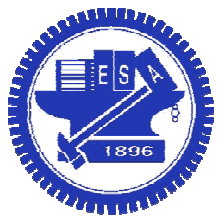


# Data Storage



National Chiao Tung University

Chun-Jen Tsai

2/24/2012

# Bits and Their Meaning

---

- ❑ First, we must consider how information can be stored inside computers
  - By “Information,” we mean numbers, text, images, sound, video, ..., etc.
- ❑ For today’s computers, information is encoded as patterns of 0’s and 1’s called bits (BInary digiTs)
  - The reason we use only two symbols (0 and 1) for information encoding is because it’s *simple*, not because it’s powerful.
  - *Non-binary computers* made of multi-level electrical components are possible but not popular yet.

# Digital Data Representation

---

- ❑ Some possible meanings for a single bit
  - Numeric value (1 or 0)
  - Boolean value (true or false)
  - Voltage (high or low)
- ❑ A bit can only represent one of two values, for more values, we need a long string of bits to represent them

# Boolean Operations

- ❑ Human beings are made of cells; Computers are made of small devices that can compute Boolean functions extremely fast
  - Boolean operations are operations that manipulate one or more 1/0 (or true/false) values
- ❑ Some examples of Boolean operations:

- AND:  $\frac{\text{AND } 0}{0}$   $\frac{\text{AND } 1}{0}$   $\frac{\text{AND } 0}{1}$   $\frac{\text{AND } 1}{1}$

- OR:  $\frac{\text{OR } 0}{0}$   $\frac{\text{OR } 1}{1}$   $\frac{\text{OR } 0}{1}$   $\frac{\text{OR } 1}{1}$

- XOR:  $\frac{\text{XOR } 0}{0}$   $\frac{\text{XOR } 1}{1}$   $\frac{\text{XOR } 0}{1}$   $\frac{\text{XOR } 1}{0}$

- NOT:  $\frac{\text{NOT } 0}{1}$   $\frac{\text{NOT } 1}{0}$

# Gates

- The small devices that compute Boolean operations are called “gates”
  - Often implemented as electronic circuits
  - The building blocks from which computers are constructed

