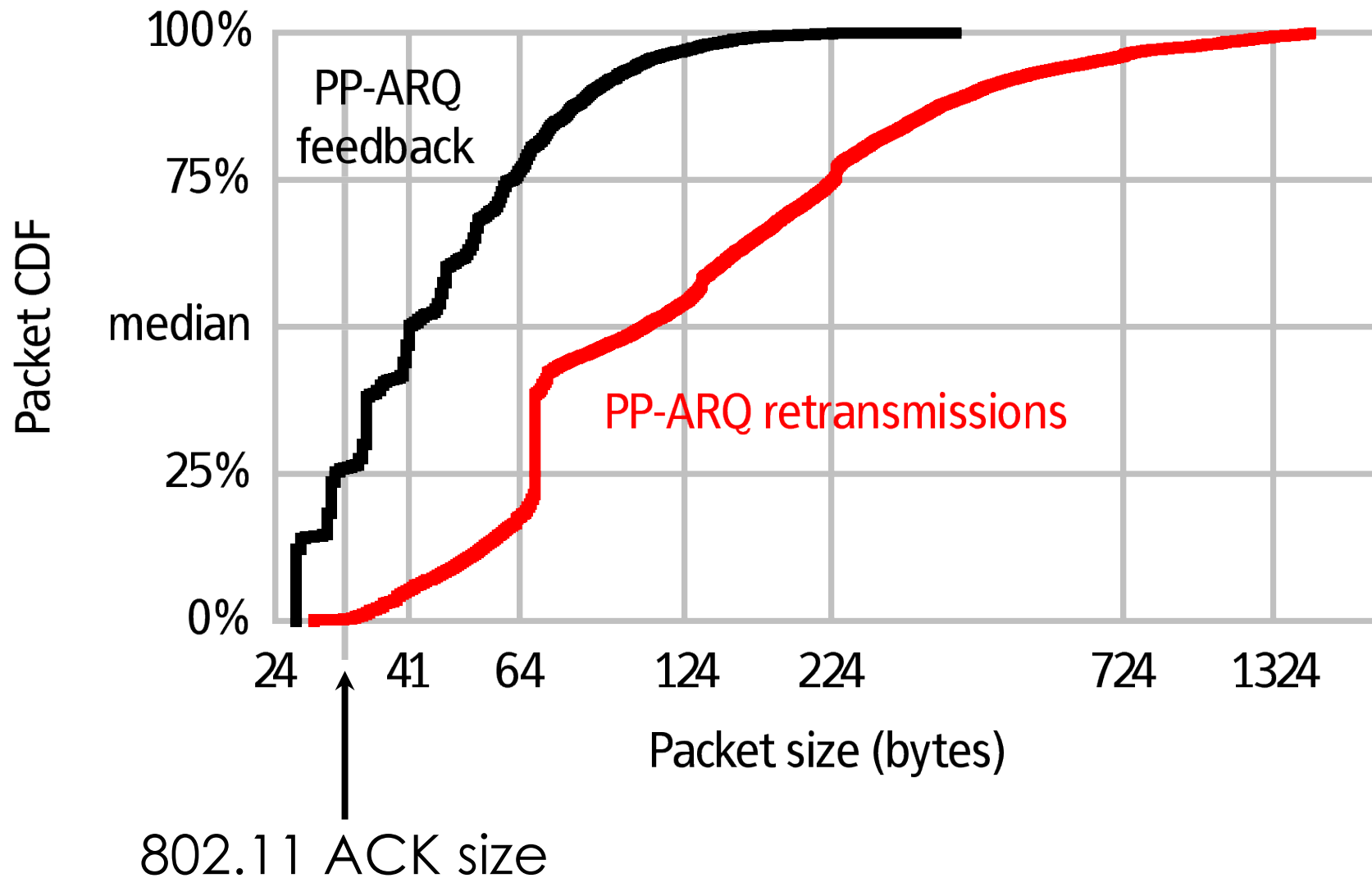
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# PP-ARQ Retransmissions are Short



# Low PP-ARQ Feedback Overhead





# Related work

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- ARQ with memory [Sindhu, IEEE Trans. On Comm. '77]
  - Incremental redundancy [Metzner, IEEE Trans. On Comm. '79]
  - Code combining [Chase, IEEE Trans. On Comm. '85]
- Combining retransmissions
  - **SPaC** [Dubois-Ferrière, Estrin, Vetterli; SenSys '05]
- Diversity combining
  - Reliability exchanging [Avudainayagam et al., IEEE WCNC '03]
  - **MRD** [Miu, Balakrishnan, Koksal; MobiCom '05]
  - **SOFT** [Woo et al.; MobiCom '07]
- Fragmented CRC
  - **Seda** [Ganti et al.; SenSys '06], 802.11 fragmentation

# Conclusion

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- Mechanisms for recovering correct bits from parts of packets
  - SoftPHY interface (PHY-independent)
  - Postamble decoding
- PP-ARQ improves throughput 2.3–2.8× over the status quo
- PPR Useful in other apps, e.g. opportunistic forwarding