**AND**



Inputs	Output
0 0	0
0 1	0
1 0	0
1 1	1

**OR**



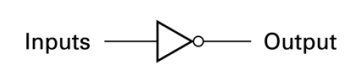
Inputs	Output
0 0	0
0 1	1
1 0	1
1 1	1

**XOR**



Inputs	Output
0 0	0
0 1	1
1 0	1
1 1	0

**NOT**

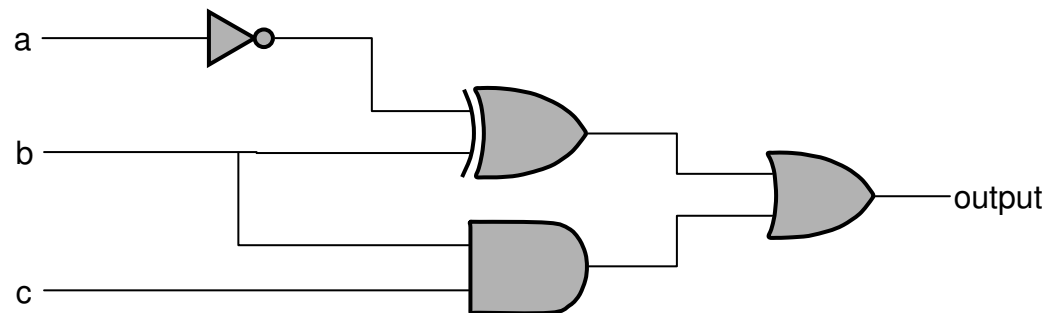


Inputs	Output
0	1
1	0

# Example of Simple Circuit

- A digital circuit's operation can be summarized by a *truth table*

**Circuit**

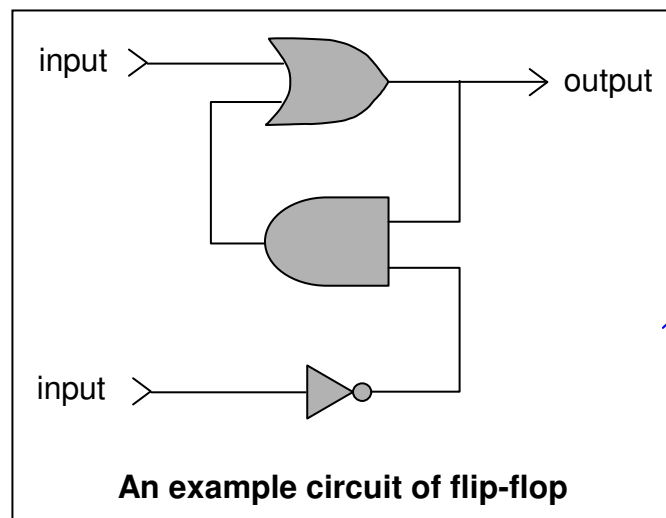


**Truth Table**

Input a, b, c	Output
000	1
001	1
010	0
011	1
100	0
101	0
110	1
111	1

# Flip-Flops

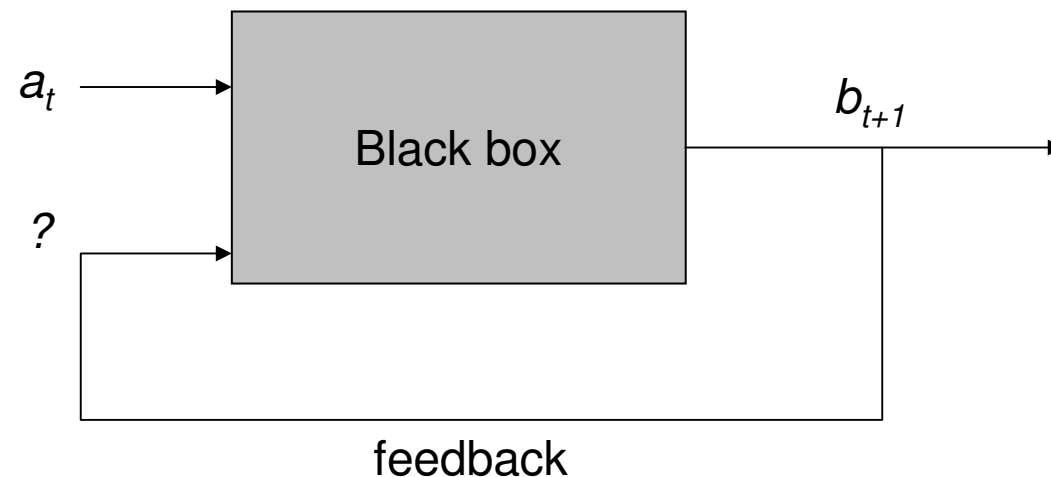
- ❑ In addition to performing Boolean operations, gates can also be used to store information
- ❑ A flip-flop is built from gates that stores one bit of data
  - Has an input line which sets its stored value to 1
  - Has an input line which sets its stored value to 0
  - When both inputs = 0, most recently stored value is preserved



Timing is important  
when analyzing the  
operation of a flip-flop!

# Note: Feedback Systems

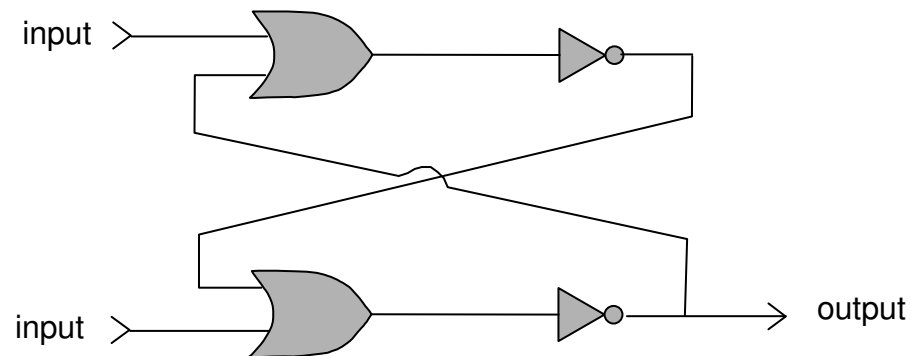
- To understand the operation of flip-flops, you must understand the concept of a feedback system:
  - At time  $t$ , we have two inputs to the “black box”, one is  $a_t$ , what is the other one?





# Alternative Flip-Flops

- There is more than one way to implement a flip-flop:



- Compare this to previous design, which one is better?

# Data Storage Hardware

---

- ❑ In addition to flip-flops, there are other fundamental bit-level data-storage devices based on, for example, magnetic or optical technologies
- ❑ A data storage device can be volatile or non-volatile:
  - Volatile memory – holds its value until the power is turned off – Example: flip-flops
  - Non-volatile memory – holds its value after the power is off – Example: magnetic storage

# Radix-N Number Coding

---

- ❑ Bit numbers are coded in radix-2 (a.k.a. base-2) while human-friendly numbers are coded in radix-10
- ❑ The conversion between radix-2 and radix-10 is not trivial, sometimes, it is easier to use a number coding system that is:
  - Easily readable by human
  - Can be converted to/from radix-2 easily (why?)
- ❑ Popular radix-N notation that fulfill these two points are radix-8 (octal notation) and radix-16 (hexadecimal notation)

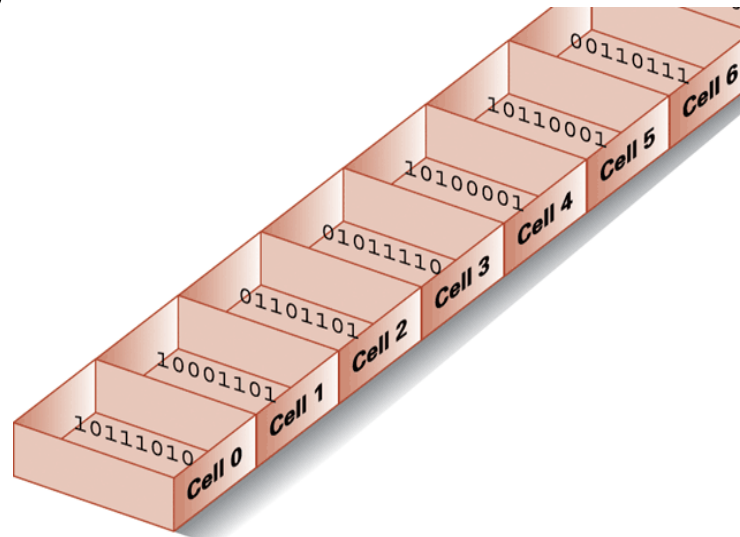
# Different Coding Systems

Binary	Octal	Hexadecimal	Decimal
0000	0	0	0
0001	1	1	1
0010	2	2	2
0011	3	3	3
0100	4	4	4
0101	5	5	5
0110	6	6	6
0111	7	7	7
1000	10	8	8
1001	11	9	9
1010	12	A	10
1011	13	B	11
1100	14	C	12
1101	15	D	13
1110	16	E	14
1111	17	F	15



# Main Memory Addresses

- ❑ An *address* is a “name” to uniquely identify one **cell** in the computer’s main memory
  - The addresses for cells in a computer are consecutive numbers, usually starting at zero
- ❑ A Random Access Memory (RAM) is a memory device where any cell can be accessed independently



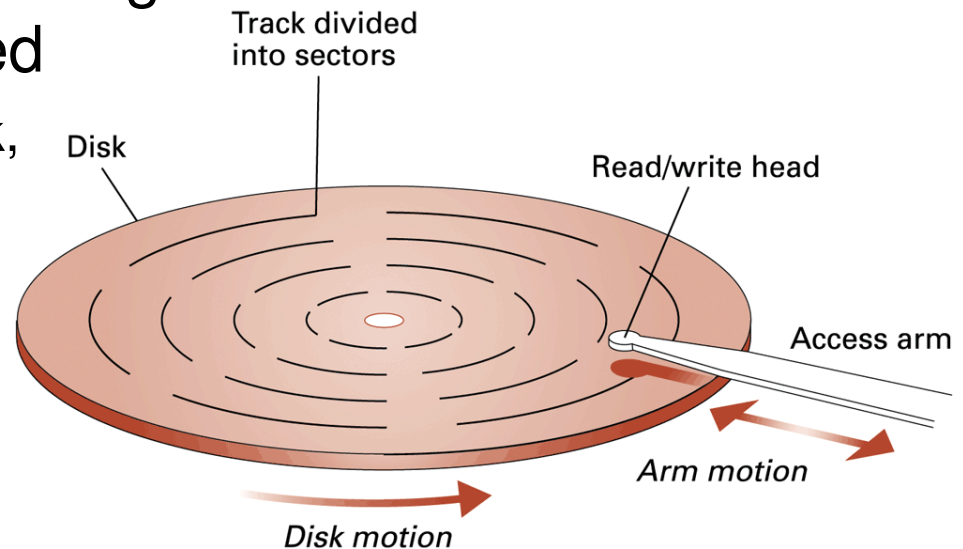
# Memory Capacity Measure

- ❑ For computer memory/storage, the unit prefixes are slightly different from those of the metric system:
  - “kB” means 1,000 byte; but KB (or more precisely, KiB) means  $2^{10} = 1024$  bytes
  - “MB” normally means 1,000,000 byte; but MiB means  $2^{20} = 1,048,576$  bytes
  - Note: the unit KiB (kibi), MiB (mebi), GiB (gibi), etc., were adopted by IEC in 1999.
- ❑ For any other computer related units, we stick only to the metric prefix system:
  - The bandwidth unit Mbps always means  $10^6$  bits-per-second

# Mass Storage Systems (1/3)

□ For mass storage, rotating disks are usually used

- Hard disk, floppy disk, CD-ROM



■ Characteristics:

- Seek time – the time to move the head to the right track
- Latency time – the time to rotate the disk for half a cycle
- Access time – seek time + latency time
- Transfer rate – the speed data transferred to or from the disk



# Mass Storage Systems (2/3)

- ❑ Back in 1956, disk storage are huge†



IBM 305 Computer System

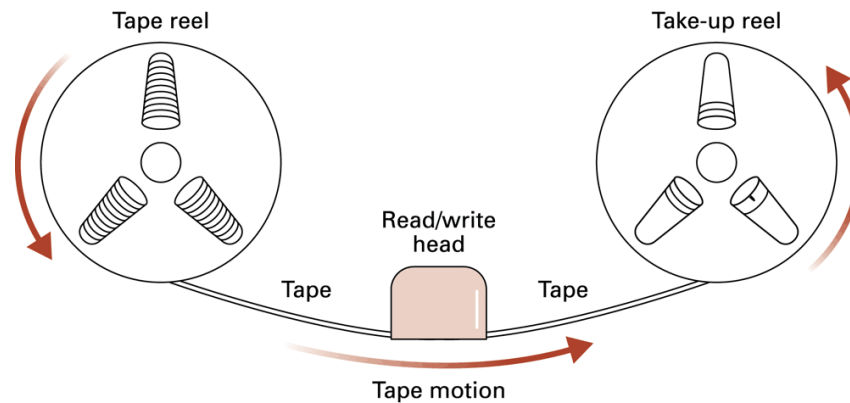


IBM 350 Disk Storage Unit  
(1200 rpm, 5MB storage)

† <http://www.ibm.com/ibm/history>

# Mass Storage Systems (3/3)

- ❑ Some “old” mass storage devices use magnetic tapes:



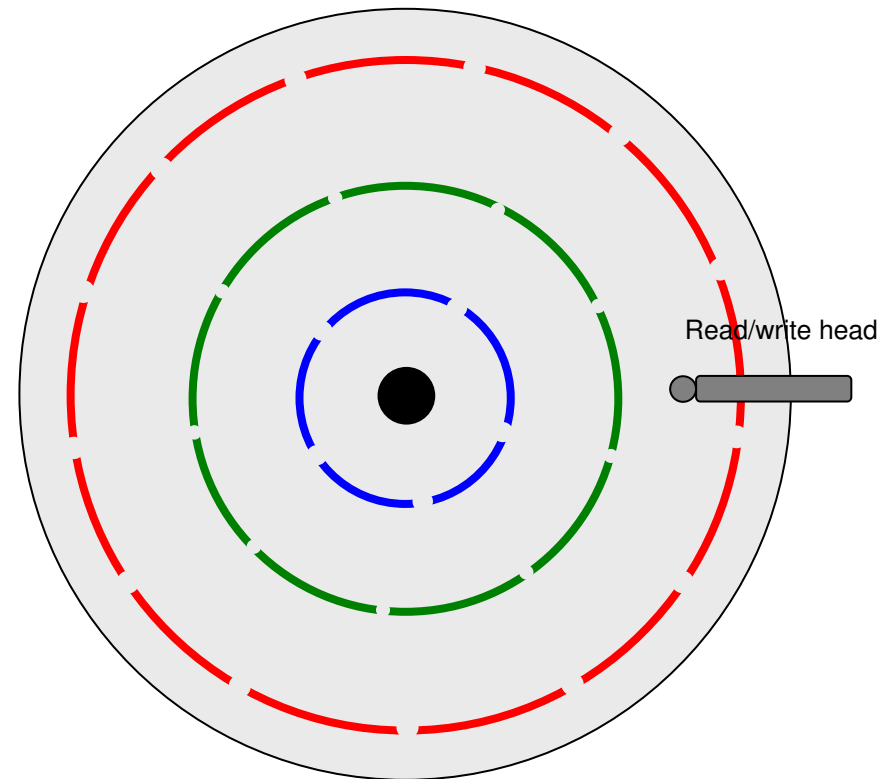
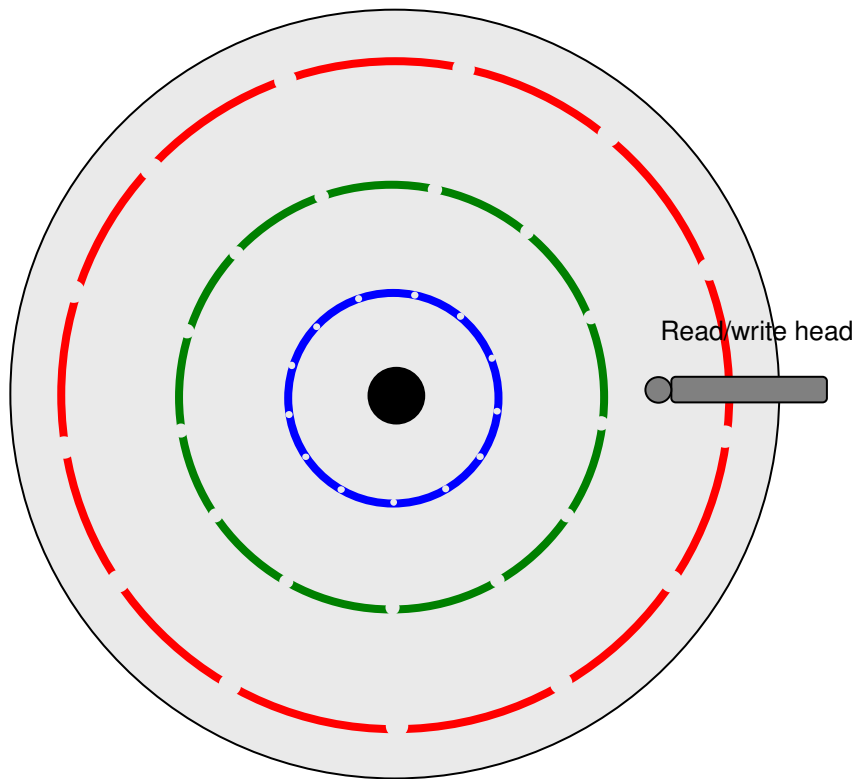
- ❑ Some “newer” mass storages use flash memory:



# Disk Rotating Speed & Sector Size

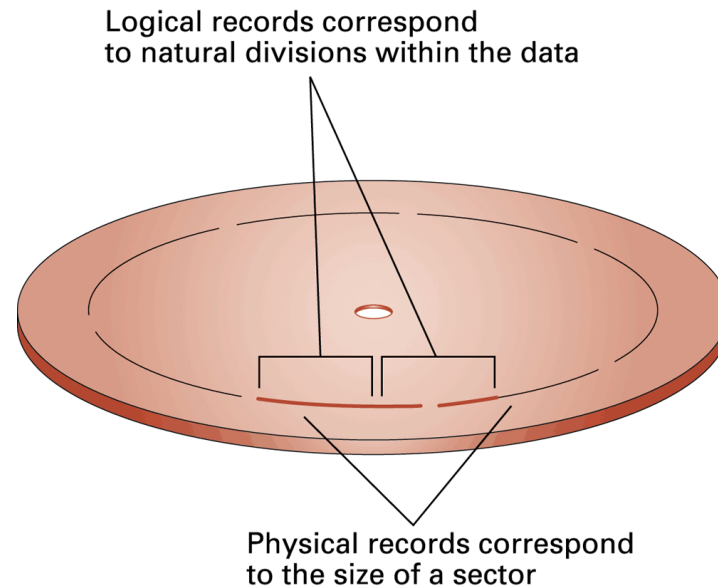
Constant angular speed

Variable angular speed



# Logical Record vs. Physical Record

- ❑ Information stored in a mass storage is organized into files; each file is composed of many smaller units, called logical records
- ❑ The mapping between logical record and physical record is not one-to-one:



# Representing Text

- For text-representation, each character (letter, punctuation, etc.) is assigned a unique bit pattern.

There are several coding standards:

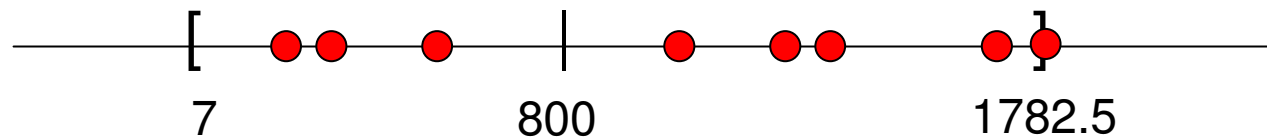
- ASCII: 7-bit coding for most symbols used in written English

01001000	01100101	01101100	01101100	01101111	00101110
H	e	l	l	o	.

- Extended ASCII (ISO 8859-1): 8-bit coding, including more western language symbols
- Unicode: 16-bit coding for most symbols used in most world languages today
- ISO 10646 Universal Character Set (UCS): a text coding system which uses 32-bit values for each characters

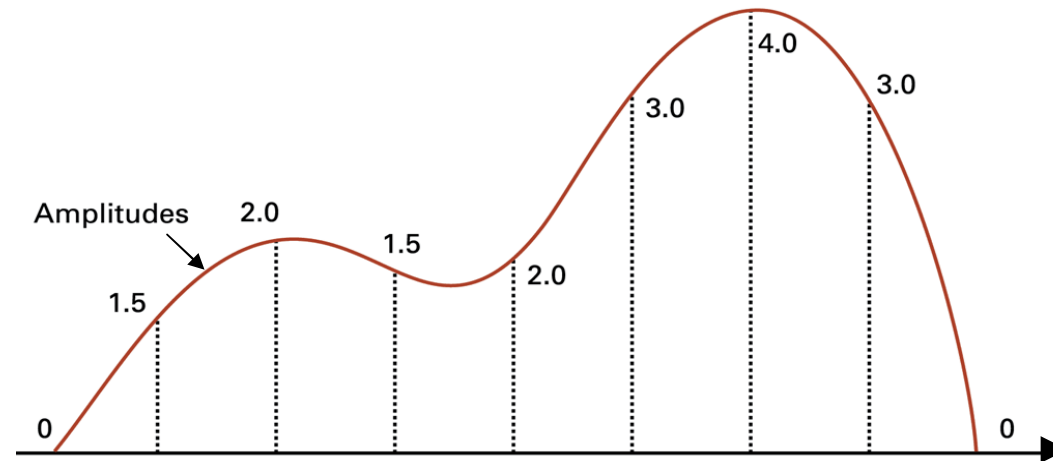
# Representing Numeric Values

- ❑ There are some limitations of computer representations of numeric values
  - *Overflow* happens if a number is too big to be represented
  - *Underflow* happens when the number is too small
  - *Truncation* happens when a number is fallen between two representable numbers
- ❑ For example, 3-bit encoding of numbers can only represent 8 different values; each red dot in the following interval is a numerical values we can represent *accurately*



# Representing Natural Phenomena

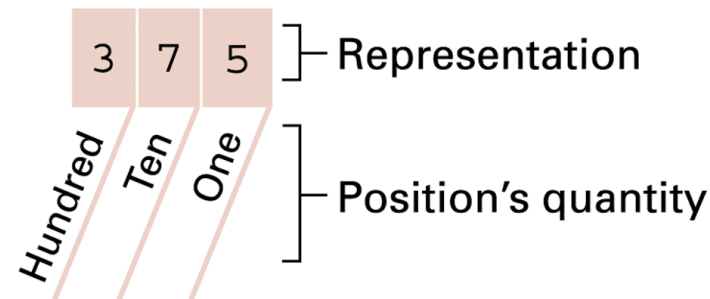
- ❑ Natural phenomena is usually detected in *analog* form; must be converted to *digital* form for computer to store and/or process
- ❑ Example: a sound wave can be represented by a sequence of numbers: 0, 1.5, 2.0, 1.5, 2.0, 3.0, 4.0, 3.0, 0



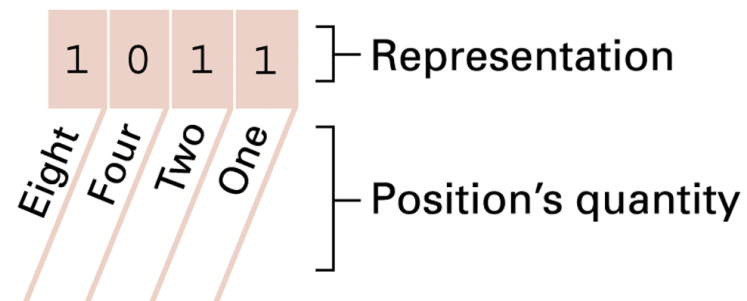
# Radix Conversion

- ❑ The conversion of number from one radix (base) to another is important
- ❑ Each digit in a number has a base quantity called position's quantity
- ❑ Position's quantity is the power of the base

a. Base ten system



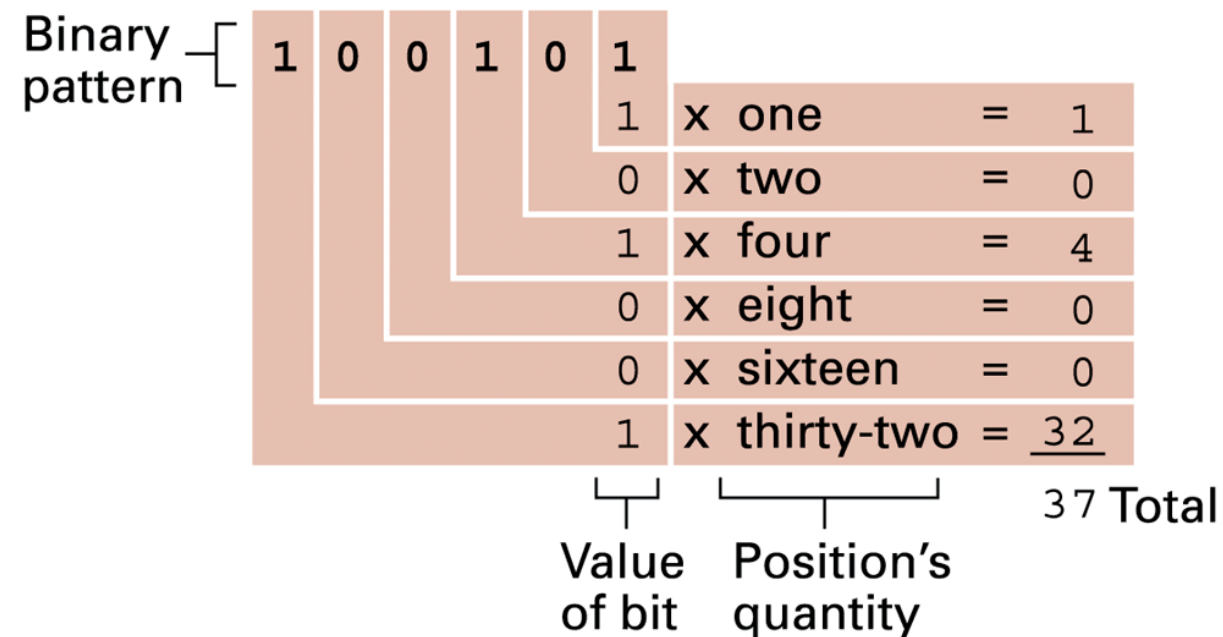
b. Base two system





# Conversion from Radix-2 to Radix-10

□ For example, 100101 can be decoded as follows:



# Conversion from Radix-10 to Radix-2

- Method to convert from Radix-10 to radix-2:
  - Divide the value by 2 and record the remainder
  - Continue to divide the quotient by two and record the remainder until the quotient is zero
- For example, convert  $(13)_{10}$  to radix-2

$$\begin{array}{l} 13 = 2 \times 6 + 1 \\ 6 = 2 \times 3 + 0 \\ 3 = 2 \times 1 + 1 \\ 1 = 2 \times 0 + 1 \end{array} \longrightarrow (13)_{10} = (1101)_2$$

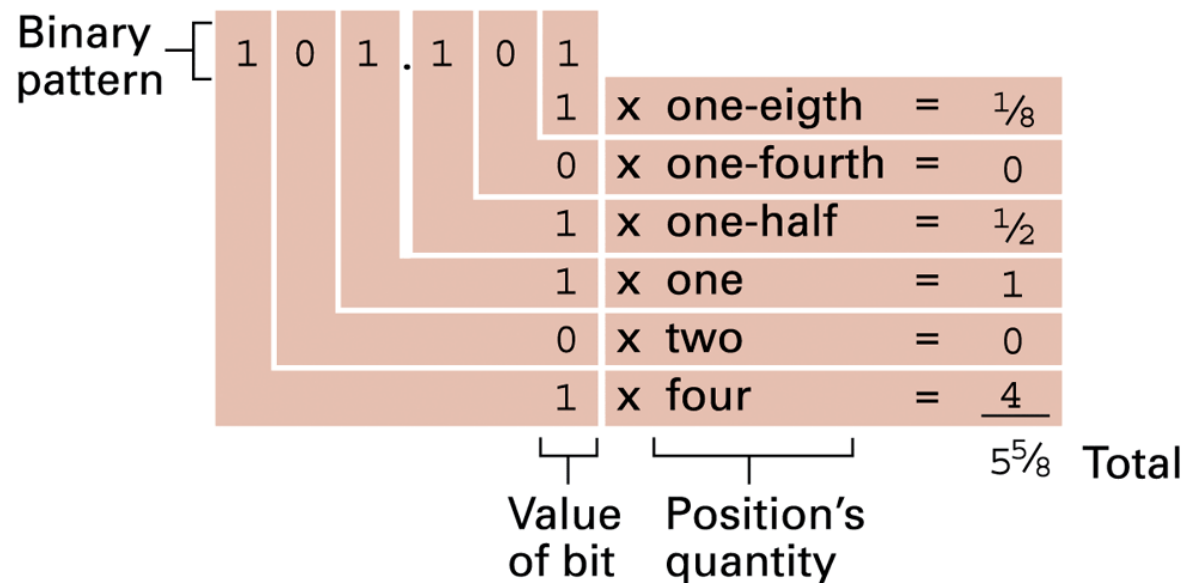
# Binary Arithmetic

- ❑ Modern computers usually use binary representation of numerical data, therefore, numerical calculations must be done in radix-2 as well
- ❑ Binary calculations are simple, but remembering a long sequence of binary numbers is the difficult part (good thing that computers have photographic memory):

$$\begin{array}{r} 0 \\ +0 \\ \hline 0 \end{array} \quad \begin{array}{r} 1 \\ +0 \\ \hline 1 \end{array} \quad \begin{array}{r} 0 \\ +1 \\ \hline 1 \end{array} \quad \begin{array}{r} 1 \\ +1 \\ \hline 10 \end{array}$$

# Binary Representation of Fractions

- A simple way to represent numbers with fractions is to use a fixed point representation:



# Converting Fractions to Binary Rep.

- To convert  $\frac{5}{8}$  to binary digits 0.101, you can inverse the above-mentioned process:

$$\frac{5}{8} \times 2 = \frac{5}{4} = \textcircled{1} + \frac{1}{4} \quad \rightarrow \quad \frac{5}{8} = 1 \times 2^{-1} + 2^{-1} \times \frac{1}{4}$$

$$\frac{1}{4} \times 2 = \frac{1}{2} = \textcircled{0} + \frac{1}{2} \quad \rightarrow \quad \frac{1}{4} = 0 \times 2^{-1} + 2^{-1} \times \frac{1}{2}$$

$$\rightarrow \quad \frac{5}{8} = 1 \times 2^{-1} + 2^{-1} \times (0 \times 2^{-1} + 2^{-1} \times \frac{1}{2})$$

$$\rightarrow \quad \frac{5}{8} = \textcircled{1} \times 2^{-1} + \textcircled{0} \times 2^{-2} + \textcircled{1} \times 2^{-3}$$

# Representing Integers

---

- ❑ Unsigned integers can be represented in base two
- ❑ Signed integers are numbers that can be positive or negative. For negative numbers, there are several ways to represent it:
  - One's complement –  
bitwise “NOT” of a positive number is the negative representation of the number
  - Two's complement notation –  
the most popular representation of negative numbers
  - Excess notation –  
often used in representing floating point exponents

# Two's Complement Representation

a. Using patterns of length three

Bit pattern	Value represented
011	3
010	2
001	1
000	0
111	-1
110	-2
101	-3
100	-4

b. Using patterns of length four

Bit pattern	Value represented
0111	7
0110	6
0101	5
0100	4
0011	3
0010	2
0001	1
0000	0
1111	-1
1110	-2
1101	-3
1100	-4
1011	-5
1010	-6
1001	-7
1000	-8

# Coding Rules of 2's Complement

## □ Algorithm:

- First, complement all the digits of the number
- Then, add one to the complemented number

## □ For example, $-6$ can be represented as follows:

- $(6)_{10} = (0110)_2$
- $0110 \rightarrow 1001$
- $1001 + 1 = 1010 = (-6)_{10}$

## □ The nice part about 2's Complement representation of negative numbers is that addition matches natural representation



# 2's Complement Arithmetic

Problem in base ten		Problem in two's complement		Answer in base ten
$\begin{array}{r} 3 \\ + 2 \\ \hline \end{array}$	→	$\begin{array}{r} 0011 \\ + 0010 \\ \hline 0101 \end{array}$	→	5
$\begin{array}{r} -3 \\ + -2 \\ \hline \end{array}$	→	$\begin{array}{r} 1101 \\ + 1110 \\ \hline 1011 \end{array}$	→	-5
$\begin{array}{r} 7 \\ + -5 \\ \hline \end{array}$	→	$\begin{array}{r} 0111 \\ + 1011 \\ \hline 0010 \end{array}$	→	2

# Excess Notation Systems

- ❑ Excess notation preserves the natural ordering of (negative) numbers
- ❑ Arithmetic operations become less trivial
- ❑ Examples of negative numbers represented using excess system:

Bit pattern	Value represented
111	3
110	2
101	1
100	0
011	-1
010	-2
001	-3
000	-4

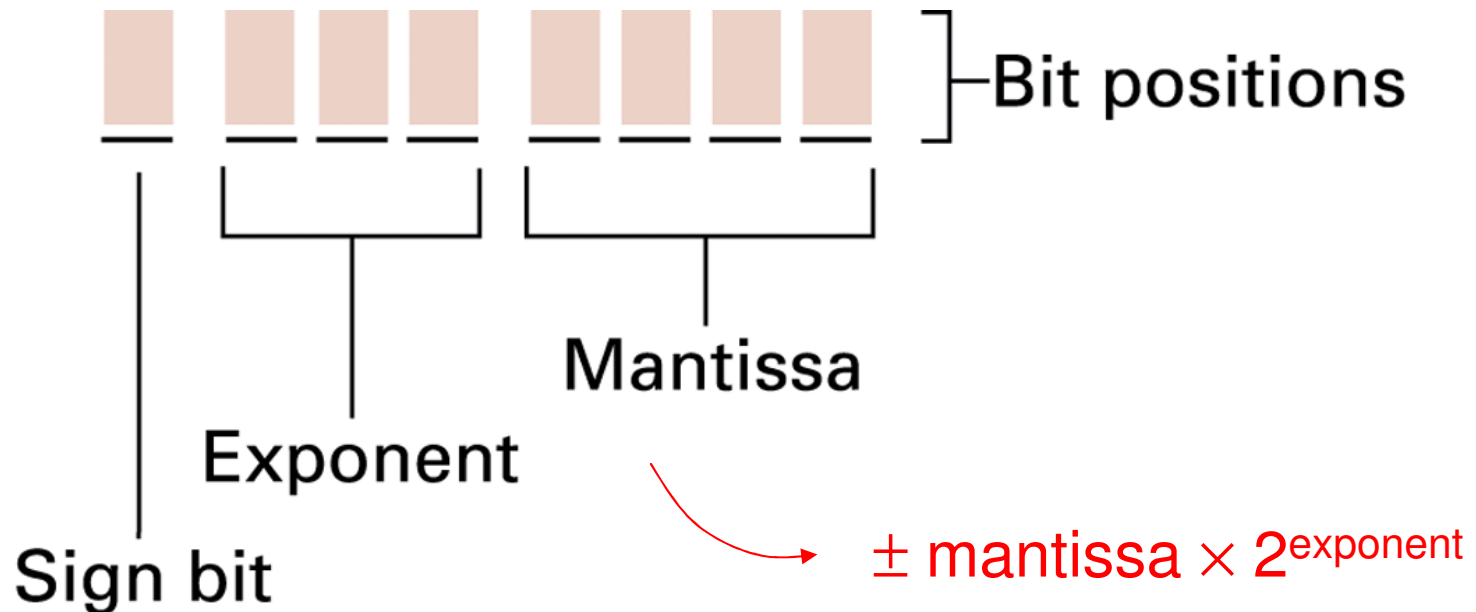
excess 3

Bit pattern	Value represented
1111	7
1110	6
1101	5
1100	4
1011	3
1010	2
1001	1
1000	0
0111	-1
0110	-2
0101	-3
0100	-4
0011	-5
0010	-6
0001	-7
0000	-8

excess 8

# Floating Point Representation (1/2)

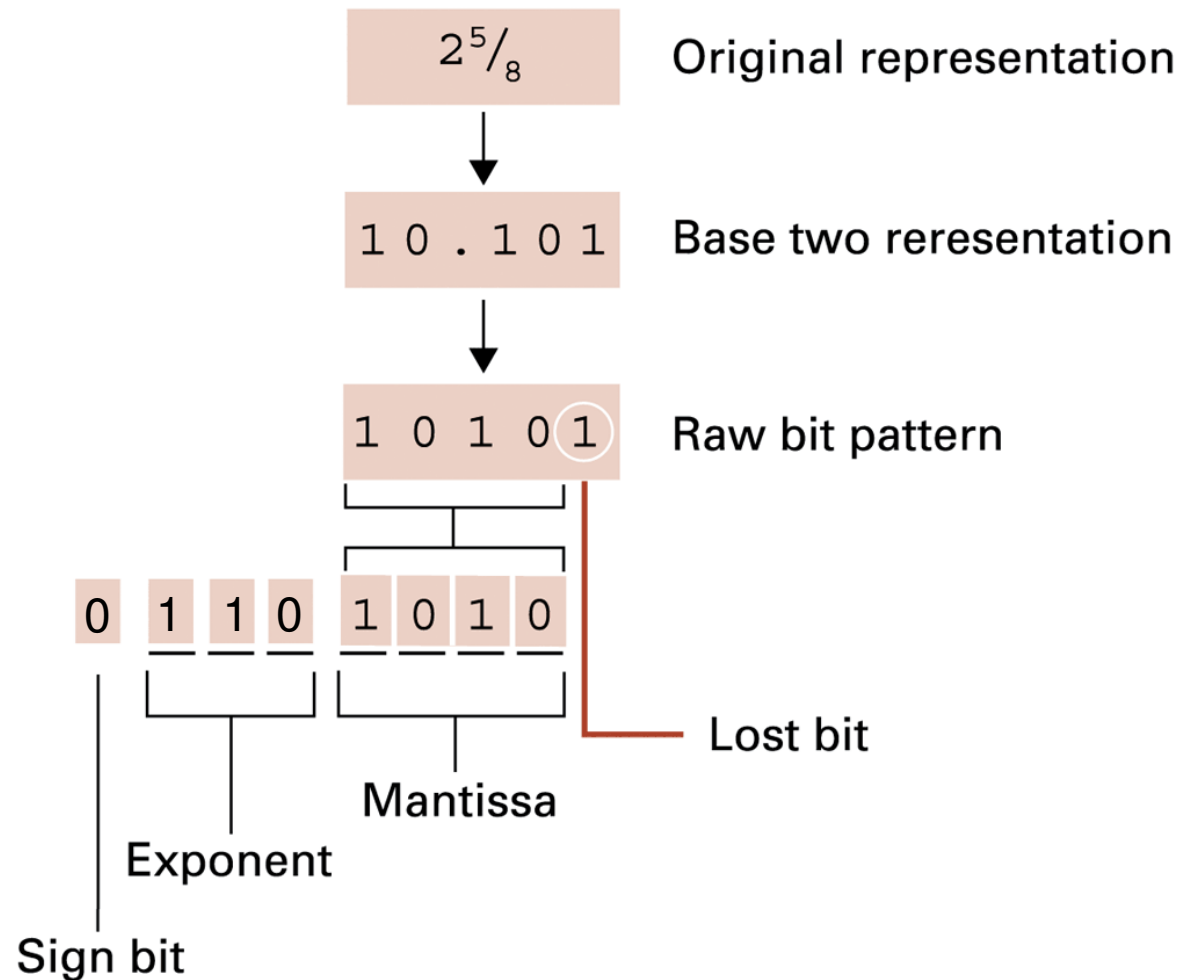
- ❑ A number with fractions can be represented in fixed point, as well as in floating point
- ❑ A floating point representation is composed of four parts:



# Floating Point Representation (2/2)

- ❑ Coding of each field of a floating point number:
  - For mantissa, fixed point representation is obviously the logical choice (remember scientific notations?)
  - For exponent, we need to represent positive and negative numbers → excess notation is often used (why?)
  - For sign bit, a single bit can be used to record the sign
- ❑ There is an international standard, IEEE-754, that defines the representation as well as the arithmetic of floating point computations
  - It is different (more complicated) from the one described in our textbook, but the concept is the same

# Example: Coding of $2\frac{5}{8}$



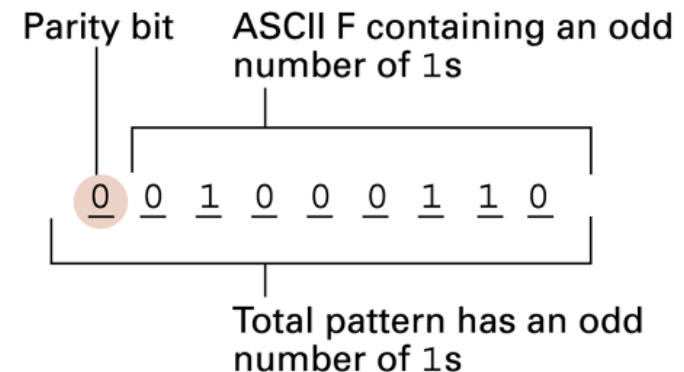
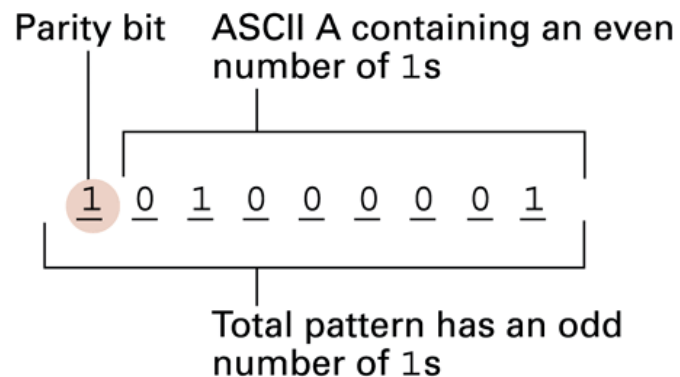
# Coding for Error Protection

---

- ❑ Sometimes, computing systems are not reliable. Numbers stored in a computing system can have random bit errors over time.
- ❑ To detect and/or correct errors, the number representation has to be robust to random bit errors
- ❑ Two types of error-robust coding are possible:
  - Error detection coding (e.g. parity bits)
  - Error correcting coding (add redundant bits to increase robustness)

# Parity Bit Coding

- ❑ Each number can be appended with one extra bit (called parity bit) for error detection
  - The extra bit force the total bit pattern has an odd or even number of 1's.
  - Odd parity – the number of 1's is odd
  - Even parity – the number of 1's is even



# Error Correction Coding (ECC)

- ❑ For ECC, we don't use all possible representations, instead, we increase the "distance" between meaningful representations:

Symbol	Code
A	000000
B	001111
C	010011
D	011100
E	100110
F	101001
G	110101
H	111010

distance = 4

distance = 4

distance = 4

distance = 3



# Example: Decoding of 010100

Character	Distance between the received pattern and the character being considered
A	2
B	4
C	3
D	1 <i>Smallest distance</i>
E	3
F	5
G	2
H	4